目录

[头文件 3](#_Toc23321)

[杂物 4](#_Toc9344)

[并查集(维护链) 4](#_Toc25683)

[读入挂 4](#_Toc10803)

[其他挂 5](#_Toc31801)

[平板电视 5](#_Toc16424)

[Dancing Links 6](#_Toc6020)

[快速乘法(就那个long double的) 9](#_Toc29798)

[快速乘法(mod加速的) 9](#_Toc3981)

[极小mex区间(r-l=max-min) 10](#_Toc22692)

[set暴力修改区间 10](#_Toc24037)

[螺旋数列value 11](#_Toc23555)

[DP 12](#_Toc27062)

[决策单调性优化 12](#_Toc21347)

[斜率优化 13](#_Toc7482)

[四边形不等式优化 13](#_Toc11660)

[wqs二分(带权二分) 14](#_Toc27587)

[数位DP 15](#_Toc31113)

[树形依赖背包 15](#_Toc29993)

[DP套DP 16](#_Toc22752)

[插头DP 17](#_Toc6564)

[斯坦纳树, 子集卷积的计数DP 21](#_Toc1169)

[字符串 23](#_Toc10572)

[KMP|最小表示法 23](#_Toc769)

[字典树 23](#_Toc29921)

[AC自动机 23](#_Toc24119)

[AC自动机 另一种写法 24](#_Toc29061)

[后缀数组 24](#_Toc25882)

[后缀自动机 25](#_Toc28556)

[马拉车 29](#_Toc6037)

[回文自动机 29](#_Toc9362)

[二分hash 31](#_Toc7361)

[一些hashset|hashmap 32](#_Toc2932)

[后缀平衡树 33](#_Toc23436)

[子序列自动机 34](#_Toc28897)

[区间border 35](#_Toc20857)

[数据结构 40](#_Toc31131)

[时间分治 40](#_Toc24677)

[二维树状数组 41](#_Toc6053)

[树状数组 不大于k的最大值 42](#_Toc21371)

[BIT\_差分 42](#_Toc19946)

[二维线段树 42](#_Toc7352)

[吉爷爷线段树 43](#_Toc24105)

[扫描线 矩形周长并 44](#_Toc6070)

[主席树 45](#_Toc6948)

[可持久化数组(主席树维护) 45](#_Toc17293)

[树套树 46](#_Toc10596)

[CDQ分治(套线段树) 47](#_Toc13588)

[SPLAY 48](#_Toc2841)

[SPLAY启发式合并 50](#_Toc19394)

[左偏树 51](#_Toc30671)

[可持久化非旋treap 52](#_Toc18175)

[LCT 53](#_Toc18204)

[KD树 56](#_Toc1709)

[莫队 58](#_Toc11635)

[树上莫队(套分块) 58](#_Toc22820)

[回滚莫队套分块 58](#_Toc17590)

[带修改莫队 60](#_Toc2561)

[二次离线莫队 61](#_Toc26182)

[用set维护凸包 62](#_Toc16502)

[李超树 62](#_Toc19100)

[线性基(套路) 63](#_Toc28918)

[手写BITSET 63](#_Toc17079)

[杨表 64](#_Toc1398)

[图论 65](#_Toc16863)

[二分图匹配 65](#_Toc14500)

[Hall定理 65](#_Toc25224)

[KM 二分图最大权匹配 66](#_Toc17330)

[最短路 67](#_Toc10083)

[差分约束系统 67](#_Toc15280)

[01分数规划 67](#_Toc476)

[切比雪夫(曼哈顿)距离最小生成树 68](#_Toc8844)

[笛卡尔树 69](#_Toc12690)

[强连通分量tarjan 70](#_Toc4395)

[支配树 70](#_Toc27793)

[边双连通分量 仙人掌图 72](#_Toc19261)

[环套外向树 73](#_Toc32337)

[网络流 74](#_Toc21175)

[无向图全局最小割 76](#_Toc15311)

[无向图最小割树 GH-tree 77](#_Toc26)

[最小费用流 77](#_Toc23859)

[上下界网络流 79](#_Toc3564)

[树分治 80](#_Toc4577)

[动态点分治 80](#_Toc14613)

[部分树上dp 82](#_Toc2541)

[2-sat 82](#_Toc25850)

[可持久化的2-sat 输出方案 83](#_Toc18500)

[dfs序\_换根的讨论 85](#_Toc27014)

[01序列转树 86](#_Toc27403)

[有方向的树链剖分 86](#_Toc8847)

[轻重儿子分开维护 87](#_Toc17406)

[链分治, 动态维护树上dp 89](#_Toc21292)

[DSU on tree 90](#_Toc10465)

[LCA 91](#_Toc17659)

[倍增 92](#_Toc25649)

[虚树ST表求lca；四毛子定理 92](#_Toc2670)

[Ladder长链剖分 k级祖先 94](#_Toc12483)

[最大团 95](#_Toc17554)

[最小树形图(mlogn) 95](#_Toc16112)

[一般图最大匹配 带花树 98](#_Toc4071)

[树分块 高度分块 99](#_Toc30841)

[树哈希 100](#_Toc6481)

[kruskal树 102](#_Toc8273)

[数学相关 104](#_Toc8994)

[牛顿迭代 开根 104](#_Toc5613)

[逆元, kummer等基础 104](#_Toc17349)

[BSGS 104](#_Toc32708)

[Pell方程 104](#_Toc12962)

[博弈：NIM,SG 104](#_Toc32551)

[Exgcd 105](#_Toc23849)

[K次方和, 伯努利数 105](#_Toc24483)

[求原根 二次三次剩余(无板子) 106](#_Toc3366)

[常系数线性递推 106](#_Toc1659)

[多项式暴力求积分 107](#_Toc29391)

[五边形数定理 108](#_Toc19210)

[FFT 109](#_Toc14379)

[多项式单log无穷背包(MTT) 109](#_Toc26369)

[多项式开根求逆,除法取模(NTT) 111](#_Toc12481)

[fwt, 子集卷积 113](#_Toc26088)

[子集和(高维前缀和, 分治) 114](#_Toc4017)

[子集卷积 114](#_Toc30022)

[高斯消元 115](#_Toc8103)

[矩阵树定理|拉格朗日插值 116](#_Toc17238)

[Polya定理| Burnside引理 117](#_Toc4930)

[Miller\_Rabin素性测试+pollard\_rho因数分解 117](#_Toc18497)

[中国剩余定理(不一定互质) 118](#_Toc31211)

[广义容斥 119](#_Toc30248)

[Prime-counting function 119](#_Toc1073)

[Min\_25筛 119](#_Toc28297)

[积性函数 前缀和 杜教筛 120](#_Toc23580)

[类欧几里得 122](#_Toc4034)

[欧拉降幂公式 123](#_Toc14624)

[单纯形法 123](#_Toc965)

[其他的东西 126](#_Toc21648)

[杜教线性递推BM板子 126](#_Toc302)

[任意模数BM板子 127](#_Toc19140)

[自适应 simpson积分 130](#_Toc14947)

[杜教多项式插值 130](#_Toc5710)

[求x^2+y^2=n的(x,y)对数 131](#_Toc4715)

头文件

#include <sstream>

#include <fstream>

#include <cstdio>

#include <iostream>

#include <algorithm>

#include <vector>

#include <set>

#include <map>

#include <string>

#include <cstring>

#include <stack>

#include <queue>

#include <cmath>

#include <ctime>

#include <utility>

#include <cassert>

#include <bitset>

**using** **namespace** std;

#define REP(I,N) for (I=0;I<N;I++)

#define rREP(I,N) for (I=N-1;I>=0;I--)

#define rep(I,S,N) for (I=S;I<N;I++)

#define rrep(I,S,N) for (I=N-1;I>=S;I--)

#define FOR(I,S,N) for (I=S;I<=N;I++)

#define rFOR(I,S,N) for (I=N;I>=S;I--)

#define DEBUG

#ifdef DEBUG

#define debug(...) fprintf(stderr, \_\_VA\_ARGS\_\_)

#define deputs(str) fprintf(stderr, "%s\n",str)

#else

#define debug(...)

#define deputs(str)

#endif *// DEBUG*

**typedef** unsigned long long ULL;

**typedef** unsigned long long ull;

**typedef** unsigned int ui;

**typedef** long long LL;

**typedef** long long ll;

**typedef** pair<int,int> pii;

**typedef** pair<ll,ll> pll;

**const** int INF=0x3f3f3f3f;

**const** LL INFF=0x3f3f3f3f3f3f3f3fll;

**const** LL M=1e9+7;

**const** LL maxn=1e6+107;

**const** double pi=acos(-1.0);

**const** double eps=0.0000000001;

LL gcd(LL a, LL b) {**return** b?gcd(b,a%b):a;}

**template**<**typename** T>**inline** void pr2(T x,int k=64) {ll i; REP(i,k) debug("%d",(x>>i)&1); putchar(' ');}

**template**<**typename** T>**inline** void add\_(T &A,int B,ll MOD=M) {A+=B; (A>=MOD) &&(A-=MOD);}

**template**<**typename** T>**inline** void mul\_(T &A,ll B,ll MOD=M) {A=(A\*B)%MOD;}

**template**<**typename** T>**inline** void mod\_(T &A,ll MOD=M) {A%=MOD; A+=MOD; A%=MOD;}

**template**<**typename** T>**inline** void max\_(T &A,T B) {(A<B) &&(A=B);}

**template**<**typename** T>**inline** void min\_(T &A,T B) {(A>B) &&(A=B);}

**template**<**typename** T>**inline** T abs(T a) {**return** a>0?a:-a;}

**inline** ll powMM(ll a, ll b, ll mod=M) {

ll ret=1;

**for** (; b; b>>=1ll,a=a\*a%mod)

**if** (b&1) ret=ret\*a%mod;

**return** ret;

}

int startTime;

void startTimer() {startTime=clock();}

void printTimer() {debug("/--- Time: %ld milliseconds ---/**\n**",clock()-startTime);}

# 杂物

**首先是没啥用的两个板子**

void msort(int le,int ri) {*//逆序对*

**if** (le==ri) **return**;

int mid=(le+ri)>>1,i=le,j=mid+1,k=i;

msort(le,mid); msort(j,ri);

**while** (i<=mid||j<=ri) {

**if** (i==mid+1) {b[k++]=a[j++]; ans+=mid-i+1;}

**else** **if** (j==ri+1) b[k++]=a[i++];

**else** **if** (a[i]<=a[j]) b[k++]=a[i++];

**else** {b[k++]=a[j++]; ans+=mid-i+1;}

}

**for** (i=le; i<=ri; i++) a[i]=b[i];

}

void fqsort(int l,int r) {*//O(n)第k大数*

int le=l,ri=r,m;

m=a[le];

**while** (le<ri) {

**while** (le<ri&&a[ri]<=m) ri--; a[le]=a[ri];

**while** (le<ri&&a[le]>=m) le++; a[ri]=a[le];

} **if** (le==k) printf("%d**\n**",m);

**else** **if** (le>k) fqsort(l,le-1);

**else** fqsort(le+1,r);

}

****并查集(维护链)****

**inline** int getfa(int x){

**if** (fa[x]==x) **return** x;

int y=getfa(fa[x]);

**if** (fa[x]!=y) sum[x]+=sum[fa[x]];

fa[x]=y;

**return** y;

}

读入挂

**namespace** fastIO {*//感觉没问题, 测试几次*

#define BUF\_SIZE 100000

**namespace** Istream {

bool IOerror = 0;

**inline** char ic() {

**static** char buf[BUF\_SIZE],\*p1=buf+BUF\_SIZE,\*pend=buf+BUF\_SIZE;

**if** (p1==pend) {

p1=buf;

pend=buf+fread(buf,1,BUF\_SIZE,stdin);

**if** (pend == p1) {IOerror = 1; **return** -1;}

} **return** \*p1++;

}

**inline** bool blank(char ch) {

**return** ch == ' ' || ch == '\n' || ch == '\r' || ch == '\t';

}

**template**<**typename** T>

**inline** void readPositive(T &x) {*//no*

char ch; x=0;

**while** (blank(ch=ic()));

**if** (IOerror) **return**;

**for** (x=0; '0'<=ch&&ch<='9'; ch=ic()) x=x\*10+ch-'0';

}

**template**<**typename** T>

**inline** void read(T &x) {

char ch; T op=1; x=0;

**while** (blank(ch=ic()));

**if** (IOerror) **return**;

**if** (ch=='-') op=-1,ch=ic();

**for** (x=0; '0'<=ch&&ch<='9'; ch=ic()) x=x\*10+ch-'0';

x\*=op;

}

**inline** void read(char &c) {

c=ic();

}

**inline** void read(char \*s) { *//len*

char ch;

**while** (blank(ch=ic()));

**if** (IOerror) **return**;

**for** (; !blank(ch)&&!IOerror; ch=ic()) \*s++=ch;

\*s='\0';

}

}

**namespace** Ostream {

char buf[BUF\_SIZE], \*p1 = buf, \*pend = buf + BUF\_SIZE;

**inline** void flush() {

fwrite(buf,1,p1-buf,stdout);

p1=buf;

}

**inline** void oc(char ch) {

**if** (p1 == pend) flush();

\*p1++=ch;

}

**inline** void println() {

oc('\n');

}

**template**<**typename** T>

**inline** void print(T x) {

**static** char s[27],\*s1=s;

**if** (!x) \*s1++='0';

**if** (x<0) oc('-'),x=-x;

**while** (x) \*s1++=x%10+'0',x/=10;

**do** {s1--; oc(\*s1);} **while** (s1!=s);

}

**inline** void print(char s) {

oc(s);

}

**inline** void print(char \*s) {

**for** (; \*s; oc(\*s++));

}

**inline** void print(**const** char \*s) {

**for** (; \*s; oc(\*s++));

}

**inline** void print(string s) {

**for** (unsigned i=0; i<s.length(); i++) oc(s[i]);

}

**struct** \_flush {

~\_flush() {flush();}

} fflush;

};

**template**<**typename** T>

**inline** void read(T &x) {Istream::readPositive(x);}

**inline** void read(char \*x) {Istream::read(x);}

**template**<**typename** T>

**inline** void print(T x) {Ostream::print(x);}

**template**<**typename** T>

**inline** void println(T x) {print(x); Ostream::oc('\n');}

}

**using** **namespace** fastIO;

其他挂

扩栈挂

#ifdef OPENSTACK

int size = 256 << 20; *// 256MB*

char \*p = (char\*)malloc(size) + size;

#if (defined \_WIN64) or (defined \_\_unix)

\_\_asm\_\_("movq %0, %%rsp**\n**" :: "r"(p));

#else

\_\_asm\_\_("movl %0, %%esp**\n**" :: "r"(p));

#endif

#endif

注意最后加exit(0);

玄学挂

#pragma comment(linker, "/stack:200000000")

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm,mmx,avx,tune=native")

然后加上并行计算(计组)

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC target("avx")

平板电视

1. **红黑树**

#include<ext/pb\_ds/assoc\_container.hpp>

#include<ext/pb\_ds/tree\_policy.hpp>

**using** **namespace** std;

**using** **namespace** \_\_gnu\_cxx;

**using** **namespace** \_\_gnu\_pbds;

**typedef** tree<int,null\_type,less<int>,rb\_tree\_tag,tree\_order\_statistics\_node\_update> rbtree;

*/\**

*定义一颗红黑树*

*int 关键字类型*

*null\_type无映射(低版本g++为null\_mapped\_type)*

*less<int>从小到大排序*

*rb\_tree\_tag 红黑树（splay\_tree\_tag）*

*tree\_order\_statistics\_node\_update结点更新*

*插入t.insert();*

*删除t.erase();*

*Rank:t.order\_of\_key();*

*第K值:t.find\_by\_order();*

*前驱:t.lower\_bound();*

*后继t.upper\_bound();*

*a.join(b)b并入a 前提是两棵树的key的取值范围不相交*

*a.split(v,b)key小于等于v的元素属于a，其余的属于b*

*T.lower\_bound(x) >=x的min的迭代器*

*T.upper\_bound((x) >x的min的迭代器*

*T.find\_by\_order(k) 有k个数比它小的数*

*\*/*

rbtree T;

rbtree::iterator it;

1. **Rope**

**using** **namespace** std;

**using** **namespace** \_\_gnu\_cxx;

/\*

1）运算符：rope支持operator += -= + - < ==

2）输入输出：可以用<<运算符由输入输出流读入或输出。

3）长度/大小：调用length()，size()都可以哦

4）插入/添加等：

append(const string&)

substr(start,length)

push\_back(x);//在末尾添加x

insert(pos,x);//在pos插入x，自然支持整个char数组的一次插入

erase(pos,x);//从pos开始删除x个

copy(pos,len,x);//从pos开始到pos+len为止用x代替

replace(pos,x);//从pos开始换成x

substr(pos,x);//提取pos开始x个

at(x)/[x];//访问第x个元素

\*/

rope<int> V;

1. **二项堆(这里是dijkstra)**

#include<iostream>

#include<cstdio>

#include<cstring>

#include<ext/pb\_ds/priority\_queue.hpp>

#define ll long long

#define pa pair<ll,int>

#define llinf 9000000000000000000LL

**using** **namespace** std;

**using** **namespace** \_\_gnu\_pbds;

**typedef** \_\_gnu\_pbds::priority\_queue<pa,greater<pa>,pairing\_heap\_tag > heap;

int n,m,cnt,last[1000005];

int T,rxa,rxc,rya,ryc,rp;

heap::point\_iterator id[1000005];

int x,y,z;

ll dis[1000005];

**struct** data {int to,next,v;} e[10000005];

**inline** int read() {

int x=0,f=1; char ch=getchar();

**while** (ch<'0'||ch>'9') {**if** (ch=='-')f=-1; ch=getchar();}

**while** (ch>='0'&&ch<='9') {x=x\*10+ch-'0'; ch=getchar();}

**return** x\*f;

}

void insert(int u,int v,int w) {

e[++cnt].to=v; e[cnt].next=last[u]; last[u]=cnt; e[cnt].v=w;

}

void dijkstra() {

heap q;

**for** (int i=1; i<=n; i++)dis[i]=llinf;

dis[1]=0; id[1]=q.push(make\_pair(0,1));

**while** (!q.empty()) {

int now=q.top().second; q.pop();

**for** (int i=last[now]; i; i=e[i].next)

**if** (e[i].v+dis[now]<dis[e[i].to]) {

dis[e[i].to]=e[i].v+dis[now];

**if** (id[e[i].to]!=0)

q.modify(id[e[i].to],make\_pair(dis[e[i].to],e[i].to));

**else** id[e[i].to]=q.push(make\_pair(dis[e[i].to],e[i].to));

}

}

}

Dancing Links

1. **不可重复**

*//dlx:求解精确覆盖*

*//link的意思是,r覆盖了c*

*//暴力枚举,n个点覆盖m个格子; 注意一定要init*

**struct** DLX {

**const** **static** int maxn=1e5+7;

**const** **static** int maxd=1e4+7;

int n,m,size;

int U[maxn],D[maxn],R[maxn],L[maxn],col[maxn],row[maxn];

int H[maxd],S[maxd];*//S:cnt*

int ans[maxn];

void init(int \_n,int \_m) {

n=\_n; m=\_m; int i;

FOR(i,0,m) {

S[i]=0;

U[i]=D[i]=i;

L[i]=i-1,R[i]=i+1;

} R[m]=0; L[0]=m;

size=m;

FOR(i,0,n) H[i]=-1;

}

void link(int r,int c) {

S[col[++size]=c]++; row[size]=r;

D[size]=D[c]; U[D[c]]=size;

D[c]=size; U[size]=c;

**if** (H[r]<0) H[r]=L[size]=R[size]=size;

**else** {

R[size]=R[H[r]];

L[R[H[r]]]=size;

L[size]=H[r];

R[H[r]]=size;

}

}

void remove(int c) {

L[R[c]]=L[c]; R[L[c]]=R[c];

**for** (int i=D[c]; i!=c; i=D[i])

**for** (int j=R[i]; j!=i; j=R[j])

U[D[j]]=U[j],D[U[j]]=D[j],S[col[j]]--;

}

void resume(int c) {

**for** (int i=U[c]; i!=c; i=U[i])

**for** (int j=L[i]; j!=i; j=L[j])

U[D[j]]=D[U[j]]=j,S[col[j]]++;

L[R[c]]=R[L[c]]=c;

}

*//这里是找可行解; 最优解无法加估价函数剪枝*

char g[maxn];

bool dance(int pos) {

**if** (R[0]==0) {

int i,j;

REP(i,pos) g[(ans[i]-1)/16]=(ans[i]-1)%16+'A';

REP(i,16) {REP(j,16) putchar(g[i\*16+j]); puts("");}

**return** 1;

}

*// cnt,pos:选择的个数; ans:选择的值(列)*

*// if (pos>=cnt&&cnt!=INF) return;*

*// if (R[0]==0) {cnt=min(cnt,pos); return;}*

int c=R[0];

**for** (int i=R[0]; i; i=R[i])

**if** (S[i]<S[c]) c=i;

remove(c);

**for** (int i=D[c]; i!=c; i=D[i]) {

ans[pos]=row[i];

**for** (int j=R[i]; j!=i; j=R[j]) remove(col[j]);

**if** (dance(pos+1)) **return** 1;

**for** (int j=L[i]; j!=i; j=L[j]) resume(col[j]);

} resume(c);

**return** 0;

}

} dlx;

char g[27][27];

int n,m;

void add(int x,int y,int k) {

int r=(x\*16+y)\*16+k;

dlx.link(r,16\*16\*0+x\*16+y+1);*//position*

dlx.link(r,16\*16\*1+x\*16+k);*//行*

dlx.link(r,16\*16\*2+y\*16+k);*//列*

dlx.link(r,16\*16\*3+(x/4\*4+y/4)\*16+k);*//块*

}

int main() {

int i,j,k;

**while** (~scanf("%s",g[0])) {

rep(i,1,16) scanf("%s",g[i]);

dlx.init(16\*16\*16,16\*16\*4);

REP(i,16) REP(j,16) FOR(k,1,16) {

**if** (g[i][j]=='-'||g[i][j]=='A'-1+k)

add(i,j,k);

}

**static** int x=0;

**if** (x) puts(""); **else** x=1;

dlx.dance(0);

}

}

1. **可重复**

*//dlx:求解精确覆盖*

*//link的意思是,r覆盖了c*

*//暴力枚举,n个点覆盖m个格子; 注意一定要init*

**struct** DLX {

**const** **static** int maxn=1e5+7;

**const** **static** int maxd=1e4+7;

int n,m,size;

int U[maxn],D[maxn],R[maxn],L[maxn],col[maxn],row[maxn];

int H[maxd],S[maxd];*//S:cnt*

int ans[maxn];

void init(int \_n,int \_m) {

n=\_n; m=\_m; int i;

FOR(i,0,m) {

S[i]=0;

U[i]=D[i]=i;

L[i]=i-1,R[i]=i+1;

} R[m]=0; L[0]=m;

size=m;

FOR(i,0,n) H[i]=-1;

}

void link(int r,int c) {

S[col[++size]=c]++; row[size]=r;

D[size]=D[c]; U[D[c]]=size;

D[c]=size; U[size]=c;

**if** (H[r]<0) H[r]=L[size]=R[size]=size;

**else** {

R[size]=R[H[r]];

L[R[H[r]]]=size;

L[size]=H[r];

R[H[r]]=size;

}

}

void remove(int c) {

**for** (int i=D[c]; i!=c; i=D[i])

L[R[i]]=L[i],R[L[i]]=R[i];

}

void resume(int c) {

**for** (int i=U[c]; i!=c; i=U[i])

L[R[i]]=R[L[i]]=i;

}

bool v[maxd];

*//估价函数,函数返回的是至少还需要多少行才能完成重复覆盖*

*//如果max的话可以直接cnt{R[]},也就是最多个数*

int f() {

int ret=0;

**for** (int c=R[0]; c; c=R[c]) v[c]=1;

**for** (int c=R[0]; c; c=R[c]) {

**if** (v[c]) {

ret++; v[c]=0;

**for** (int i=D[c]; i!=c; i=D[i])

**for** (int j=R[i]; j!=i; j=R[j])

v[col[j]]=0;

}

}

**return** ret;

}

int cnt;

void dance(int pos) {

**if** (R[0]==0) {cnt=min(cnt,pos); **return**;}

**if** (pos+f()>=cnt) **return**;

int c=R[0];

**for** (int i=R[0]; i; i=R[i])

**if** (S[i]<S[c]) c=i;

**for** (int i=D[c]; i!=c; i=D[i]) {

ans[pos]=row[i];

remove(i);

**for** (int j=R[i]; j!=i; j=R[j]) remove(j);

dance(pos+1);

**for** (int j=L[i]; j!=i; j=L[j]) resume(j);

resume(i);

}

}

} dlx;

int n,m;

int check(int x,int y,int a,int b,double d) {

**return** (x-a)\*(x-a)+(y-b)\*(y-b)<d\*d;

}

int x1[maxn],x2[maxn],y1[maxn],y2[maxn];

int main() {

int T;

scanf("%d",&T);

**while** (T--) {

int k,i;

scanf("%d%d%d",&n,&m,&k);

FOR(i,1,n) scanf("%d%d",&x1[i],&y1[i]);

FOR(i,1,m) scanf("%d%d",&x2[i],&y2[i]);

double l=0,r=1500;

**while** (r-l>1e-7) {

int i,j;

double mid=(l+r)/2;

*// printf("%f %f %f\n",l,mid,r);*

dlx.init(m,n);

FOR(i,1,n) FOR(j,1,m) {

**if** (check(x1[i],y1[i],x2[j],y2[j],mid))

dlx.link(j,i);

}

dlx.cnt=k+1;

dlx.dance(0);

**if** (dlx.cnt>k) l=mid;

**else** r=mid;

} printf("%.6f**\n**",l);

}

}

快速乘法(就那个long double的)

**return** ( x \* y - ( long long ) ( x / ( long double ) MOD \* y + 1e-8 ) \* MOD + MOD ) % MOD ;

快速乘法(mod加速的)

**typedef** unsigned long long u64;

**typedef** \_\_int128\_t i128;

**typedef** \_\_uint128\_t u128;

int \_,k;

u64 A0,A1,M0,M1,C,M;

**struct** Mod64 {

Mod64():n\_(0) {}

Mod64(u64 n):n\_(init(n)) {}

**static** u64 init(u64 w) { **return** reduce(u128(w) \* r2); }

**static** void set\_mod(u64 m) {

mod=m; assert(mod&1);

inv=m; rep(i,0,5) inv\*=2-inv\*m;

r2=-u128(m)%m;

}

**static** u64 reduce(u128 x) {

u64 y=u64(x>>64)-u64((u128(u64(x)\*inv)\*mod)>>64);

**return** ll(y)<0?y+mod:y;

}

Mod64& **operator** += (Mod64 rhs) { n\_+=rhs.n\_-mod; **if** (ll(n\_)<0) n\_+=mod; **return** \***this**; }

Mod64 **operator** + (Mod64 rhs) **const** { **return** Mod64(\***this**)+=rhs; }

Mod64& **operator** -= (Mod64 rhs) { n\_-=rhs.n\_; **if** (ll(n\_)<0) n\_+=mod; **return** \***this**; }

Mod64 **operator** - (Mod64 rhs) **const** { **return** Mod64(\***this**)-=rhs; }

Mod64& **operator** \*= (Mod64 rhs) { n\_=reduce(u128(n\_)\*rhs.n\_); **return** \***this**; }

Mod64 **operator** \* (Mod64 rhs) **const** { **return** Mod64(\***this**)\*=rhs; }

u64 get() **const** { **return** reduce(n\_); }

**static** u64 mod,inv,r2;

u64 n\_;

};

u64 Mod64::mod,Mod64::inv,Mod64::r2;

u64 pmod(u64 a,u64 b,u64 p) {

u64 d=(u64)floor(a\*(long double)b/p+0.5);

ll ret=a\*b-d\*p;

**if** (ret<0) ret+=p;

**return** ret;

}

void bruteforce() {

u64 ans=1;

**for** (int i=0;i<=k;i++) {

ans=pmod(ans,A0,M);

u64 A2=pmod(M0,A1,M)+pmod(M1,A0,M)+C;

**while** (A2>=M) A2-=M;

A0=A1; A1=A2;

}

printf("%llu**\n**",ans);

}

int main() {

**for** (scanf("%d",&\_);\_;\_--) {

scanf("%llu%llu%llu%llu%llu%llu%d",&A0,&A1,&M0,&M1,&C,&M,&k);

Mod64::set\_mod(M);

Mod64 a0(A0),a1(A1),m0(M0),m1(M1),c(C),ans(1),a2(0);

**for** (int i=0;i<=k;i++) {

ans=ans\*a0;

a2=m0\*a1+m1\*a0+c;

a0=a1; a1=a2;

}

printf("%llu**\n**",ans.get());

}

}

极小mex区间(r-l=max-min)

*//题意是个极小mex区间取得最值，求个什么东西*

int n;

**struct** frac {

int x,y;*//x/y*

frac(int \_x=1,int \_y=1):x(\_x),y(\_y) {}

bool **operator** < (**const** frac &A) **const** {

**return** (ll)x\*A.y<(ll)y\*A.x;

}

}

frac query(int x) {

} MIN[maxn]; *//len*

**inline** int lowbit(int x) {

**return** x&-x;

}

void update(int x,frac val) {

**for** (; x<=n; x+=lowbit(x))

MIN[x]=min(val,MIN[x]);

frac ans;

**for** (; x; x-=lowbit(x))

ans=min(ans,MIN[x]);

**return** ans;

}

vector<pair<int,frac> > segment[maxn];

map<int,int> MP;*//value,pos*

int R[maxn],cnt[maxn];

int A[maxn],pos[maxn];

vector<pair<int,int> > queries[maxn];

frac ans[maxn];

int main() {

int i,q;

scanf("%d%d",&n,&q);

FOR(i,1,n) {

scanf("%d",&A[i]);

assert(0<=A[i]&&A[i]<=n);

} FOR(i,0,n) pos[i]=n+1;

rFOR(i,1,n) {

R[i]=pos[A[i]]; pos[A[i]]=i;

} int now=0,tot=0;

FOR(i,1,n) {*//f[1]*

cnt[A[i]]++;

**if** (cnt[now]) {

**while** (cnt[now]) now++;

MP[now]=i;*//pos*

}

} MP[n+1]=n+1;*//为了简化操作*

FOR(i,1,n) { *//remove this 的贡献*

int position=i;*//should\_add*

**while** (1) {*//不能直接remove*

**auto** it=MP.lower\_bound(A[i]);

**if** (it==MP.end()) **break**;

**if** (it->second>=R[i]) **break**;*//del this; position*

int len=it->second-i+1,x=it->first;*//should+*

**if** (position==i) position=it->second;*//first\_change*

segment[i].push\_back(make\_pair(it->second,frac(len+1-x,len+1+x))); tot++;

**auto** itt=it; itt++;

int \_R=itt->second; MP.erase(it);

**if** (\_R>R[i]) MP[x]=R[i];*//insert\_more*

}*//not!*

**if** (position!=i&&A[i]) MP[A[i]]=position;

} assert(tot<=3\*n);*//最大3n以内; 应该说2n以内*

FOR(i,1,q) {*//为了用树状数组倒着查emmmm*

int l,r;

scanf("%d%d",&l,&r);

assert(1<=l&&l<=r&&r<=n);

queries[l].push\_back(make\_pair(r,i));

} rFOR(i,1,n) {

**for** (**auto** now:segment[i])

update(now.first,now.second);

**for** (**auto** now:queries[i])

ans[now.second]=query(now.first);

} FOR(i,1,q) {

int g=gcd(ans[i].x,ans[i].y);

printf("%d/%d**\n**",ans[i].x/g,ans[i].y/g);

}

}

set暴力修改区间

*//UVALive 8191 区间same*

set<pair<pii,int> > POS;

int cnt[maxn];

void update(int col,int x){

cnt[col]+=x;

}

void update(int l,int r,int x){

**auto** **final**=make\_pair(make\_pair(l,r),x);

**while** (l<=r){

**auto** it=POS.upper\_bound(make\_pair(make\_pair(l,INF),0)); it--;

**auto** now=\*it; POS.erase(it);

int nxtl=now.first.second+1;

**if** (now.first.first<l){

pair<int,int> remain;

remain.first=now.first.first;

remain.second=l-1;

**if** (remain.first<=remain.second)

POS.insert(make\_pair(remain,now.second));

}

**if** (now.first.second>r){

pair<int,int> remain;

remain.first=r+1;

remain.second=now.first.second;

**if** (remain.first<=remain.second)

POS.insert(make\_pair(remain,now.second));

nxtl=r+1;

}

update(now.second,-(nxtl-l));

update(x,nxtl-l);

l=nxtl;

} POS.insert(**final**);

}

int main() {

int l,c;

**while**(~scanf("%d%d%d",&l,&c,&n)){

int i; POS.clear();

FOR(i,1,c) cnt[i]=0; cnt[1]=l; int ans=0;

POS.insert(make\_pair(make\_pair(1,l),1));

FOR(i,1,n){

int p,x,a,b;

scanf("%d%d%d%d",&p,&x,&a,&b);

int S=cnt[p];

int m1=(a+(ll)S\*S)%l+1;

int m2=(a+(ll)(S+b)\*(S+b))%l+1;

**if** (m1>m2) swap(m1,m2);

update(m1,m2,x);

} FOR(i,1,c) ans=max(ans,cnt[i]);

printf("%d**\n**",ans);

}

}

螺旋数列value

*//123*

*//894*

*//765*

**inline** ll getValue(ll n,ll x,ll y) {

ll r = 0;

x=n-x+1; y=n-y+1;

**if** (x<=y && x+y <= n+1) {

r = x;

**return** 4\*(r-1)\*n - 4\*(r-1)\*(r-1) +1 + y-r;

}

**if** (x<=y && x+y >= n+1) {

r = n- y + 1;

**return** 4\*(r-1)\*n - 4\*(r-1)\*(r-1) + 1 + n-2\*r + 1 + x - r;

}

**if** (x>=y && x+y >= n+1) {

r = n - x +1;

**return** 4\*(r-1)\*n - 4\*(r-1)\*(r-1) + 1 + 3\*n-6\*r + 3 - y + r;

}

**if** (x>=y && x+y <= n+1) {

r = y;

**return** 4\*(r-1)\*n - 4\*(r-1)\*(r-1) + 1 + 4\*n-8\*r + 4 - x + r;

}

assert(0);

**return** -1;

}

# DP

决策单调性优化

**//决策单调性优化可以处理所有斜率优化的题**

**//题意:sum{A[l]->A[k],{1<=l<r<=n,k是l->r的路径上最近的标记点}}**

**//做法:DP; 注意有时DP[0]甚至DP[1]都要预处理的**

**//注意先写好DP方程**

**//注意DP方程上代表的意义!**

**//注意不能转移的地方!一定continue,否则可能破坏可以优化的性质**

**//我的理解:从左往右来看,如果l++,那么切的点只会向右移动,xl,xr是指转折点可能出现的位置;**

**//CDQ分治,传递下去了解可能存在的区间**

**//每次更新的是mid节点**

**//bfs,dfs均可,时间均为log(莫队不影响,莫队时间可证明nlogn)**

**//CF868F题意:切区间k段,每段数字出现个数sigma{n(n-1)/2}最小的个数**

LL L1[maxn],L2[maxn],R1[maxn],R2[maxn];*//前缀和之和,小技巧*

LL getL(int l,int r) { *//一个求l->r的点到l的sum和*

**return** (L2[r]-L2[l])-L1[l]\*(r-l);

}

LL getR(int l,int r) {

**return** (R2[l]-R2[r])-R1[r]\*(r-l);

}

LL pre[maxn],dp[maxn];

**struct** node {int l,r,xl,xr;};

LL cnt,sum,sum\_sum;

queue<node> Q;

void changel(LL val,int seg) {

sum\_sum+=sum\*seg\*2;

sum\_sum-=cnt\*val\*seg\*2;

cnt+=seg; sum+=val\*seg;

}

void changer(LL val,int seg) {

sum\_sum-=sum\*seg\*2;

sum\_sum+=cnt\*val\*seg\*2;

cnt+=seg; sum+=val\*seg;

}

int \_l,\_r;

LL A[maxn];

void changeto(int l,int r) {

**while** (\_r<r) \_r++,changer(A[\_r],1);

**while** (\_l>l) \_l--,changel(A[\_l],1);

**while** (\_l<l) changel(A[\_l],-1),\_l++;

**while** (\_r>r) changer(A[\_r],-1),\_r--;

}

void solve(int n) {

int i;

Q.push(node{1,n,0,n-1});

**while** (Q.size()) {

**auto** F=Q.front(); Q.pop();

int l=F.l,r=F.r,L=F.xl,R=F.xr;*//l,r,check\_l,check\_r*

int m=(l+r)/2,M=L;

LL &now=dp[m];

FOR(i,L,min(m-1,R)) {

*//这里changeto不会改变复杂度*

LL msum=(m-i)\*getL(m,n);

LL rsum=(n-m+1)\*(getR(i+1,m)+i\*(A[m]-A[i]));

**if** (now>pre[i]-msum-rsum)

now=pre[i]-msum-rsum,M=i;

}

**if** (l<m) Q.push(node{l,m-1,L,M});

**if** (r>m) Q.push(node{m+1,r,M,R});

}

}

*//DP[i]:i\_chosen; contains [i]->[i]; [i]->R(i+1->n)*

*//update:m [i-m]->[i], [i-m]->[m-n] [i-m]->[i-m]*

int T;

int n,m,k;

int i,j;

int main() {

**while** (~scanf("%d%d",&n,&k)) {

FOR(i,1,n) scanf("%lld",&A[i]);

A[0]=A[1]; A[n+1]=A[n];

FOR(i,1,n) L1[i]=A[i]-A[i-1]+L1[i-1];

FOR(i,1,n) L2[i]=L2[i-1]+L1[i];

rFOR(i,1,n) R1[i]=A[i+1]-A[i]+R1[i+1];

rFOR(i,1,n) R2[i]=R2[i+1]+R1[i];

\_l=1; \_r=0; sum=sum\_sum=cnt=0;

changeto(1,n);

FOR(i,0,n) dp[i]=sum\_sum;

FOR(i,1,k) {

int i;

FOR(i,0,n) pre[i]=dp[i];

solve(n);

*// FOR(m,1,n) FOR(i,0,m-1){*

*//// changeto(i+1,m);*

*//// cal:-=[m,n]->[i](differ)+[i+1-m](to m)*

*//// cal:-=[i+1,m]->[m,n](to m)*

*// LL msum=(m-i)\*getL(m,n);*

*// LL rsum=(n-m+1)\*(getR(i+1,m)+i\*(A[m]-A[i]));*

*// dp[m]=min(dp[m],pre[i]-msum-rsum);*

*// }*

}

LL ans=dp[0];

FOR(i,1,n) ans=min(ans,dp[i]);

printf("%lld**\n**",ans);

}

}

斜率优化

**//HDU 3480//斜率优化**

**//题意:一堆数字,切成k份,每块的代价为(max-min)^2**

**//dp方程:dp[i][j]=min{dp[k][j-1]+(a[i]-a[k+1])^2};**

**//dp方程:**

**//dp[i][j]=min{dp[k][j-1]+a[k+1]^2-2\*a[i]\*a[k+1]}+a[i]^2**

**//k=(dp[k][j-1](pre)+a[k+1]^2)/(a[k+1]),常数2\*a[i]**

**//斜率优化本质是维护一个下凸壳**

int n,m,i,j,k,t;

int a[maxn],pre[maxn],dp[maxn];

int head,tail;

int Q[maxn];*//id*

**inline** int getY(int id){

**return** pre[id]+a[id+1]\*a[id+1];

}

**inline** int getX(int id){

**return** a[id+1];

}

int main(){

int T,X=0;

scanf("%d",&T);

**while** (T--){

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%d",&a[i]);

sort(a+1,a+1+n);

int qi,qj,qk;

FOR(i,1,n) dp[i]=(a[i]-a[1])\*(a[i]-a[1]);

FOR(j,2,m){

FOR(i,1,n) pre[i]=dp[i];

head=tail=0;

dp[0]=0;Q[tail++]=0;

FOR(i,1,n){

**while** (head+1<tail){

qi=Q[head],qj=Q[head+1];

**if** (getY(qj)-getY(qi)<=2\*a[i]\*(getX(qj)-getX(qi))) head++;

**else** **break**;

}qi=Q[head];

dp[i]=pre[qi]+(a[i]-a[qi+1])\*(a[i]-a[qi+1]);

**while** (head+1<tail){

qi=Q[tail-2];qj=Q[tail-1];qk=i;

int y1=getY(qj)-getY(qi),x1=getX(qj)-getX(qi);

int y2=getY(qk)-getY(qj),x2=getX(qk)-getX(qj);

**if** (y2\*x1<=y1\*x2) tail--;*//y2/x2>y1/x1*

**else** **break**;

}Q[tail++]=i;

}

}

printf("Case %d: %d**\n**",++X,dp[n]);

}

}

四边形不等式优化

**//HDU 3516//四边形不等式优化**

**//题意:给定一个从左上往右下的图，只能往下往右连，求一个构造使得所有的边长度总和最小**

**//dp方程:**

**//dp[i][j]=max{dp[i][k]+dp[k+1][j]+x[k+1]-x[i]+y[k]-y[j]};**

**//能用：满足:**

**//w[i][j]+w[i'][j']<=w[i][j']+w[i'][j];**

**//w[i'][j']<=w[i][j],那么决策区间包含**

**struct** node{

int x,y;

}a[maxn];

int n,m,i,j,k,t;

int dp[maxn][maxn],pos[maxn][maxn];

int main(){

**while** (~scanf("%d",&n)){

FOR(i,1,n) scanf("%d%d",&a[i].x,&a[i].y),pos[i][i]=i;

FOR(i,1,n) FOR(j,i+1,n) dp[i][j]=INF;

FOR(t,1,n-1){

FOR(i,1,n-t){

j=i+t;

FOR(k,pos[i][j-1],min(j-1,pos[i+1][j])){

int now=dp[i][k]+dp[k+1][j]+a[k+1].x-a[i].x+a[k].y-a[j].y;

**if** (dp[i][j]>now){

dp[i][j]=now;

pos[i][j]=k;

}

}

}

}

printf("%d**\n**",dp[1][n]);

}

}

wqs二分(带权二分)

**//2018南京B, 使用的性质是每次更新的value斜率会下降**

**//所以二分这个斜率即可; 这个题套了个斜率优化**

**struct** node {

ll ans; int c;

} dp[maxn];

**struct** range {

int l,p;

} Q[maxn];

ll sum[maxn];

int A[maxn];

**inline** ll calc(int l,int r) {

ll ret=dp[l].ans;

l++; int mid=(l+r+1)/2;

ret+=(sum[r]-sum[mid]-(ll)(r+l–mid\*2)\*A[mid])-

(sum[mid-1]-sum[l-1]);

*// printf("%d-%d: ret=%lld = %lld - %lld\n",l,r,ret,*

*// (sum[r]-sum[mid]-(r-mid)\*A[mid]),*

*// (sum[mid-1]-sum[l-1]-(mid-l)\*A[mid]));*

**return** ret;

}

**inline** bool check(int a,int b,int p) {

**return** calc(a,p)>calc(b,p);

}

**inline** int solve\_slope(ll x) {

int st=0,ed=-1,i;

Q[++ed]=range{1,0}; Q[ed+1].l=n+1;

FOR(i,1,n) {

**if** (st<=ed&&Q[st+1].l-1<i) st++;

dp[i]= {calc(Q[st].p,i)+x,dp[Q[st].p].c+1};

*// printf("%d->%d : %lld\n",Q[st].p,i,dp[i].ans-dp[Q[st].p].ans);*

**if** (i==n) **return** dp[n].c;

**if** (!check(i,Q[ed].p,n)) {

**while** (st<=ed&&!check(i,Q[ed].p,Q[ed].l)) ed--;

**if** (st>ed) {

Q[++ed]= {i+1,i}; Q[ed+1].l=n+1;

} **else** {

*// printf("CHECK;\n");*

int l=Q[ed].l,r=n+1;

**while** (l+1<r) {

int mid=l+(r-l)/2;

**if** (check(i,Q[ed].p,mid)) l=mid;

**else** r=mid;

} Q[++ed]=range{r,i};

Q[ed+1].l=n+1;

}

}

*// int k; FOR(i,st,ed)*

*// printf(" x=%lld; %d - %d : %d; dp=%lld-%d\n",x,Q[ed].l,n,Q[ed].p,dp[i].ans,dp[i].c);*

}

**return** 0;

}

**inline** ll solve(int m) {

ll l=-100000000,r=100000001000000000;

**while** (l+1<r) {

ll mid=l+(r-l)/2;

**if** (solve\_slope(mid)<m) r=mid;

**else** l=mid;

} solve\_slope(l);

*// printf(" ans=%lld %d\n",dp[n].ans,dp[n].c);*

**return** dp[n].ans-l\*m;

}

int main() {

int i,k;

scanf("%d%d",&n,&k);

FOR(i,1,n) scanf("%d",&A[i]),sum[i]=A[i]+sum[i-1];

*// solve\_slope(100);*

printf("%lld**\n**",solve(k));

}

数位DP

**//当板子了**

**//这道题是连续的差最大是1**

**//需要注意时间空间限制,有时需要hash**

**//注意取模时底下calc也要取-\_-**

LL f[27][17][2];

int value[27];

LL calc(int x,int prev,int not\_0,int flag) {

**if** (x==0) **return** 1;

**if** (!flag&&f[x][prev][not\_0]!=-1)

**return** f[x][prev][not\_0];

LL ret=0; int i,maxi=9;

**if** (flag) maxi=min(maxi,value[x]);

FOR(i,0,maxi) {

*// if (not\_0||i)//这是与lead\_0有关的写法*

**if** (not\_0&&abs(prev-i)<2) **continue**;

**else** ret+=calc(x-1,i,not\_0||i,flag&&(i==maxi));

} **if** (!flag) f[x][prev][not\_0]=ret;

**return** ret;

} LL calc(LL x) {

int length=0;

**while** (x) value[++length]=x%10,x/=10;

**return** calc(length,0,0,1);

} LL calc(LL l,LL r) {

**return** calc(r)-calc(l-1);

}

int n,m,i,j,T;

int main() {

memset(f,0xff,**sizeof**(f));

FOR(i,1,10000)

**if** (calc(i,i)) printf("%d ",i);

puts("");

LL l,r;

scanf("%lld%lld",&l,&r);

printf("%lld**\n**",calc(l,r));

}

树形依赖背包

**// 树形依赖背包**

**// 题意: 是否存在块的val=i**

**// 做法: 先树分治变成必须包含top**

**// 然后往下dp, 按照dfs序看, 有一段是不能用的**

**// 所以倒着来dp或, 从下往上算贡献**

**// 大概做法是考虑这个点必选, 所以整体往右移val[x]来dp**

int A[maxn];

vector<int> edge[maxn];

int sz[maxn];

bool mark[maxn];

int minweight,root;

void dfs1(int x,int fa,int n) {

int weight=0;

sz[x]=1;

**for** (int v:edge[x]) {

**if** (v==fa||mark[v]) **continue**;

dfs1(v,x,n); sz[x]+=sz[v];

weight=max(weight,sz[v]);

} weight=max(weight,n-sz[x]);

**if** (weight<minweight) root=x,minweight=weight;

}

bitset<100007> now[3007],ans;*//depth*

void dfs2(int x,int fa,int dep) {

now[dep]=now[dep-1]; sz[x]=1;

**for** (int v:edge[x]) {

**if** (v==fa||mark[v]) **continue**;

dfs2(v,x,dep+1); sz[x]+=sz[v];

} now[dep-1]|=now[dep]<<A[x];

}

void dfs3(int x) {

debug("dfs3:%d**\n**",x);

now[0].reset(); now[0].set(0);

dfs2(x,0,1); mark[x]=1;

ans|=now[0];

**for** (int v:edge[x]) {

**if** (mark[v]) **continue**;

minweight=sz[v];

dfs1(v,0,sz[v]);

dfs3(root);

}

}

int main() {

int n,m,T;

int i;

scanf("%d",&T);

**while** (T--) {

scanf("%d%d",&n,&m);

REP(i,n-1) {

int u,v;

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

} FOR(i,1,n) scanf("%d",&A[i]);

minweight=n;

dfs1(1,0,n); dfs3(root);

FOR(i,1,m) printf("%d",(int)ans[i]);

puts("");

ans.reset();

FOR(i,1,n) edge[i].clear(),mark[i]=0;

}

**return** 0;

}

DP套DP

**//题意:麻将胡牌的可能种数**

**//为了不数漏,方法是这样的:**

**//首先考虑每个可能情况选择的个数,只可能有3\*3\*2=18种**

**//然后我们把状态压一下,每种牌型可能的1<<18的状态!**

**//对这个1<<18的状态进行转移**

void print2(int x) {

int i;

rREP(i,18) putchar(((x>>i)&1)+'0');

} int encode(int n\_2,int n\_1,int have2) { *//start from n-2 | n-1*

int ret=0;

ret=ret\*3+n\_2;

ret=ret\*3+n\_1;

ret=ret\*2+have2;

**return** ret;

} void decode(int e,int &n\_2,int &n\_1,int &have2) {

have2=e%2; e/=2;

n\_1=e%3; e/=3;

n\_2=e%3; e/=3;

}

void printstatus(int e) {

int n\_2,n\_1,have2;

decode(e,n\_2,n\_1,have2);

printf(" %d %d %d ",n\_2,n\_1,have2);

}

int getnextstatus(int status,int k) {

int nxtstatus=0,n;

int n\_2,n\_1,have2;

int x\_2,x\_1,xave2;

REP(n,18) **if** ((status>>n)&1) {

decode(n,n\_2,n\_1,have2);

x\_2=n\_1; x\_1=k-n\_2-n\_1; xave2=have2;

**if** (x\_1>=0) {

int x=encode(x\_2,x\_1%3,xave2);

nxtstatus|=(1<<x);

*// printstatus(n);printf("->");printstatus(x);printf("(+%d)",k);puts("");*

} **if** (!have2&&x\_1-2>=0) {

int x=encode(x\_2,x\_1-2,1);

nxtstatus|=(1<<x);

*// printstatus(n);printf("->");printstatus(x);printf("(+%d)",k);puts("");*

}

}

*// printf("get:%d->%d (k=%d)\n",status,nxtstatus,k);*

**return** nxtstatus;

}

queue<int> Q;

int id[1<<18|7],val[1007];

int tot;

int nxt[1007][7];

void initDP() {

int i,j; tot=0;

int k;*//this\_number*

Q.push(1); id[0]=++tot;

**while** (Q.size()) {

int status=Q.front(); Q.pop();

FOR(k,0,4) { *//只考虑这里产生2~*

int nxtstatus=getnextstatus(status,k);

**if** (!id[nxtstatus]) id[nxtstatus]=++tot,val[tot]=nxtstatus,Q.push(nxtstatus);

nxt[id[status]][k]=id[nxtstatus];

}

}

*// printf("%d\n",tot);*

*// REP(i,(1<<18)) if (id[i]){*

*// printf("(%-2d): ",id[i]);*

*// print2(i);puts("");*

*// REP(j,18) if ((i>>j)&1) printstatus(j);puts("");*

*// }*

*// FOR(i,1,tot){*

*// printf(" %-2d : ",i);*

*// print2(val[i]);puts("");*

*// REP(j,18) if ((val[i]>>j)&1) printstatus(j);puts("");*

*// }*

}

int dp[207][207][78];

**inline** void update(int &x,int y) {

((x+=y)>M)&&(x-=M);

}

int solve(int n,int m) {

int i,j,k,t;

FOR(i,0,n+3) FOR(j,0,m) FOR(t,0,68) dp[i][j][t]=0;

dp[0][0][1<<id[encode(0,0,0)]]=1;

FOR(i,0,n+3) {

int MAX;

**if** (i<n) MAX=4; **else** MAX=0;

FOR(j,0,m) {

FOR(t,1,tot) **if** (dp[i][j][t]) {

FOR(k,0,MAX) {

int nxtpos=nxt[t][k];

*// printf("%d->%d; k=%d\n",t,id[nxtstauts],k);*

update(dp[i+1][j+k][nxtpos],dp[i][j][t]);

}

}

}

} int ret=0;

*// FOR(t,1,tot) printf("%d: %d\n",t,dp[n+3][m][t]);*

FOR(t,1,tot) {

**if** ((val[t]>>encode(0,0,1))&1) {

update(ret,dp[n+3][m][t]);

*// printf("t=%d\n",t);*

}

}

**return** ret;

}

int main() {

int T;

initDP();

scanf("%d",&T);

**while** (T--) {

int n,m;

**static** int x=0;

scanf("%d%d",&n,&m);

printf("Case #%d: %d**\n**",++x,solve(n,m));

}

**return** 0;

}

插头DP

主要分两种dp方式

最小表示法:

**//插头dp长这样**

**// \_\_\_\_\_\_**

**//\_\_\_\_|o|->**

**//L==U就是环个数**

**template**<**typename** T1,**typename** T2> **struct** hashmap {

**const** **static** int seed=199991;

**const** **static** int maxn=1e6+7;

**struct** node {

T1 key; T2 val; int next;

node() {};

node(T1 k,T2 v,int n):key(k),val(v),next(n) {};

} T[maxn]; *//更好地空间局部性?(雾)*

int head[seed],size;

void clear() {

memset(head,-1,**sizeof**(head));

size=0;

}

void insert(T1 pos,T2 val) {

int x=pos%seed;

T[size]=node(pos,val,head[x]);

head[x]=size++;

}

bool find(int x) {

**for** (int i=head[x%seed]; ~i; i=T[i].next)

**if** (T[i].key==x) **return** 1;

**return** 0;

}

T2 &**operator** [](T1 x) {

**for** (int i=head[x%seed]; ~i; i=T[i].next)

**if** (T[i].key==x) **return** T[i].val;

insert(x,INF);

**return** T[size-1].val;

}

};

**typedef** hashmap<int,int> HASHMAP;

HASHMAP MP[2];

**inline** int getBit(int x,int k) {

**return** (x>>(k+k))&3;

}

**inline** int setBit(int x,int k,int v) {

**return** (x&~(3<<(k+k)))|(v<<(k+k));

}

**inline** void insert(HASHMAP &nxt,int k,int val) {

int &nxtval=nxt[k];

nxtval=min(nxtval,val);*//down*

}

**inline** void insert(HASHMAP &nxt,int k,int j,int down,int right,int val) {

k=setBit(k,j-1,down);

k=setBit(k,j,right);

int &nxtval=nxt[k];

nxtval=min(nxtval,val);*//down*

}

*//题意: 要从上往下划个线,把L和R分开; 不能有环*

char str[27][27];

int main() {

int n,m;

int i,j;

**while** (~scanf("%d%d",&n,&m)) {

FOR(i,1,n) scanf("%s",str[i]+1);

int now=0,nxt=1;

MP[now].clear();

FOR(j,1,m) {

int x=setBit(0,j,1);

MP[now].insert(x,0);

};*//top\_插头*

FOR(i,1,n) {

FOR(j,1,m) {*//check\_position; to\_right*

MP[nxt].clear();

int more=str[i][j]-'0';

**for** (int it=0; it<MP[now].size; it++) { *//x\_left; y:down*

int k=MP[now].T[it].key,val=MP[now].T[it].val;

int L=getBit(k,j-1),U=getBit(k,j);*//v=value*

int z=0;*//from left; downval\_count*

{int t; REP(t,j) **if** (getBit(k,t)) z^=1;}

**if** ((str[i][j]=='#'||str[i][j]=='W'||str[i][j]=='L')&&(L||U)) {

**continue**;*//有插头*

} **else** **if** (str[i][j]=='W') {

**if** (z) **continue**;

insert(MP[nxt],k,MP[now][k]);*//no way*

} **else** **if** (str[i][j]=='L') {

**if** (!z) **continue**;*//no!*

insert(MP[nxt],k,MP[now][k]);*//no way*

} **else** **if** (str[i][j]=='#') {*//all is ok*

insert(MP[nxt],k,MP[now][k]);*//no way*

} **else** {

**if** (L&&U) {*//value:(left and up)*

**if** (L==U) **continue**;

{

*//merge\_多个*

int t,\_k=0,c=0;

int v[10],id[4]={0,0,0,0};

REP(t,m+1) {

v[t]=getBit(k,t);

**if** (v[t]==L) v[t]=U;

} v[j-1]=v[j]=0;

REP(t,m+1) **if** (v[t]){

**if** (!id[v[t]]) id[v[t]]=++c;

\_k=setBit(\_k,t,id[v[t]]);

} insert(MP[nxt],\_k,val+more);

}

} **else** **if** (L||U) {*//left or up*

insert(MP[nxt],k,j,L|U,0,val+more);

insert(MP[nxt],k,j,0,L|U,val+more);

} **else** {*//circle*

insert(MP[nxt],k,val);*//not\_choose*

{

*//get\_value*

int t,\_k=0,c=0;

int v[10],id[4]={0,0,0,0};

REP(t,m+1) {

v[t]=getBit(k,t);

**if** (v[t]==L) v[t]=U;

} v[j-1]=v[j]=3;*//insert\_new*

REP(t,m+1) **if** (v[t]){

**if** (!id[v[t]]) id[v[t]]=++c;

\_k=setBit(\_k,t,id[v[t]]);

} insert(MP[nxt],\_k,val+more);*//not choose*

}

}

}

} now^=1; nxt^=1;

}

MP[nxt].clear();

**for** (int it=0; it<MP[now].size; it++) {

int k=MP[now].T[it].key; int w=MP[now].T[it].val;

**if** (!getBit(k,m)) insert(MP[nxt],k<<2,w);

} now^=1; nxt^=1;

}

int ans=INF;

FOR(j,1,m) {

int x=setBit(0,j,1);

ans=min(ans,MP[now][x]);

}

**if** (ans<INF) printf("%d**\n**",ans);

**else** puts("-1");

}

}

/\*

6 8

88W888L8

888#W888

888888L8

8W8LL#88

8888888L

00000W88

括号序列:

**//插头dp长这样**

**// \_\_\_\_\_\_**

**//\_\_\_\_|o|->**

**template**<**typename** T1,**typename** T2> **struct** hashmap {

**const** **static** int seed=199991;*//seed最好设置小点=\_=! 要clear*

**const** **static** int maxn=1e6+7;

**struct** node {

T1 key; T2 val; int next;

node() {};

node(T1 k,T2 v,int n):key(k),val(v),next(n) {};

} T[maxn]; *//更好地空间局部性?(雾)*

int head[seed],size;

void clear() {

memset(head,-1,**sizeof**(head));

size=0;

}

void insert(T1 pos,T2 val) {

int x=pos%seed;

T[size]=node(pos,val,head[x]);

head[x]=size++;

}

bool find(int x) {

**for** (int i=head[x%seed]; ~i; i=T[i].next)

**if** (T[i].key==x) **return** 1;

**return** 0;

}

T2 &**operator** [](T1 x) {

**for** (int i=head[x%seed]; ~i; i=T[i].next)

**if** (T[i].key==x) **return** T[i].val;

insert(x,INF);

**return** T[size-1].val;

}

};

**typedef** hashmap<int,int> HASHMAP;

HASHMAP MP[2];

**inline** int getBit(int x,int k) {

**return** (x>>(k+k))&3;

}

**inline** int setBit(int x,int k,int v) {*//注意这里是返回=\_=*

**return** (x&~(3<<(k+k)))|(v<<(k+k));

}

**inline** void insert(HASHMAP &nxt,int k,int val) {

int &nxtval=nxt[k];

nxtval=min(nxtval,val);*//down*

}

**inline** void insert(HASHMAP &nxt,int k,int j,int down,int right,int val) {

k=setBit(k,j-1,down);

k=setBit(k,j,right);

int &nxtval=nxt[k];

nxtval=min(nxtval,val);*//down*

}

*//题意: 要从上往下划个线,把L和R分开; 不能有环*

char str[27][27];

int n,m;

void printstatus(int k,int i,int j,**const** char str[]=""){

printf("%s: %d %d; status=",str,i,j);

REP(i,m+1) printf("%d",getBit(k,i));

*// system("pause");*

}

int main() {

int i,j;

**while** (~scanf("%d%d",&n,&m)) {

FOR(i,1,n) scanf("%s",str[i]+1);

int now=0,nxt=1;

MP[now].clear();

FOR(j,1,m) {

int x=setBit(0,j,3);

MP[now].insert(x,0);

};*//top\_插头*

FOR(i,1,n) {

FOR(j,1,m) {*//check\_position; to\_right*

MP[nxt].clear();

int more=str[i][j]-'0';

**for** (int it=0; it<MP[now].size; it++) { *//x\_left; y:down*

int k=MP[now].T[it].key,val=MP[now].T[it].val;

int L=getBit(k,j-1),U=getBit(k,j);*//v=value*

int z=0;*//from left; downval\_count*

{int t; REP(t,j) **if** (getBit(k,t)) z^=1;}

**if** ((str[i][j]=='#'||str[i][j]=='W'||str[i][j]=='L')&&(L||U)) {

**continue**;*//有插头*

} **else** **if** (str[i][j]=='W') {

**if** (z) **continue**;

insert(MP[nxt],k,MP[now][k]);*//no way*

} **else** **if** (str[i][j]=='L') {

**if** (!z) **continue**;*//no!*

insert(MP[nxt],k,MP[now][k]);*//no way*

} **else** **if** (str[i][j]=='#') {*//all is ok*

insert(MP[nxt],k,MP[now][k]);*//no way*

} **else** {

**if** (L&&U) {*//value:(left and up)*

**if** (L!=U) {

**if** (L==2&&U==1) {*// (value= \_|)*

insert(MP[nxt],k,j,0,0,val+more);

} **else** **if** (L==3&&U==1) {

int pos,\_k=setBit(k,j-1,0);

\_k=setBit(\_k,j,0);

FOR(pos,j+1,m) **if** (getBit(k,pos)==2) **break**;

**if** (0<=pos&&pos<=m) {

\_k=setBit(\_k,pos,3);

insert(MP[nxt],\_k,val+more);

} **else** printstatus(k,i,j,"bug1");

} **else** **if** (L==2&&U==3) {

int pos,\_k=setBit(k,j-1,0);

\_k=setBit(\_k,j,0);

rFOR(pos,0,j-2) **if** (getBit(k,pos)==1) **break**;

**if** (0<=pos&&pos<=m) {

\_k=setBit(\_k,pos,3);

insert(MP[nxt],\_k,val+more);

} **else** printstatus(k,i,j,"bug2");

} **else** **continue**;*//L=1&&U=2;*

*// merge\_circle((i,j)=bottom\_right(ex,ey))*

} **else** {

int pos,\_k=setBit(k,j,0); \_k=setBit(\_k,j-1,0);

**if** (L==1) {

FOR(pos,j+1,m) **if** (getBit(k,pos)==(L^3)) **break**;

} **else** {

rFOR(pos,0,j-2) **if** (getBit(k,pos)==(L^3)) **break**;

} **if** (0<=pos&&pos<=m) {

\_k=setBit(\_k,pos,L);

insert(MP[nxt],\_k,val+more);

} **else** printstatus(k,i,j,"bug3");

}

} **else** **if** (L||U) {*//left or up*

insert(MP[nxt],k,j,L|U,0,val+more);

insert(MP[nxt],k,j,0,L|U,val+more);

} **else** {

insert(MP[nxt],k,val);*//not\_choose*

insert(MP[nxt],k,j,1,2,val+more);*//new*

}

}

} now^=1; nxt^=1;

}

MP[nxt].clear(); *//to\_next(->)*

**for** (int it=0; it<MP[now].size; it++) {

int k=MP[now].T[it].key; int w=MP[now].T[it].val;

**if** (!getBit(k,m)) insert(MP[nxt],k<<2,w);

} now^=1; nxt^=1;

}

int ans=INF;

FOR(j,1,m) {

int x=setBit(0,j,3);

ans=min(ans,MP[now][x]);

}

**if** (ans<INF) printf("%d**\n**",ans);

**else** puts("-1");

}

}

斯坦纳树, 子集卷积的计数DP

**斯坦纳树:**

**//题意: 有几个点必须连接**

**//每个边的长度是1, 问你斯坦纳树有几个**

**// 斯坦纳树, 求min\_length很简单.. min\_cnt会重复计算, 所以从小到大计算**

**// len=1, 求方案数**

**struct** info {

int min,cnt;

info(int \_min=INF,int \_cnt=0):min(\_min),cnt(\_cnt) {};

} f[1<<12|7][57],g[1<<12|7][57];

**inline** void add(info &A,info B) {

**if** (A.min>B.min) A=info(B.min,0);

**if** (A.min==B.min) add\_(A.cnt,B.cnt);

}

**inline** info merge(info A,info B) {

info ret(A.min+B.min,(ll)A.cnt\*B.cnt%M);

**if** (ret.min>n) ret.min=n,ret.cnt=0;

**return** ret;

}

vector<int> edge[maxn];

vector<int> have[maxn];

int now[maxn],dep[maxn],vis[maxn];

int TaskA() {

int i,j,\_,maxs;

scanf("%d",&\_); maxs=1<<\_;

REP(i,n) edge[i].clear();

REP(i,m) {

int u,v;

scanf("%d%d",&u,&v);

u--; v--;

edge[u].push\_back(v);

edge[v].push\_back(u);

}

REP(i,maxs) REP(j,n) f[i][j]=g[i][j]=info(n,0);

REP(i,n) {

int cur=i<\_?1<<i:0; vis[i]=-1;

f[cur][i]=g[cur][i]=info(0,1);

}

int sta;

REP(sta,maxs) {

REP(i,n) {*//f:last\_op:addedge; g:no\_limit*

**if** (i<\_&&!((sta>>i)&1)) **continue**;

int remove=i<\_?1<<i:0; int remain=sta^remove;

int lowbit=remain&-remain; *// 防止重复计算, 一定注意这里是remain!*

**if** (remain)

**for** (int pre=remain&(remain-1); pre; pre=remain&(pre-1)) **if** (pre&lowbit)

add(g[sta][i],merge(f[pre|remove][i],g[(sta^pre)|remove][i]));

dep[i]=g[sta][i].min;

**if** (dep[i]<n) have[dep[i]].push\_back(i);

} *//?被卡常了?*

vector<int> Q;

REP(i,n) {

**for** (**auto** x:have[i]) {

**if** (vis[x]==sta) **continue**;

Q.push\_back(x); vis[x]=sta;

} **for** (**auto** x:Q) {

info now=info(g[sta][x].min+1,g[sta][x].cnt);

**for** (**auto** v:edge[x]) {

**if** (!(v<\_&&!((sta>>v)&1))) {

**if** (dep[v]>dep[x]+1) {

dep[v]=dep[x]+1;

have[dep[v]].push\_back(v);

}

} int nxtsta=v<\_?sta|(1<<v):sta;

add(g[nxtsta][v],now); add(f[nxtsta][v],now);

}

} Q.clear(); have[i].clear();

}

} *// printf("%d %d\n",g[maxs-1][1].min,g[maxs-1][1].cnt);*

printf("%d**\n**",g[maxs-1][1].cnt);

**return** 0;

}

**另一个题:**

**//题意:**

**//给一堆边, 每个生成树上的边贡献w[i]\*max(dep[u],dep[v])**

**//问你生成树总贡献**

**//做法: 枚举生成树, 然后直接dp 两边cnt和len得到答案**

//f:\sum{dep} g:\sum{cnt}

int e[17][17]; int ew[17][17];

int f[17][1<<12|7],g[17][1<<12|7];

int F[17][1<<12|7],G[17][1<<12|7];*//F,G:link*

int bit[1<<12|7];

int main() {

int i,j;

scanf("%d%d",&n,&m);

REP(i,m) {

int u,v,w;

scanf("%d%d%d",&u,&v,&w);

u--; v--; e[u][v]++; e[v][u]++;

ew[u][v]+=w; ew[v][u]+=w;

} int sta;

REP(i,n) g[i][1<<i]=1;

REP(sta,(1<<n)) bit[sta]=bit[sta>>1]+(sta&1);

REP(sta,(1<<n)) {

REP(i,n) **if** ((sta>>i)&1) { *//this\_root*

int remain=sta^(1<<i);

**if** (remain){

int low=remain&-remain;*//low写错了 =\_=*

**for** (int now=remain; now ; now=(now-1)&remain) **if** (now&low){

int sta1=now,sta2=sta^sta1;

add\_(f[i][sta],(ll)F[i][sta1]\*g[i][sta2]%M);

add\_(f[i][sta],(ll)G[i][sta1]\*f[i][sta2]%M);

add\_(g[i][sta],(ll)G[i][sta1]\*g[i][sta2]%M);

}

} **else** g[i][sta]=1;

REP(j,n) **if** (!((sta>>j)&1)&&e[i][j]){

add\_(F[j][sta],e[i][j]\*(f[i][sta]+(ll)g[i][sta]\*bit[sta]%M)%M);

add\_(G[j][sta],(ll)e[i][j]\*g[i][sta]%M);

}

}

} sta=(1<<n)-1; int ans=0;

REP(i,n) REP(j,n) **if** (ew[i][j]&&i!=j){

int s=sta^(1<<j);

**for** (int now=s; now; now=(now-1)&s) **if** ((now>>i)&1){

int sta1=now,sta2=sta^sta1;

int cnt=(f[i][sta1]+(ll)bit[sta1]\*g[i][sta1]%M)%M\*g[j][sta2]%M;

add\_(ans,(ll)ew[i][j]\*cnt%M);

}

} printf("%d**\n**",ans);

}

# 字符串

KMP|最小表示法

//记得border是个等差数列

int fail[maxn];

int check(char a[],int n){

fail[0]=fail[1]=0;

int i,j;

FOR(i,2,n){

j=fail[i-1];

**while** (j&&a[j+1]!=a[i]) j=fail[j];

**if** (a[j+1]==a[i]) fail[i]=j+1;

**else** fail[i]=0;

}**if** (n%(n-fail[n])==0) **return** n/(n-fail[n]);

**return** 1;

}

*//最小表示暴力法*

int getmin(char a[],int n){*//1-start*

int i,j,l;

FOR(i,1,n) a[i+n]=a[i];

i=1,j=2;

**while** (i<=n&&j<=n){

REP(l,n) **if** (a[i+l]!=a[j+l]) **break**;

**if** (l==n) **break**;

**if** (a[i+l]>a[j+l]) swap(i,j);

j=max(j+l+1,i+1);

}**return** i;

}

字典树

**//x xor v->max;**

**//没注释的是v<limit**

**//注释的是xor后小于limit**

**//计数问题有个套路:**

**//先算出全部,然后for一边容斥**

int nxt[maxn\*20\*10][2],tot;

int cnt[maxn\*20\*10];

LL xornum,limit;

void Ins(int &now,int k,int val) {

**if** (!now) now=++tot;

cnt[now]+=val;

**if** (k==-1) **return**;

int c=(xornum>>k)&1;

Ins(nxt[now][c],k-1,val);

}

LL Que(int now,int k,bool mark) { *//mark:have limit*

**if** (!now||!cnt[now]) **return** -INFF;

**if** (k==-1) **return** 0;

int c=(xornum>>k)&1,lim=(limit>>k)&1;

LL ret=-INFF;

**if** (!lim&&mark) {

**return** (c<<k)+Que(nxt[now][0],k-1,mark);

*// return Que(nxt[now][c],k-1,mark);*

} **else** {

ret=(1ll<<k)+Que(nxt[now][c^1],k-1,mark&&!(c&1));

**if** (ret<0) ret=Que(nxt[now][c],k-1,mark&&(c&1));

*// ret=(1ll<<k)+Que(nxt[now][c^1],k-1,mark);*

*// if (ret<0) ret=Que(nxt[now][c],k-1,0);*

} **return** ret;

}

AC自动机

**//HDU2896,匹配多串,查询id**

**namespace** ACM {

**const** int maxn=505\*140;

int next[maxn][98],fail[maxn],len[maxn],tot;

vector<int> have[maxn];

void init() {

tot=0; len[0]=0; fail[0]=0;

memset(next[0],0,**sizeof**(next[0]));

}

void insert(char s[],int id) {

int i,n=strlen(s),p=0;

REP(i,n) {

int c=s[i]-33;

**if** (!next[p][c]) {

next[p][c]=++tot; len[tot]=len[p]+1;

have[tot].clear(); fail[tot]=0;

memset(next[tot],0,**sizeof**(next[tot]));

} p=next[p][c];

} have[p].push\_back(id);

}

int Q[maxn],ST,ED;

void buildAC() {

ST=0; ED=-1; Q[++ED]=0;

**while** (ST<=ED) {

int p=Q[ST++],c;

REP(c,98) {

**if** (next[p][c]) {

fail[next[p][c]]=p?next[fail[p]][c]:0;

Q[++ED]=next[p][c];

} **else** next[p][c]=p?next[fail[p]][c]:0;*//否则可能fail=self*

}

**for** (int v:have[fail[p]])

have[p].push\_back(v);

}

}

void query(char a[],vector<int> &ans) {

int p=0;

int n=strlen(a),i;

REP(i,n) {

int c=a[i]-33; p=next[p][c];

**for** (int v:have[p]) ans.push\_back(v);

}

}

}

AC自动机 另一种写法

**// 2016南宁D**

**// 复杂度是所有串的len和**

**// 题意: 是否存在一个排列, 使得能一一对应**

**// 做法: 求每个点前相同val的len差, 然后直接AC自动机**

**// 修改fail的写法**

**namespace** ACM {

**const** int maxn=1e6+7;

map<int,int> next[maxn];

int fail[maxn],len[maxn],tot;

bool mark[maxn];

void init() {

tot=0; len[0]=0; fail[0]=0; mark[0]=0; next[0].clear();

}

void insert(int s[],int n) {

int i,p=0;

REP(i,n) {

int c=s[i];

**if** (!next[p].count(c)) {

next[p][c]=++tot; len[tot]=len[p]+1;

fail[tot]=0; mark[tot]=0;

next[tot].clear();

} p=next[p][c];

} mark[p]=1;

}

int Q[maxn],ST,ED;

**inline** int getnext(int x,int c){

**for** (;;x=fail[x]){

**if** (len[x]+1<=c) c=0;

**if** (!x||next[x].count(c)) **break**;

} **if** (next[x].count(c)) **return** next[x][c];

**return** x;

}

void buildAC() {

ST=0; ED=-1; Q[++ED]=0;

**while** (ST<=ED) {

int p=Q[ST++];

**for** (**auto** now:next[p]){

int c=now.first,nxt=now.second;

**if** (p) fail[nxt]=getnext(fail[p],c);

**else** fail[nxt]=0;

Q[++ED]=nxt;

} mark[p]|=mark[fail[p]];

}

}

bool query(int a[],int n) {

int p=0,have=0,i;

REP(i,n) {

int c=a[i]; p=getnext(p,c);

have|=mark[p];

} **return** have;

}

}

后缀数组

int wa[maxn],wb[maxn],wv[maxn],ws1[maxn];

int cmp(int \*r,int a,int b,int l) {

**return** r[a]==r[b]&&r[a+l]==r[b+l];

}

*//sa->pos(后缀排名->pos)*

void da(int \*r,int \*sa,int n,int m) {

r[n++]=0;*//使rank从1开始(sa[0]=n)*

int i,j,p,\*x=wa,\*y=wb,\*t;

REP(i,m) ws1[i]=0;*//pre-cmp*

REP(i,n) ws1[x[i]=r[i]]++;*//r->x*

rep(i,1,m) ws1[i]+=ws1[i-1];

rREP(i,n) sa[--ws1[x[i]]]=i;*//sort(计数排序)*

**for** (j=1,p=1; p<n; j<<=1,m=p) { *//j->2^x*

p=0; rep(i,n-j,n) y[p++]=i; *//最后j个是不用加(显然)*

REP(i,n) **if** (sa[i]>=j) y[p++]=sa[i]-j;*//后缀顺序*

REP(i,n) wv[i]=x[y[i]];*//x+y->wv(由于后缀顺序)*

REP(i,m) ws1[i]=0;

REP(i,n) ws1[wv[i]]++;

rep(i,1,m) ws1[i]+=ws1[i-1];

rREP(i,n) sa[--ws1[wv[i]]]=y[i];*//sort(计数排序)*

t=x,x=y,y=t;

p=1; x[sa[0]]=0;

rep(i,1,n) x[sa[i]]=cmp(y,sa[i-1],sa[i],j)?p-1:p++;

}

}

int rnk[maxn],height[maxn];

void calheight(int \*r,int \*sa,int n) {

int i,j,k=0;

FOR(i,1,n) rnk[sa[i]]=i;

REP(i,n) {

**if** (k) k--;

j=sa[rnk[i]-1];

**while** (r[i+k]==r[j+k]) k++;

height[rnk[i]]=k;

}

}

后缀自动机

**// 1题意:至少在k个子串中出现的子串数量**

**// 2题意:sigma{循环后匹配cnt}**

**// 这里的len不可以直接使用~ 原因是这里的len指的是原串len**

**// fail过后,len是可以直接使用的~ (会fail到确定的节点上)**

**// 这个fail的含义是说后缀相同,向前拓展的val(一个一个拓展len差项)**

**// sam反向不为拓扑序!注意自己进行拓扑排序**

**// 更新时注意len的限制!(因为更新时可能根本没有考虑前缀len)**

**// 注意nq在更新时更新时val和q是相等的,也就是说,维护值时nq要完全和q一样**

**// sum{len[x]-len[fail[x]]}=不同串个数,每个串代表fail->this的len**

**// 每个串的位置建议存的时候就保留下来~ 要不就有点麻烦了**

**// 复制出来的虚拟节点在计算次数时不参与计算~**

**// 也就是说计算相同串个数时,复制出来的只是个虚拟的节点**

**// query时在末尾加个0可以去掉很多的判断!**

**// 加空字符时注意len,这个len有两个作用:避免topo排错,减少add特判**

**// 加的不是root,就是个空字符,dfs的话只能dfs一个串!从后往前递推可行**

**// 如果是在一颗树上建,那么直接计数排序按len排是错的!一定注意!**

**// 注意看子串时的重复~**

**// 小技巧:由于每个节点对应的len是一定的,如果想要找l->r对应串可以倍增来找到对应的串**

**// 用fail建后缀树时,压缩路径第一个位置为pos[i]-len[fail[i]]**

**// 注意一件事:我这样做是并不能保证len[fail]!=len的**

**// 只有bfs trie可以保证,这样来进行按fail排序建立后缀树**

**// dfs trie的时间复杂度是trie叶结点深度和=\_=!证明..直接当多个**

**// 只有bfs能稳定的保证复杂度,但是好像没人这样卡人**

**struct** SAM{

int next[maxn][26],fail[maxn],len[maxn];

int cnt,last;

void init(){

cnt=last=0;fail[0]=-1;len[0]=0;

memset(next[0],0,**sizeof**(next[0]));

}

void add(int c){

int np=++cnt,p=last;

memset(next[np],0,**sizeof**(next[np]));

len[np]=len[p]+1;

**for** (;p!=-1&&!next[p][c];p=fail[p]) next[p][c]=np;

**if** (p==-1) fail[np]=0;

**else** {

int q=next[p][c];

**if** (len[p]+1==len[q]) fail[np]=q;

**else**{

int nq=++cnt;len[nq]=len[p]+1;

memcpy(next[nq],next[q],**sizeof**(next[q]));

fail[nq]=fail[q];

fail[np]=fail[q]=nq;

**for** (;p!=-1&&next[p][c]==q;p=fail[p]) next[p][c]=nq;

}

}

last=np;

}

*// 1:trie上建树,启发式合并set*

map<int,int> have[maxn];

int Next[maxn][26],Last[maxn],tot;

void add(char a[],int id){

int n=strlen(a),i,p=0;last=0;

REP(i,n) {

int c=a[i]-'a';

**if** (Next[p][c]) p=Next[p][c],last=Last[p];

**else** add(c),Last[p=Next[p][c]=++tot]=last;

have[last][id]++;

}

}

void merge(map<int,int> &A,map<int,int> &B){

**if** (A.size()<B.size()) swap(A,B);

**for** (**auto** now:B) A[now.first]+=now.second;

B.clear();*//delete &B;*

}

vector<int> edge[maxn];

LL Ans[maxn];

void DFS(int x,int k){

**for** (int v:edge[x]){DFS(v,k);merge(have[x],have[v]);}

**if** (have[x].size()>=k)

**for** (**auto** v:have[x])

Ans[v.first]+=(LL)v.second\*(len[x]-len[fail[x]]);

}

void solve(int k){

int i;

FOR(i,0,cnt) edge[i].clear();

FOR(i,1,cnt) edge[fail[i]].push\_back(i);

DFS(0,k);

}

*// 2:在query前进行了cnt[np]++和沿fail增加*

set<int> A;int CNT[maxn];

LL query(char a[]){

int i;LL ret=0;

int n=strlen(a),p=0,l=0;A.clear();

REP(i,n+n-1){

int c=a[i%n]-'a';

**if** (next[p][c]) l++,p=next[p][c];

**else** {

**while** (p!=-1&&!next[p][c]) p=fail[p];

**if** (p==-1) p=l=0;

**else** l=len[p]+1,p=next[p][c];

}**while** (len[fail[p]]>=n) p=fail[p],l=len[p];

**if** (l>=n){

**if** (A.count(p)) **continue**;

A.insert(p);

ret+=CNT[p];

}

*// if (l>=n) printf("i=%2d ret+id(%2d); l=%2d; +=%d\n",i,p,l,CNT[p]);*

}**return** ret;

}

void print(){

int i;

FOR(i,1,cnt) {

}

}

char a[maxn];

void dfs(int x=0,int len=0){

int i;

**for** (**auto** v:have[x])

printf("%2d(%2d) ",v.first,v.second);

puts("");

*// printf("%-3d(fail:%-3d,len=%-2d):%s\n",x,fail[x],this->len[x],a);*

REP(i,26){

**if** (next[x][i]){

a[len]=i+'a';

dfs(next[x][i],len+1);

a[len]=0;

}

}

}

}sam;

***后缀自动机+主席树合并***

**// 查询某串部分在串l->r的最大出现次数及位置**

**// SAM(这个套路)**

**// 做法:求出后缀树然后直接找到对应位置merge**

**// 这里可以看出, fail的含义就是说**

**// 某个位置往前len差长度的所有子串**

**// 然后对后缀树来建树然后对len倍增**

**// 就能求出对应的最短对应点来**

**// SPOJ COT4**

**// 题意:S串后面接字符生成新字符串**

**// T串由两个字符串(T串)接起来**

**// 问你Tj在Si出现次数**

**// (由于是Si所以dfsfail树,log查改)**

**// 做法是,S串建SAM,然后dfs出rank**

**// T串直接考虑在S串中的rank范围**

**struct** SAM{

**const** **static** int maxn=2e5+7;

int next[maxn][26],fail[maxn],len[maxn],cnt,pos[maxn],Pos[maxn];

void init(){

cnt=0;fail[0]=-1;len[0]=0;pos[0]=0;

memset(next[0],0,**sizeof**(next[0]));

}

int add(int p,int c,int id){

int np=++cnt;pos[np]=Pos[np]=id;

memset(next[np],0,**sizeof**(next[np]));

len[np]=len[p]+1;

**for** (;p!=-1&&!next[p][c];p=fail[p]) next[p][c]=np;

**if** (p==-1) fail[np]=0;

**else** {

int q=next[p][c];

**if** (len[p]+1==len[q]) fail[np]=q;

**else**{

int nq=++cnt;len[nq]=len[p]+1;

memcpy(next[nq],next[q],**sizeof**(next[q]));

fail[nq]=fail[q];pos[nq]=pos[q];Pos[nq]=0;

fail[np]=fail[q]=nq;

**for** (;p!=-1&&next[p][c]==q;p=fail[p]) next[p][c]=nq;

}

}

**return** np;

}

int failnext[maxn][26];

};

**struct** TRIE{

SAM sam;

**const** **static** int maxn=1e5+26+7;

void init(){

sam.init();

tot=0;ToT=1;id[ToT]=0;val[0]=-1;*//1:空*

memset(next[0],0,**sizeof**(next[0]));

}

*//1:trie上建树*

int id[maxn],ToT;*//queries*

int next[maxn][26],last[maxn],tot;*//on the trie*

int val[maxn];

void Add(int i,int c){

int p=id[i];ToT++;

**if** (!next[p][c]) {

next[p][c]=++tot;val[tot]=c;

memset(next[tot],0,**sizeof**(next[tot]));

fa[tot][0]=p;

}id[ToT]=next[p][c];

}

int Q[maxn],st,ed;

void buildSAM(){

st=ed=0;Q[ed++]=0;

**while** (st!=ed){

int p=Q[st++];char c;

REP(c,26) **if** (next[p][c]){

int nxt=next[p][c];

last[nxt]=sam.add(last[p],c,nxt);

Q[ed++]=nxt;

}

}

}

*//2:get L-R*

int failtot;

int rank[maxn],sa[maxn];

void dfsrank(int x){

**if** (sam.Pos[x]) rank[sam.Pos[x]]=++failtot,sa[failtot]=sam.Pos[x];

char c;

REP(c,26) **if** (sam.failnext[x][c]) dfsrank(sam.failnext[x][c]);

}

void linkfail(){

int i;

memset(sam.failnext,0,**sizeof**(sam.failnext[0])\*(sam.cnt+1));

FOR(i,1,sam.cnt)

sam.failnext[sam.fail[i]][val[prev(sam.pos[i],sam.len[sam.fail[i]])]]=i;

dfsrank(0);

}

*//3:build\_fa; ladder长链剖分*

int fa[maxn][21],son[maxn],top[maxn],len[maxn],dep[maxn];

vector<int> ladder[maxn],upper[maxn];

int upp[maxn];

void buildfa(){

int i,j,c;

dep[0]=0;

FOR(i,1,tot) rep(j,1,21)

fa[i][j]=fa[fa[i][j-1]][j-1],dep[i]=dep[fa[i][0]]+1;

rFOR(i,0,tot){

int o=0;top[i]=i;

ladder[i].clear();

REP(c,26) **if** (next[i][c]){

int p=next[i][c];

**if** (!o||len[o]<len[p]) o=p;

}**if** (o) len[i]=len[o]+1;**else** o=0;

son[i]=o;top[i]=i;

}FOR(i,0,tot) **if** (son[i]) top[son[i]]=top[i];

rFOR(i,0,tot) ladder[top[i]].push\_back(i);

FOR(i,1,tot) **if** (top[i]==i){

int u=i;

REP(j,len[i]){

u=fa[u][0];

ladder[i].push\_back(u);

**if** (!u) **break**;

}

}upp[0]=-1;

FOR(i,1,tot) upp[i]=upp[i-1]+(i==(i&-i));

}

int prev(int x,int k){;

**if** (!k) **return** x;

x=fa[x][upp[k]];k-=1<<upp[k];

k-=dep[x]-dep[top[x]];x=top[x];

**return** ladder[x][len[x]+k];

}

}trie;

**struct** queries{

int op,i,j;char c;

}q[maxn];

*//3:get Ans\_L-R*

int QAQ;

**struct** node{

int l,r,len,next;

node(){l=r=len=0;}

node(int \_l,int \_r,int \_len):l(\_l),r(\_r),len(\_len){};

}A[maxn],C[27];

node merge(node A,node B){*//反着来的,B在后*

**if** (A.len==0) **return** B;

**if** (B.len==0) **return** A;

**if** (B.l>B.r||A.l>A.r) **return** node(0,-1,A.len+B.len);

int l,r,L,R;

l=B.r+1;r=B.l-1;

**for**(L=B.l,R=B.r;L<=R;){

int mid=(L+R)/2;

**if** (trie.rank[trie.prev(trie.sa[mid],B.len)]<A.l) L=mid+1;

**else** R=mid-1,l=mid;

}

**for**(L=B.l,R=B.r;L<=R;){

int mid=(L+R)/2;

**if** (trie.rank[trie.prev(trie.sa[mid],B.len)]>A.r) R=mid-1;

**else** L=mid+1,r=mid;

}

**return** node(l,r,A.len+B.len);

}

*//4:solve*

int F[maxn];

**inline** int lowbit(int x){**return** x&-x;}

void update(int x,int val){

**for** (;x<=trie.failtot;x+=lowbit(x)) F[x]+=val;

}int query(int x){

int ret=0;

**for** (;x;x-=lowbit(x)) ret+=F[x];

**return** ret;

}int query(int l,int r){

**if** (l>r) **return** 0;

**return** query(r)-query(l-1);

}

node B[maxn];

int Ans[maxn],head[maxn];

void addnode(int x,int pos,int i){

x=trie.id[x];B[i]=A[pos];

B[i].next=head[x];head[x]=i;

}

void getans(int x){

int i;

**if** (x) update(trie.rank[x],1);

**for** (i=head[x];~i;i=B[i].next){

**if** (B[i].len&&B[i].l<=B[i].r) Ans[i]=query(B[i].l,B[i].r);

**else** Ans[i]=0;

}

REP(i,26) **if** (trie.next[x][i]) getans(trie.next[x][i]);

**if** (x) update(trie.rank[x],-1);

}

int n,m,Q;

int i,j,k;

char c;

int main(){

scanf("%d",&Q);

trie.init();

FOR(i,1,Q){

scanf("%d",&q[i].op);

**if** (q[i].op==1){

scanf("%d %c",&q[i].i,&c);q[i].c=c-'a';

trie.Add(q[i].i,q[i].c);

}**else** **if** (q[i].op==2){

scanf("%d%d %c",&q[i].i,&q[i].j,&c);q[i].c=c-'a';

}**else** scanf("%d%d",&q[i].i,&q[i].j);

}

trie.buildfa();

trie.buildSAM();

trie.linkfail();

REP(i,26){

int l,r,L,R;

l=trie.failtot+1;r=0;

**for**(L=1,R=trie.failtot;L<=R;){

int mid=(L+R)/2;

**if** (trie.val[trie.sa[mid]]<i) L=mid+1;

**else** R=mid-1,l=mid;

}

**for**(L=1,R=trie.failtot;L<=R;){

int mid=(L+R)/2;

**if** (trie.val[trie.sa[mid]]>i) R=mid-1;

**else** L=mid+1,r=mid;

}

C[i]=node(l,r,1);

}

FOR(i,0,trie.tot) head[i]=-1;QAQ=1;

FOR(i,1,Q){

**if** (q[i].op==1){

head[i]=-1;

}**else** **if** (q[i].op==2){

**if** (q[i].i==0) A[++QAQ]=merge(C[q[i].c],A[q[i].j]);

**if** (q[i].i==1) A[++QAQ]=merge(A[q[i].j],C[q[i].c]);

}**else** **if** (q[i].op==3)

A[++QAQ]=merge(A[q[i].i],A[q[i].j]);

**else** addnode(q[i].j,q[i].i,i);

}

getans(0);

FOR(i,1,Q) **if** (q[i].op==4) printf("%d**\n**",Ans[i]);

**return** 0;

}

马拉车

**//p是每个点为中心的延伸最长回文子串长度，-1就是原串以这个点为中心的长度**

**//看到题先去想这种方法，再说其他方法**

**//区间的最长子串做法是分成两段，然后直接考虑线段树分开算**

int n,m;

char s[maxn],str[maxn];

int len1,len2,p[maxn],ans;

void init() {

ans=0; int i;

str[0]='+'; str[1]='%';

REP(i,len1+1) {

str[i\*2+2]=s[i];

str[i\*2+3]='%';

} len2=len1\*2+2;

}

*// 主要是说已经对称匹配过的不用再进行*

void manacher() {

int id=0,mx=0; int i;

FOR(i,1,len2-1) {

**if** (mx>i) p[i]=min(p[2\*id-i],mx-i);

**else** p[i]=1;

**while** (str[i+p[i]]==str[i-p[i]]) p[i]++;

**if** (p[i]+i>mx) {

mx=p[i]+i; id=i;

}

}

}

回文自动机

**//next是将字符拼接到两端产生的字符串!**

**//一定注意这一点!**

**//也就是说,如果从上到下累积的话,可以很容易的将其与位置联系到一起!**

**//注意last是可以在线的,但是如果加了个其他的可以从fail上爬的,**

**//在讨论外边也要向上爬,或者一次过后就保存下来下次接着使用**

**//对于sans,diff,slink:**

**//sans是把之前的series\_ans保留下来**

**//diff相同时,sans一定会与上一个相同(由于对称的特殊性)**

**//所以只需改变diff改变时的ans即可**

**//区间本质不同回文串数**

**//由于border的特性, 可以通过等差数列的方法来分类更新答案**

**//bzoj5384,跳border+bit可以做到两个log**

**struct** PAM {

int next[maxn][27];

int fail[maxn];

int len[maxn];*//长度*

int diff[maxn];*//length(this-fail)*

int anc[maxn];*//diff不同的fail,共log个*

int S[maxn];*//字符*

int last;*//上一个字符节点*

int n,tot;*//n表示字符位置*

int newnode(int l) {

memset(next[tot],0,**sizeof**(next[tot]));

len[tot]=l;*//不是1...*

**return** tot++;

}

void init() {

tot=0; last=n=0;

newnode(0); newnode(-1);

S[n]=-1; fail[0]=1;

}

int getfail(int x) {

**while** (S[n-len[x]-1]!=S[n]) x=fail[x];

**return** x;

}

int add(int c) {

S[++n]=c;

int cur=getfail(last);

**if** (!next[cur][c]) {

int now=newnode(len[cur]+2);

fail[now]=next[getfail(fail[cur])][c];

diff[now]=len[now]-len[fail[now]];

**if** (diff[now]==diff[fail[now]])

anc[now]=anc[fail[now]];

**else** anc[now]=now;

next[cur][c]=now;*//这里一定要在fail后边=\_=*

} **return** last=next[cur][c];

}

char a[maxn];

void dfs(int p,int len=0) {

int c;

printf("%-20s (p=%-5d, length=%-5d fail=%-5d anc=%-5d diff=%-5d)",a,p,**this**->len[p],fail[p],anc[p],diff[p]);

*// if (p>=2) printf("%d len=%lld\n",);*

puts("");

REP(c,26) **if** (next[p][c]) {

a[len]=c+'a';

dfs(next[p][c],len+1);

a[len]=0;

}

}

} pam;

int dfn[maxn],out[maxn],tot;

vector<int> edge[maxn];

void getdfn(int x) {

dfn[x]=++tot;

**for** (int \_=0; \_<(int)edge[x].size(); \_++)

getdfn(edge[x][\_]);

out[x]=tot;

}

**namespace** SEG {

int MAX[maxn<<2];

void init(int val) {

memset(MAX,0,(val+1)\***sizeof**(int)\*4);

}

int query(int x,int l,int r,int L,int R) {

**if** (l<=L&&R<=r) **return** MAX[x];

int mid=(L+R)/2,ret=0;

**if** (l<=mid) ret=max(ret,query(x<<1,l,r,L,mid));

**if** (mid<r) ret=max(ret,query(x<<1|1,l,r,mid+1,R));

**return** ret;

}

void update(int x,int pos,int val,int L,int R) {

**if** (L==R) {MAX[x]=val; **return**;}

int mid=(L+R)/2;

**if** (pos<=mid) update(x<<1,pos,val,L,mid);

**else** update(x<<1|1,pos,val,mid+1,R);

MAX[x]=max(MAX[x<<1],MAX[x<<1|1]);

}

}

**namespace** BIT {

int sum[maxn],n;

void init(int val) {

memset(sum,0,(val+1)\***sizeof**(int)); n=val;

}

**inline** int lowbit(int x) {**return** x&-x;}

void add(int x,int val) {

**for** (; x<=n; x+=lowbit(x)) sum[x]+=val;

}

**inline** int get(int x) {

int ret=0;

**for** (; x; x-=lowbit(x)) ret+=sum[x];

**return** ret;

}

}

vector<pair<int,int> > queries[maxn];

char str[maxn];

int id[maxn];

ll ans[maxn];

int main() {

int n,q;

scanf("%d%d%s",&n,&q,str);

int i;

pam.init();

REP(i,n) id[i+1]=pam.add(str[i]-'a');

*// pam.dfs(0); puts("0");//2*

*// pam.dfs(1); puts("1");//1*

REP(i,pam.tot) edge[i].clear();

REP(i,pam.tot) **if** (i!=1) edge[pam.fail[i]].push\_back(i);

tot=0; getdfn(1);

FOR(i,1,q) {

int l,r; scanf("%d%d",&l,&r);

queries[r].push\_back(make\_pair(l,i));

} BIT::init(n);

SEG::init(tot);

FOR(i,1,n) {

*// for (int v=T.last;T.len[v]>0;v=T.slink[v]){*

*// sans[v]=f[i-(T.len[T.slink[v]]+T.diff[v])];*

*// if (T.diff[v]==T.diff[T.fail[v]])*

*// (sans[v]+=sans[T.fail[v]])%=M;*

*// if (!(i&1)) (f[i]+=sans[v])%=M;//f[x]*

*// }*

**for** (int p=id[i]; pam.len[p]>0; p=pam.fail[pam.anc[p]]) {

int l=max(1,SEG::query(1,dfn[p],out[p],1,tot)-pam.len[p]+1+1);

int r=i-pam.len[pam.anc[p]]+1+1;*//+1:start; 等差数列一起算*

BIT::add(l,1); BIT::add(r,-1);

} SEG::update(1,dfn[id[i]],i,1,tot);

**for** (int \_=0; \_<(int)queries[i].size(); \_++)

ans[queries[i][\_].second]=BIT::get(queries[i][\_].first);

}

int Ans=0;

FOR(i,1,q) Ans=(Ans+(ll)ans[i]\*i)%M;

printf("%d**\n**",Ans);

}

二分hash

**// wannafly挑战赛11D**

**// 题意:求上下拼接后的最长回文串长度(很坑)**

**struct** hashset{

**const** **static** int seed=1e7+7;

**const** **static** int maxn=2e6+7;

**struct** node{

int x,y;int next;

node(){};

node(int \_x,int \_y,int n):x(\_x),y(\_y),next(n){};

}T[maxn];*//更好地空间局部性?(雾)*

int head[seed],size;

void clear(){

memset(head,-1,**sizeof**(head));

size=0;

}

void insert(int x,int y){

int& h=head[x%seed];

**for** (int i=h;~i;i=T[i].next)

**if** (T[i].x==x&&T[i].y==y) **return**;

T[size]=node(x,y,h);h=size++;

}

bool count(int x,int y){

**for** (int i=head[x%seed];~i;i=T[i].next)

**if** (T[i].x==x&&T[i].y==y) **return** 1;

**return** 0;

}

}have;

**struct** hash{

int px[maxn],val[maxn],p;

void setp(int P,int n=200000){

int i;px[0]=1;p=P;

FOR(i,1,n) px[i]=(LL)px[i-1]\*p%M;

}

void set(char a[],int n){

int i;val[0]=0;

FOR(i,1,n) val[i]=((LL)val[i-1]\*p+a[i-1])%M;

}

int get(int l,int r){

l++;r++;

int ret=val[r]-(LL)val[l-1]\*px[r-l+1]%M;

(ret<0)&&(ret+=M);**return** ret;

}

}HA,RB;

void manacher(char A[],int p[],int len){

int id=0,mx=0,i;

rep(i,1,len){

**if** (mx>i) p[i]=min(p[2\*id-i],mx-i);

**else** p[i]=1;

**while** (A[i+p[i]]==A[i-p[i]]) p[i]++;

**if** (p[i]+i>mx) mx=p[i]+i,id=i;

}

}

int n,i;

int s[maxn];

char a[maxn],b[maxn],A[maxn\*2],B[maxn\*2];

int PA[maxn\*2],PB[maxn\*2];*//id*

int len,ans;

int main(){

scanf("%d",&n);

scanf("%s%s",a,b+1);

a[n]='(';b[0]=')';n++;

A[len]='+';B[len]='-';len++;

A[len]='%';B[len]='%';len++;

REP(i,n){

A[len]=a[i];B[len]=b[i];len++;

A[len]='%'; B[len]='%'; len++;

}A[len]='\*';B[len]='/';len++;

n=len;

manacher(A,PA,len);

manacher(B,PB,len);

HA.setp(19);RB.setp(19);

HA.set(A,n);reverse(B,B+n);RB.set(B,n);

reverse(B,B+n);

rep(i,1,n){

*//min(i-1-PA[i]+1,n-1-i-PA[i]+1)+1*

*//PA和PB的判断相同 (只需一个最大即可)*

PA[i]=max(PA[i],PB[i]);

int l=0,r=min(i-PA[i],n-1-i-PA[i])+1;*//r:not*

**while** (l+1<r){

int mid=(l+r)/2;

int hash\_A=HA.get(i-PA[i]-mid+1,i-PA[i]);

int hash\_B=RB.get(n-(i+PA[i]+mid),n-1-(i+PA[i]));

**if** (hash\_A==hash\_B) l=mid;

**else** r=mid;

}ans=max(ans,PA[i]+l);

}printf("%d**\n**",ans-1);

}

一些hashset|hashmap

**template**<**typename** T1,**typename** T2> **struct** hashmap{

**const** **static** int seed=999991;

**const** **static** int maxn=1e6+7;

**struct** node{

T1 key;T2 val;int next;

node(){};

node(T1 k,T2 v,int n):key(k),val(v),next(n){};

}T[maxn];*//更好地空间局部性?(雾)*

int head[seed],size;

void clear(){

memset(head,-1,**sizeof**(head));

size=0;

}

void insert(T1 pos,T2 val){

int x=pos%seed;

T[size]=node(pos,val,head[x]);

head[x]=size++;

}

T2 &**operator** [](T1 x){

**for** (int i=head[x%seed];~i;i=T[i].next)

**if** (T[i].key==x) **return** T[i].val;

insert(x,0);

**return** T[size-1].val;

}

};

**//用于字典树啥的空间优化**

**struct** linknode{

**struct** node{

int key,val;int next;

node(){};

node(int k,int v,int n):key(k),val(v),next(n){};

}T[maxn];*//更好地空间局部性?(雾)*

int head[maxn],size;

void clear(){

memset(head,-1,**sizeof**(head));

size=0;

}

int get(int x,int y){

**for** (int i=head[x];~i;i=T[i].next)

**if** (T[i].key==y) **return** T[i].val;

**return** 0;

}

void insert(int pos,int key,int val){

T[size]=node(key,val,head[pos]);

head[pos]=size++;

}

};

后缀平衡树

**// 替罪羊树...这道题卡splay,treap**

**// 题意：加字符，减字符，query子串个数**

**// 做法：建后缀自动机+LCT; right集个数**

**// 后缀自动机做法是直接链加链减**

**// 或者后缀顺序建平衡树然后树上query**

**// 后缀平衡树的顺序是倒着的, 倒着的后缀rank**

**// 以上是https://www.nowcoder.net/acm/contest/59/C**

**// 由于这个是倒着的rank, 反过来的情况非常常见(往前加)**

**// 这个直接用这个板子insert, query即可**

**const** double alpha=0.75;

**namespace** SAT {

**const** ull MAX=(1ull<<63)-1;

**struct** node {

int son[2]; int pre,size;

int sum,val; ull rank; char c;

void initval(char \_c) {

son[0]=son[1]=0; pre=0;

size=sum=val=1; rank=0; c=\_c;

}

} T[maxn];

int cnt,root,last;

**inline** bool cmp(int x,int y) {*//x<y*

assert(x!=y);

**if** (T[x].c!=T[y].c) **return** T[x].c<T[y].c;

**return** T[T[x].pre].rank<T[T[y].pre].rank;*//same:*

}

void pushup(int x){

T[x].size=1; T[x].sum=T[x].val;

**if** (T[x].son[0]) {

T[x].size+=T[T[x].son[0]].size;

T[x].sum+=T[T[x].son[0]].sum;

} **if** (T[x].son[1]) {

T[x].size+=T[T[x].son[1]].size;

T[x].sum+=T[T[x].son[1]].sum;

}

}

int id[maxn],tot;

bool rebuildRoot;*//手动rebuild\_{root}*

void getrank(int x) {

**if** (T[x].son[0]) getrank(T[x].son[0]);

**if** (!rebuildRoot||T[x].val) id[++tot]=x;

**if** (T[x].son[1]) getrank(T[x].son[1]);

}

void rerank(int &x,int l,int r,ull L,ull R) {

x=0; **if** (l>r) **return**;

ull mid=(L+R)/2; int m=(l+r)/2;

x=id[m]; T[x].rank=mid;

rerank(T[x].son[0],l,m-1,L,mid-1);

rerank(T[x].son[1],m+1,r,mid+1,R);

pushup(x);

}

void rebuild(int &x,ull l,ull r) {

**if** (!x) **return**;

tot=0; getrank(x);

rerank(x,1,tot,l,r);

}

void ins(int &x,ull l,ull r) {

ull mid=(l+r)/2;

**if** (!x) {x=cnt; **if** (l<=r) T[x].rank=mid; **return**;}

int p=cmp(x,cnt);

int &son=T[x].son[p];

**if** (p==0) ins(son,l,mid-1);

**else** ins(son,mid+1,r);

pushup(x); *//changes*

**if** (max(T[T[x].son[0]].size,T[T[x].son[1]].size)>

T[x].size\*alpha) rebuild(x,l,r);

}

void insert(char c) {

T[++cnt].initval(c);

T[cnt].pre=last; last=cnt;

ins(root,1,MAX);

**if** (!T[cnt].rank) {

rebuildRoot=true;

rebuild(root,1,MAX);

rebuildRoot=false;

}

}

void insert(char s[]) {

int len=strlen(s),i;

REP(i,len) insert(s[i]);

}

bool cmp(int k,char s[],int len) {*//smaller //okay!*

**for** (int i=0; i<len; i++,k=T[k].pre) {

**if** (!k) **return** 1;

**if** (s[i]!=T[k].c) **return** T[k].c<s[i];

} **return** 0;

}

int query(char s[],int len) {

int ret=0;

**for** (int now=root; now;) {

**if** (!cmp(now,s,len)) now=T[now].son[0];

**else** {

ret+=T[now].val+T[T[now].son[0]].sum,

now=T[now].son[1];

}

} **return** ret;

}

int query(char s[]) {

int len=strlen(s);

reverse(s,s+len); s[len]='Z'+1;*// s[len+1]=0;*

**return** query(s,len+1)-query(s,len);

}

void del(int k) {

**for** (; k&&last; last=T[last].pre,k--) {

int now;

**for** (now=root; now!=last;) {

T[now].sum--;

int p=T[last].rank>=T[now].rank;

now=T[now].son[p];

} assert(last==now);

T[last].val=0; T[last].sum--;

} **if** (!last) root=0;

}

void init(){

cnt=root=last=0;

}

}

*//2017icpc青岛J*

*//题意: 每个串找个后缀拼起来*

*//query 后缀最小序是多少*

*//倒着加, 然后找个最小rank把剩下的都去掉即可*

char pool[maxn],\*st=pool;

char \*A[maxn]; int len[maxn];

char ans[maxn];int L;

int main() {

int T;

scanf("%d",&T);

**while** (T--){

int i,j;

SAT::init(); L=0; st=pool;

scanf("%d",&n);

REP(i,n) {

A[i]=st,scanf("%s",A[i]);

st+=(len[i]=strlen(A[i]));

}

rREP(i,n) {

*// printf("i=%d;\n",i);*

rREP(j,len[i]) SAT::insert(A[i][j]);

int k=SAT::last; ull MIN=SAT::T[k].rank;int l=0;

REP(j,len[i]) {*//del\_cnt*

**if** (MIN>SAT::T[k].rank) MIN=SAT::T[k].rank,l=j;

k=SAT::T[k].pre;

} SAT::del(l);

rrep(j,l,len[i]) ans[L++]=A[i][j]; ans[L]=0;

} reverse(ans,ans+L);

printf("%s**\n**",ans);

}

**return** 0;

}

子序列自动机

**//序列自动机: 假设y之后第一次出现位置为nxt[x][y]**

**//f[x]=f[x]+f[nxt[x][y]]; 所以只与head是几有关**

**//trans[head][last(nextvalue)]**

**//保存的是第一个head处的值(转移)**

**//多加一个节点保存到结尾位置的ans是多少**

**//题意: 给个字符串, 每次在原串的每两个字符中间加个字符, 问你最后的子序列个数**

**//做法: 倒着发现是个复制后中间加个字符, 用子序列自动机来做**

int fa[maxn];

ll val[maxn];

**struct** mat{

int trans[27][27]; bool have[27];

mat(){memset(trans,0,**sizeof**(trans)); memset(have,0,**sizeof**(have));}

};

mat mul(mat A,mat B){

mat ret; int i,j,k;

REP(i,27) {

**if** (A.have[i]) {

REP(j,27) {

**if** (B.have[j])

REP(k,27) add\_(ret.trans[i][k],(ll)A.trans[i][j]\*B.trans[j][k]%M);

**else** add\_(ret.trans[i][j],A.trans[i][j]);

}

} **else** REP(j,27) add\_(ret.trans[i][j],B.trans[i][j]);

}

REP(i,27) ret.have[i]=A.have[i]|B.have[i];

**return** ret;

}

mat getMat(int c){

c-='a'; mat ret; int i; ret.have[c]=1;

REP(i,27) ret.trans[c][i]=1;

**return** ret;

}

int main() {

int i;

scanf("%d",&n);

scanf("%s",str);

mat ans=getMat(str[n-1]); int Ans=0;

rREP(i,n-1) {

*// int j,k;;*

*// mat mat1=mul(getMat(str[i]),ans);*

*// REP(j,27) {*

*// REP(k,27){*

*// printf("%2d ",mat1.trans[j][k]);*

*// } puts("<-mul");*

*// }*

ans=mul(ans,mul(getMat(str[i]),ans));

*// REP(j,27) {*

*// REP(k,27){*

*// printf("%2d ",ans.trans[j][k]);*

*// } puts("<-ans");*

*// }*

}

REP(i,26) add\_(Ans,ans.trans[i][26]);

printf("%d**\n**",Ans);

}

区间border

#include <vector>

#include <iostream>

#include <cstdio>

#include <cstdlib>

#include <cstring>

#include <algorithm>

#include <stack>

#define rep(i, a, n) for (int i = a; i < n; ++i)

#define per(i, a, n) for (int i = n - 1; i >= a; --i)

#define fi first

#define se second

#define SZ(x) ((int)(x).size())

**typedef** std::vector<int> VI;

**typedef** long long ll;

**typedef** std::pair<int, int> PII;

#define gcd(a, b) std::\_\_gcd((a), (b))

**const** int N = 1000010;

int ff[N], n, a[N], sql[N], sqr[N], sqr2[N];

**struct** SuffixArray {

int sa[N], rk[N], ht[N];

bool t[N << 1];

int hv[21][N];

bool islms(**const** int i, **const** bool \*t) {

**return** i > 0 && t[i] && !t[i - 1];

}

**template**<**class** **T**>

**inline** void sort(T s, int \*sa, **const** int len, **const** int sz, **const** int sigma, bool \*t,

int \*b, int \*cb, int \*p) {

memset(b, 0, **sizeof**(int) \* sigma);

memset(sa, -1, **sizeof**(int) \* len);

rep(i, 0, len) b[(int) s[i]]++;

cb[0] = b[0];

rep(i, 1, sigma) cb[i] = cb[i - 1] + b[i];

per(i, 0, sz) sa[--cb[(int) s[p[i]]]] = p[i];

rep(i, 1, sigma) cb[i] = cb[i - 1] + b[i - 1];

rep(i, 0, len) **if** (sa[i] > 0 && !t[sa[i] - 1]) sa[cb[(int) s[sa[i] - 1]]++] = sa[i] - 1;

cb[0] = b[0];

rep(i, 1, sigma) cb[i] = cb[i - 1] + b[i];

per(i, 0, len) **if** (sa[i] > 0 && t[sa[i] - 1]) sa[--cb[(int) s[sa[i] - 1]]] = sa[i] - 1;

}

**template**<**class** **T**>

**inline** void sais(T s, int \*sa, **const** int len, bool \*t, int \*b, int \*b1, **const** int sigma) {

int p = -1, \*cb = b + sigma;

t[len - 1] = 1;

per(i, 0, len - 1) t[i] = s[i] < s[i + 1] || (s[i] == s[i + 1] && t[i + 1]);

int sz = 0, cnt = 0;

rep(i, 1, len) **if** (t[i] && !t[i - 1]) b1[sz++] = i;

sort(s, sa, len, sz, sigma, t, b, cb, b1);

sz = 0;

rep(i, 0, len) **if** (islms(sa[i], t)) sa[sz++] = sa[i];

rep(i, sz, len) sa[i] = -1;

rep(i, 0, sz) {

int x = sa[i];

rep(j, 0, len) {

**if** (p == -1 || s[x + j] != s[p + j] || t[x + j] != t[p + j]) {

++cnt;

p = x;

**break**;

} **else** **if** (j > 0 && (islms(x + j, t) || islms(p + j, t))) {

**break**;

}

}

sa[sz + (x >>= 1)] = cnt - 1;

}

**for** (int i = len - 1, j = len - 1; i >= sz; --i) **if** (sa[i] >= 0) sa[j--] = sa[i];

int \*s1 = sa + len - sz, \*b2 = b1 + sz;

**if** (cnt < sz) sais(s1, sa, sz, t + len, b, b1 + sz, cnt);

**else**

rep(i, 0, sz) sa[s1[i]] = i;

rep(i, 0, sz) b2[i] = b1[sa[i]];

sort(s, sa, len, sz, sigma, t, b, cb, b2);

}

**template**<**class** **T**>

**inline** void getHeight(T s, int n) {

rep(i, 1, n + 1) rk[sa[i]] = i;

int j = 0, k = 0;

**for** (int i = 0; i < n; ht[rk[i++]] = k) {

**for** (k ? k-- : 0, j = sa[rk[i] - 1]; s[i + k] == s[j + k]; ++k);

}

}

**template**<**class** **T**>

**inline** void init(T s, **const** int len, **const** int sigma) {

sais(s, sa, len, t, rk, ht, sigma);

}

**inline** void solve(int \*s, int len) {

init(s, len + 1, 8);

getHeight(s, len);

rk[len] = 0;

rep(i, 1, len + 1) hv[0][i] = ht[i];

rep(j, 1, 20) **for** (int i = 1; i + (1 << j) - 1 <= len; ++i) {

hv[j][i] = std::min(hv[j - 1][i], hv[j - 1][i + (1 << (j - 1))]);

}

}

int lcp(int p, int q) {

**if** (q > n) **return** 0;

**if** (p == q) **return** n - p + 1;

p = rk[p - 1];

q = rk[q - 1];

**if** (p > q) std::swap(p, q);

int w = ff[q - p];

**return** std::min(hv[w][p + 1], hv[w][q - (1 << w) + 1]);

}

} s1, s2;

char s[N];

void gao(int l, int r, int ty) {

int pi = l - s2.lcp(n + 1 - r, n + 1 - (l - 1));

int pj = r + s1.lcp(l, r + 1);

int p = r - l + 1;

**if** (pj - pi + 1 >= 2 \* p && pi < l && l <= pi + p && (ty == 0 || pj != n)) {

**for** (int k = pi + 2 \* p - 1; k <= pj; ++k) {

sqr[k] = std::min(sqr[k], 2 \* p);

sql[k - 2 \* p + 1] = std::min(sql[k - 2 \* p + 1], 2 \* p);

sqr2[k] = std::max(sqr2[k], k - 2 \* p + 1);

}

}

}

void gao(int \*a) {

s1.solve(a + 1, n);

std::reverse(a + 1, n + a + 1);

a[n + 1] = 0;

s2.solve(a + 1, n);

*//std::reverse(a + 1, n + a + 1);*

std::stack<PII> st;

st.push({n + 1, n + 1});

per(i, 1, n + 1) {

int j = i;

**while** (SZ(st) > 1) {

**auto** seg = st.top();

**if** (s1.rk[i - 1] >= s1.rk[seg.fi - 1]) **break**;

j = seg.se;

st.pop();

}

st.push({i, j});

gao(i, j, 0);

}

st = std::stack<PII>();

st.push({n + 1, n + 1});

per(i, 1, n + 1) {

int j = i;

**while** (SZ(st) > 1) {

**auto** seg = st.top();

**if** (s1.rk[i - 1] <= s1.rk[seg.fi - 1]) **break**;

j = seg.se;

st.pop();

}

st.push({i, j});

gao(i, j, 1);

}

}

**namespace** border {

int tmp1[N], tmp2[N], c[N], sa[N], rk[21][N], lev;

PII pos[21][N];

void buildDict(char \*s, int \*sa, int \*x, int \*y, int n, int m = 128) {

rep(i, 0, m) c[i] = 0;

rep(i, 0, n) c[x[i] = s[i]]++;

rep(i, 1, m) c[i] += c[i - 1];

per(i, 0, n) sa[--c[x[i]]] = i;

rep(i, 0, n) rk[0][i] = x[i];

rep(i, 0, n) pos[0][i] = {rk[0][sa[i]], sa[i]};

lev = 1;

**for** (int k = 1; k < n; k <<= 1, ++lev) {

int p = 0;

per(i, n - k, n) y[p++] = i;

rep(i, 0, n) **if** (sa[i] >= k) y[p++] = sa[i] - k;

rep(i, 0, m) c[i] = 0;

rep(i, 0, n) c[x[y[i]]] ++;

rep(i, 1, m) c[i] += c[i - 1];

per(i, 0, n) sa[--c[x[y[i]]]] = y[i];

std::swap(x, y);

p = 1;

x[sa[0]] = 0;

y[n] = -1;

*//rep(i, 1, n) if (y[sa[i - 1]] == y[sa[i]]) assert(sa[i - 1] + k <= n && sa[i] + k <= n);*

rep(i, 1, n) x[sa[i]] = (y[sa[i - 1]] == y[sa[i]] && y[sa[i - 1] + k] == y[sa[i] + k]) ? p - 1 : p++;

rep(i, 0, n) rk[lev][i] = x[i];

rep(i, 0, n) pos[lev][i] = {rk[lev][sa[i]], sa[i]};

m = p;

}

}

**struct** seq {

int a, k, sz;

bool contain(int x) {

**if** (sz == 0) **return** 0;

**if** (x < a || x > a + (sz - 1) \* k) **return** 0;

**if** (x == a) **return** 1;

**return** (x - a) % k == 0;

}

};

ll Inv(ll q, ll m) {

**if** (q == 0) **return** 0;

*//assert(q >= 0);*

ll a1 = m, b1 = 0, a2 = q, b2 = 1, a3, b3, t;

**while** (a2 != 1) {

t = a1 / a2, a3 = a1 - t \* a2, b3 = b1 - t \* b2 % m,

a1 = a2, a2 = a3, b1 = b2, b2 = b3;

**if** (b2 < 0) b2 += m;

}

**return** b2;

}

std::pair<ll, ll> merge(ll a, ll b, ll c, ll d) {

c -= a;

ll dd = gcd(b, d);

**if** (c % dd != 0) **return** {-1, -1};

b /= dd;

c /= dd;

d /= dd;

ll t = c \* Inv(b, d) % d;

**if** (t < 0) t += d;

**return** {b \* t \* dd + a, b \* d \* dd};

}

seq intersect(seq a, seq b) {

**if** (a.sz > b.sz) std::swap(a, b);

**if** (a.sz == 0) **return** a;

**else** if (a.sz == 1) {

**if** (b.contain(a.a)) **return** a;

**else** **return** (seq) {0, 0, 0};

} **else** {

std::pair<ll, ll> d = merge(a.a % a.k, a.k, b.a % b.k, b.k);

**if** (d.se == -1) **return** (seq) {0, 0, 0};

int l = std::max(a.a, b.a), r = std::min(a.a + (a.sz - 1) \* a.k, b.a + (b.sz - 1) \* b.k);

int pl = (l - d.fi + d.se - 1) / d.se, pr = (r - d.fi + d.se) / d.se - 1;

**if** (pl > pr) **return** (seq) {0, 0, 0};

**else** return (seq) {(int) (d.fi + pl \* d.se), (pl == pr) ? 1 : (int) d.se, pr - pl + 1};

}

}

int findprev(int p, int lev, int r) { *// <=r start position*

PII \*ps = std::lower\_bound(pos[lev], pos[lev] + n, std::make\_pair(rk[lev][p], r + 1));

**if** (ps != pos[lev]) --ps; **else** **return** -1;

**if** (ps->fi != rk[lev][p]) **return** -1;

**else** **return** ps->se;

}

int findnxt(int p, int lev, int l) {*// >=l*

PII \*ps = std::lower\_bound(pos[lev], pos[lev] + n, std::make\_pair(rk[lev][p], l));

**if** (ps == pos[lev] + n || ps->fi != rk[lev][p]) **return** -1;

**else** **return** ps->se;

}

int bit[24];

*//#define bit(k) (1<<(k))*

#define bit(k) bit[k]

seq occur(int p, int lev, int l, int r) {

int fp = findnxt(p, lev, l);

**if** (fp == -1 || fp > r) **return** (seq) {0, 0, 0};

int fq = findnxt(p, lev, fp + 1);

**if** (fq == -1 || fq > r) **return** (seq) {fp, 1, 1};

int fr = findprev(p, lev, r);

**return** (seq) {fp, fq - fp, (fr - fp) / (fq - fp) + 1};

}

int query(int l, int r) {

--l;

--r;

**for** (int k = lev; k >= 1; k--) {

**if** ((r - l + 1) <= bit(k - 1)) **continue**;

seq a = occur(l, k - 1, std::max(r - bit(k) + 1, l), r - bit(k - 1) + 1);

seq b = occur(r - bit(k - 1) + 1, k - 1, l, std::min(l + bit(k - 1), r - bit(k - 1) + 1));

a.a = l + r - (a.a + (a.sz - 1) \* a.k);

b.a += bit(k - 1) - 1;

seq c = intersect(a, b);

**if** (c.sz != 0 && c.a + (c.sz - 1) \* c.k == r) --c.sz;

**if** (c.sz != 0) **return** c.a + (c.sz - 1) \* c.k - l + 1;

}

**return** 0;

}

void init() {

buildDict(s, sa, tmp1, tmp2, n);

**for** (int i = 0; i < 24; ++i) bit[i] = 1 << i;

}

}

std::pair<char, char> input[N];

char tmpa[5], tmpb[5];

char convert() {

**switch**(tmpb[0]) {

**case** 'd': **return** 'a';

**case** 'r': **return** 'b';

**case** 'm': **return** 'c';

**case** 'f': **return** 'd';

**case** 's':

**if** (tmpb[1] == 'o') **return** 'e';

**if** (tmpb[1] == 'i') **return** 'g';

**case** 'l': **return** 'f';

}

*//assert(0);*

}

void gets() {

std::deque<char> dq;

**for** (int i = 1; i <= n; ++i) {

scanf("%s%s", tmpa, tmpb);

input[i] = {tmpa[0], convert()};

*//input[i].first = rand() & 1 ? 'a' : 'p';*

*//input[i].second = (rand() % 7) + 'a';*

**if** (input[i].first == 'a') dq.push\_back(input[i].second);

**else** dq.push\_front(input[i].second);

}

**for** (int i = 0; i < n; ++i) {

s[i] = dq.front();

dq.pop\_front();

}

}

*// 题意:求循环节*

int cntp[N];

int main() {

*//auto bg = clock();*

rep(i, 2, 1000001) ff[i] = ff[i >> 1] + 1;

scanf("%d", &n);

*//n = 1000000;*

rep(i, 1, n + 1) sql[i] = n + 1, sqr[i] = n + 1;

gets();

rep(i, 1, n + 1) {

a[i] = s[i - 1] - 'a' + 1;

}

*//std::cout << clock() - bg << std::endl;*

gao(a);

rep(i, 1, n + 1) sqr2[i] = std::max(sqr2[i], sqr2[i - 1]);

border::init();

*//std::cout << clock() - bg << std::endl;*

**for** (int i = n; i; --i) {

cntp[i] = cntp[i + 1] + input[i].first == 'p';

}

**for** (int i = 1; i <= n; ++i) {

int l = cntp[i + 1] + 1;

int r = l + i - 1;

int bd = border::query(l, r);

int tot = r - l + 1;

int len = tot - bd;

printf("%d**\n**", (tot + len - 1) / len);

}

*//std::cout << clock() - bg << std::endl;*

**return** 0;

}

# 数据结构

时间分治

**// 题意: 动态求桥的个数**

**// 做法: 按时间分治, 然后缩边, tarjan找桥**

int m;

**struct** E {

int v,w;

E() {}

E(int \_v,int \_w) {v=\_v; w=\_w;}

};

**struct** node {

int u,v,w,l,r;

node(int \_u,int \_v,int \_w,int \_l,int \_r) {

u=\_u; v=\_v; w=\_w; l=\_l; r=\_r;

}

};

vector<E> edge[maxn];

**typedef** long long ll;

vector<node> remain;

int key[maxn],dfn[maxn];

E build(int x,int fa) {*//build vt*

vector<E> ch; dfn[x]=1;

**for** (**auto** now:edge[x]) {

**if** (now.v==fa) **continue**;

E w=build(now.v,x); w.w+=now.w;

**if** (w.v) ch.push\_back(w);

} **if** (ch.size()>=2) key[x]=1;

**if** (key[x]) {

**for** (**auto** v:ch)

remain.push\_back(node(x,v.v,v.w,1,m));*//exist!*

**return** E(x,0);

} **if** (ch.size()) **return** ch[0];

**return** E(0,0);

}

int low[maxn],vis[maxn],id[maxn];

int S[maxn],top,tot;

*//first: more;*

void tarjan(int x,int fa) { *//先缩这个, 再建虚树*

dfn[x]=low[x]=++tot;

S[++top]=x; vis[x]=1;

int cntfa=0;

**for** (**auto** now:edge[x]) {

**if** (now.v==fa&&!cntfa) {

cntfa=1; **continue**;

}

**if** (!dfn[now.v]) {

tarjan(now.v,x);

low[x]=min(low[x],low[now.v]);

} **else** **if** (vis[now.v])

low[x]=min(low[x],dfn[now.v]);

} **if** (dfn[x]==low[x]) {

**while** (1) {

int now=S[top--];

vis[now]=0; id[now]=x;

**if** (now==x) **break**;

}

}

}

void clear\_and\_set(int l,int r,**const** vector<node> &nodes) {

**for** (**auto** x:nodes) {

edge[x.u].clear();

edge[x.v].clear();

dfn[x.u]=dfn[x.v]=0;

} tot=0;*//clear*

**for** (**auto** x:nodes) {

**if** (x.l<=l&&r<=x.r) {

edge[x.u].push\_back(E(x.v,x.w));

edge[x.v].push\_back(E(x.u,x.w));

}

}*//all\_have*

}

int ans[maxn];

void solve(int l,int r,vector<node> nodes,int base) { *//m:边数*

clear\_and\_set(l,r,nodes);

**for** (**auto** x:nodes) {

**if** (!dfn[x.u]) tarjan(x.u,0);

**if** (!dfn[x.v]) tarjan(x.v,0);

}*//tarjan*

vector<node> tmp;

int all=0;

**for** (**auto** x:nodes) {

**if** (id[x.u]!=id[x.v]) {

node nxt=x;

nxt.u=id[x.u]; nxt.v=id[x.v];

tmp.push\_back(nxt);

**if** (x.l<=l&&r<=x.r) all+=x.w;

}

}*//init*

nodes.swap(tmp);

**if** (l==r) {

ans[l]=base+all;

**return**;

} int mid=(l+r)/2,div;

tmp.clear();

clear\_and\_set(l,r,nodes);

**for** (**auto** x:nodes) key[x.u]=key[x.v]=0;

**for** (**auto** x:nodes) {

**if** (x.l<=mid&&!(x.l<=l&&r<=x.r)) {

key[x.u]=key[x.v]=1;

tmp.push\_back(x);

}

} div=0; remain.clear();

**for** (**auto** x:nodes) {

**if** (x.l<=mid&&!(x.l<=l&&r<=x.r)) {

**if** (!dfn[x.u]) build(x.u,0);

**if** (!dfn[x.v]) build(x.v,0);

}

}*//tarjan*

**for** (**auto** x:remain) tmp.push\_back(x),div+=x.w;;

solve(l,mid,tmp,all-div+base);

tmp.clear();

clear\_and\_set(l,r,nodes);

**for** (**auto** x:nodes) key[x.u]=key[x.v]=0;

**for** (**auto** x:nodes) {

**if** (mid<x.r&&!(x.l<=l&&r<=x.r)) {

key[x.u]=key[x.v]=1;

tmp.push\_back(x);

}

} div=0; remain.clear();

**for** (**auto** x:nodes) {

**if** (mid<x.r&&!(x.l<=l&&r<=x.r)) {

**if** (!dfn[x.u]) build(x.u,0);

**if** (!dfn[x.v]) build(x.v,0);

}

}*//tarjan*

**for** (**auto** x:remain) tmp.push\_back(x),div+=x.w;;

solve(mid+1,r,tmp,all-div+base);

}

int main() {

int i,n;

scanf("%d%d",&n,&m);

map<pair<int,int>,int> MP;

vector<node> init;

FOR(i,1,m) {

char op[4]; int u,v;

scanf("%s%d%d",op,&u,&v);

**if** (u>v) swap(u,v);

**if** (op[0]=='A') {

MP[make\_pair(u,v)]=i;

} **else** {

init.push\_back(node(u,v,1,MP[make\_pair(u,v)],i-1));

MP.erase(make\_pair(u,v));

}

} **for** (**auto** now:MP)

init.push\_back(node(now.first.first,now.first.second,1,now.second,m));

solve(1,m,init,0);

FOR(i,1,m) printf("%d**\n**",ans[i]);

**return** 0;

}

二维树状数组

**//poj2155,修改区间01,query单点01,差分来做**

int n,m;

int c[maxn][maxn];

int lowbit(int x){**return** x&-x;}

void update(int \_x,int \_y){

**for** (int x=\_x;x<=n;x+=lowbit(x))

**for** (int y=\_y;y<=n;y+=lowbit(y)) c[x][y]^=1;

}

int sum(int \_x,int \_y){

int ret=0;

**for** (int x=\_x;x;x-=lowbit(x))

**for** (int y=\_y;y;y-=lowbit(y)) ret^=c[x][y];

**return** ret;

}

树状数组 不大于k的最大值

**const** int MAX=1000000;

**inline** int lowbit(int x) {**return** x&-x;}

**inline** void insert(int x) {

**for** (; x<=MAX; x+=lowbit(x)) a[x]++;

}

**inline** int find(int x) {

**while** (x&&!a[x]) x^=lowbit(x);

**if** (!x) **return** 0;

int t=lowbit(x)>>1,y=a[x];

**while** (t) {

**if** (y-a[x-t]) y-=a[x-t];

**else** {y=a[x-t]; x=x-t;}

t>>=1;

}

**return** x;

}

BIT\_差分

LL A[maxn],B[maxn];*//A\*i+B*

**inline** int lowbit(int x){**return** x&-x;}

void Add(int x,LL val,LL VAL){

**for** (;x<=n;x+=lowbit(x)) (A[x]+=val)%=M,(B[x]+=VAL)%=M;

}

void add(int l,int r,LL val){

Add(l,val,-((l-1)\*val%M)+M);

Add(r+1,M-val,r\*val%M);

}

LL query(int x){

LL ret=0;**for** (int i=x;x;x-=lowbit(x)) (ret+=A[x]\*i+B[x])%=M;

**return** ret;

}

LL query(int l,int r){

**return** (query(r)-query(l-1)+M)%M;

}

二维线段树

**//单点修改区间查询min,max**

**struct** node{

int left,right;

}treeX[maxn\*4],treeY[maxn\*4];

int a[maxn\*4][maxn\*4];

int mx[maxn\*4][maxn\*4],mn[maxn\*4][maxn\*4];

void buildY(int x,int y,int yl,int yr){

treeY[y].left=yl,treeY[y].right=yr;

**if** (yl==yr){

**if** (treeX[x].left==treeX[x].right)

mx[x][y]=mn[x][y]=a[treeX[x].left][yl];

**else**{

mx[x][y]=max(mx[x<<1][y],mx[x<<1|1][y]);

mn[x][y]=min(mn[x<<1][y],mn[x<<1|1][y]);

}

**return**;

}

int mid=(yl+yr)/2;

buildY(x,y<<1,yl,mid);

buildY(x,y<<1|1,mid+1,yr);

mx[x][y]=max(mx[x][y<<1],mx[x][y<<1|1]);

mn[x][y]=min(mn[x][y<<1],mn[x][y<<1|1]);

}

void buildX(int x,int n,int xl,int xr){

treeX[x].left=xl,treeX[x].right=xr;

**if** (xl==xr){

buildY(x,1,1,n);

**return**;

}

int mid=(xl+xr)/2;

buildX(x<<1,n,xl,mid);

buildX(x<<1|1,n,mid+1,xr);

buildY(x,1,1,n);

}

int querymaxY(int x,int y,int yl,int yr){

int L=treeY[y].left,R=treeY[y].right;

**if** (yl<=L&&R<=yr){

**return** mx[x][y];

}

int mid=(L+R)/2,ret=0;

**if** (mid>=yl) ret=max(ret,querymaxY(x,y<<1,yl,yr));

**if** (yr>mid) ret=max(ret,querymaxY(x,y<<1|1,yl,yr));

**return** ret;

}

int querymaxX(int x,int xl,int xr,int yl,int yr){

int L=treeX[x].left,R=treeX[x].right;

**if** (xl<=L&&R<=xr){

**return** querymaxY(x,1,yl,yr);

}

int mid=(L+R)/2,ret=0;

**if** (mid>=xl) ret=max(ret,querymaxX(x<<1,xl,xr,yl,yr));

**if** (xr>mid) ret=max(ret,querymaxX(x<<1|1,xl,xr,yl,yr));

**return** ret;

}

int queryminY(int x,int y,int yl,int yr){

int L=treeY[y].left,R=treeY[y].right;

**if** (yl<=L&&R<=yr){

**return** mn[x][y];

}

int mid=(L+R)/2,ret=INF;

**if** (mid>=yl) ret=min(ret,queryminY(x,y<<1,yl,yr));

**if** (yr>mid) ret=min(ret,queryminY(x,y<<1|1,yl,yr));

**return** ret;

}

int queryminX(int x,int xl,int xr,int yl,int yr){

int L=treeX[x].left,R=treeX[x].right;

**if** (xl<=L&&R<=xr){

**return** queryminY(x,1,yl,yr);

}

int mid=(L+R)/2,ret=INF;

**if** (mid>=xl) ret=min(ret,queryminX(x<<1,xl,xr,yl,yr));

**if** (xr>mid) ret=min(ret,queryminX(x<<1|1,xl,xr,yl,yr));

**return** ret;

}

void updateY(int x,int y,int posy,int val){

int L=treeY[y].left,R=treeY[y].right;

**if** (L==R){

**if** (treeX[x].left==treeX[x].right)

mx[x][y]=mn[x][y]=val;

**else**{

mx[x][y]=max(mx[x<<1][y],mx[x<<1|1][y]);

mn[x][y]=min(mn[x<<1][y],mn[x<<1|1][y]);

}

**return**;

}

int mid=(L+R)/2;

**if** (mid>=posy) updateY(x,y<<1,posy,val);

**else** updateY(x,y<<1|1,posy,val);

mx[x][y]=max(mx[x][y<<1],mx[x][y<<1|1]);

mn[x][y]=min(mn[x][y<<1],mn[x][y<<1|1]);

}

void updateX(int x,int posx,int posy,int val){

int L=treeX[x].left,R=treeX[x].right;

**if** (L==R){

updateY(x,1,posy,val);

**return**;

}

int mid=(L+R)/2;

**if** (mid>=posx) updateX(x<<1,posx,posy,val);

**else** updateX(x<<1|1,posx,posy,val);

updateY(x,1,posy,val);

}

int n,m,q;

int i,j;

int ans;

int main(){

int T,x=0;

scanf("%d",&T);

**while** (T--){

scanf("%d",&n);

FOR(i,1,n)

FOR(j,1,n) scanf("%d",&a[i][j]);

buildX(1,n,1,n);

scanf("%d",&q);

printf("Case #%d:**\n**",++x);

**while** (q--){

int x,y,r;

scanf("%d%d%d",&x,&y,&r);

r/=2;

int xl=max(1,x-r),xr=min(n,x+r);

int yl=max(1,y-r),yr=min(n,y+r);

int MX=querymaxX(1,xl,xr,yl,yr);

int MN=queryminX(1,xl,xr,yl,yr);

updateX(1,x,y,(MX+MN)/2);

printf("%d**\n**",(MX+MN)/2);

}

}

}

吉爷爷线段树

**// 区间取max min, 维护其他值**

**// 直接维护max,min,第二小即可, 暴力改就行了, 一个log**

**struct** node{

int MIN,MINCNT,IMIN;

ll SUM;

node(){MIN=MINCNT=SUM=0; IMIN=INF;}

}T[maxn\*4];

void min\_(int &A,int B){(A>B)&&(A=B);}

char op[2]; ll ans;

void build(int x,int l,int r){

T[x]=node(); T[x].MINCNT=r-l+1;

**if** (l==r) **return**;

int mid=(l+r)/2;

build(x<<1,l,mid);

build(x<<1|1,mid+1,r);

}

**inline** void update(int x,int l,int r,int val,int L,int R,int tag=1) {

**if** (l<=L&&R<=r){

**if** (T[x].IMIN>val){

**if** (T[x].MIN<val) {

T[x].SUM+=(ll)(val-T[x].MIN)\*T[x].MINCNT;

T[x].MIN=val;

} **if** (tag) ans+=T[x].SUM;

**return**;

}

} int mid=(L+R)/2;

update(x<<1,L,mid,T[x].MIN,L,mid,0);

update(x<<1|1,mid+1,R,T[x].MIN,mid+1,R,0);

**if** (l<=mid) update(x<<1,l,r,val,L,mid);

**if** (mid<r) update(x<<1|1,l,r,val,mid+1,R);

T[x].SUM=T[x<<1].SUM+T[x<<1|1].SUM;

T[x].MIN=T[x].IMIN=INF;

T[x].MIN=min(T[x<<1].MIN,T[x<<1|1].MIN);

T[x].MINCNT=0;

**if** (T[x<<1].MIN==T[x].MIN) {

min\_(T[x].IMIN,T[x<<1].IMIN);

T[x].MINCNT+=T[x<<1].MINCNT;

} **else** min\_(T[x].IMIN,T[x<<1].MIN);

**if** (T[x<<1|1].MIN==T[x].MIN) {

min\_(T[x].IMIN,T[x<<1|1].IMIN);

T[x].MINCNT+=T[x<<1|1].MINCNT;

} **else** min\_(T[x].IMIN,T[x<<1|1].MIN);

}

扫描线 矩形周长并

int size;

int len[maxn\*2];

int n,m,i,j,k;

**struct** Seg {

**struct** node {

int left,right;

int len,num;

bool cl,cr;*//iff*

int lazy;

void update(int x) {

lazy+=x;

}

} tree[maxn\*4];

void pushup(int x) {

**if** (tree[x].lazy) {

tree[x].len=len[tree[x].right+1]-len[tree[x].left];

tree[x].cl=tree[x].cr=1; tree[x].num=2;

} **else** **if** (tree[x].left==tree[x].right) {

tree[x].len=0;

tree[x].cl=tree[x].cr=0; tree[x].num=0;

} **else** {

tree[x].len=tree[x<<1].len+tree[x<<1|1].len;

tree[x].num=tree[x<<1].num+tree[x<<1|1].num;

**if** (tree[x<<1].cr&&tree[x<<1|1].cl) tree[x].num-=2;

tree[x].cl=tree[x<<1].cl;

tree[x].cr=tree[x<<1|1].cr;

}

};

void build(int x,int l,int r) {

tree[x].left=l; tree[x].right=r;

tree[x].len=tree[x].lazy=0;

**if** (l==r) {

} **else** {

int mid=(l+r)/2;

build(x<<1,l,mid);

build(x<<1|1,mid+1,r);

pushup(x);

}

}

void update(int x,int l,int r,LL val) {

int L=tree[x].left,R=tree[x].right;

**if** (l<=L&&R<=r) {

tree[x].update(val);

pushup(x);

} **else** {

int mid=(L+R)/2;

**if** (mid>=l) update(x<<1,l,r,val);

**if** (r>mid) update(x<<1|1,l,r,val);

pushup(x);

}

}

int query(int x,int l,int r) { *//num*

int L=tree[x].left,R=tree[x].right;

**if** (l<=L&&R<=r) {

**return** tree[x].len;

} **else** {

int mid=(L+R)/2;

int ans;

**if** (mid>=l) ans+=query(x<<1,l,r);

**if** (r>mid) ans+=query(x<<1|1,l,r);

pushup(x);

**return** ans;

}

}

} T;

**struct** point {

int x1,x2,h;

int n;

bool **operator** <(**const** point &a)**const** {

**if** (h!=a.h) **return** h<a.h;

**return** n>a.n;

}

} a[maxn];

map<int,int> Hash;

int x1,x2,y1,y2;

int ans;

int len1,len2,num;

int main() {

**while** (~scanf("%d",&n)) {

**if** (n==0) **break**;

FOR(i,1,n) {

scanf("%d%d%d%d",&x1,&y1,&x2,&y2);

len[i\*2-1]=x1; len[i\*2]=x2;

a[i\*2-1].x1=x1; a[i\*2-1].x2=x2;

a[i\*2-1].n=1; a[i\*2-1].h=y1;

a[i\*2].x1=x1; a[i\*2].x2=x2;

a[i\*2].n=-1; a[i\*2].h=y2;

}

sort(a+1,a+n\*2+1);

sort(len+1,len+n\*2+1);

Hash.clear();

FOR(i,1,2\*n) Hash[len[i]]=i;

T.build(1,1,n\*2);

ans=0;

FOR(i,1,2\*n) {

len1=T.tree[1].len; num=T.tree[1].num;

T.update(1,Hash[a[i].x1],Hash[a[i].x2]-1,a[i].n);

len2=T.tree[1].len;

ans+=abs(len2-len1);

ans+=num\*(a[i].h-a[i-1].h);

}

printf("%d**\n**",ans);

}

}

主席树

**//静态区间第k大**

vector<int> v;*//学到的hash方法*

int getid(int x){**return** lower\_bound(v.begin(),v.end(),x)-v.begin()+1;}

int root[maxn],a[maxn],cnt;

**struct** Tnode{

int left,right,sum;

}T[maxn\*40];

void update(int l,int r,int &x,int y,int pos){

T[++cnt]=T[y];T[cnt].sum++;x=cnt;

**if** (l==r) **return**;

int mid=(l+r)/2;

**if** (mid>=pos) update(l,mid,T[x].left,T[y].left,pos);

**else** update(mid+1,r,T[x].right,T[y].right,pos);

}

int query(int l,int r,int x,int y,int k){

**if** (l==r) **return** l;

int mid=(l+r)/2;

int sum=T[T[y].left].sum-T[T[x].left].sum;

**if** (sum>=k) **return** query(l,mid,T[x].left,T[y].left,k);

**else** **return** query(mid+1,r,T[x].right,T[y].right,k-sum);

}

可持久化数组(主席树维护)

**struct** Tnode {

int left,right,val;

} T[maxn\*80];

int cnt=0;

void build(int &x,int l,int r) {

**if** (!x) x=++cnt;

**if** (l==r) {T[x].val=l; **return**;}

int mid=(l+r)/2;

build(T[x].left,l,mid);

build(T[x].right,mid+1,r);

}

void update(int &x,int y,int pos,int val,int l,int r) {

T[++cnt]=T[y]; x=cnt;

**if** (l==r) {T[x].val=val; **return**;}

int mid=(l+r)/2;

**if** (mid>=pos) update(T[x].left,T[y].left,pos,val,l,mid);

**else** update(T[x].right,T[y].right,pos,val,mid+1,r);

}

int query(int x,int pos,int l,int r) {

**if** (l==r) **return** T[x].val;

int mid=(l+r)/2;

**if** (mid>=pos) **return** query(T[x].left,pos,l,mid);

**else** **return** query(T[x].right,pos,mid+1,r);

}

int root[maxn];

int n,m;

int i,j,k,t;

int a,b,ans;

**inline** int getfather(int x) {

int t=query(root[i],x,1,n);

**if** (t==x) **return** x;

int fa=getfather(t);

update(root[i],root[i],x,fa,1,n);

**return** fa;

}

int main() {

scanf("%d%d",&n,&m);

build(root[0],1,n);

FOR(i,1,m) {

scanf("%d",&k);

root[i]=root[i-1];

**if** (k==1) {

scanf("%d%d",&a,&b);

a^=ans; b^=ans;

int x=getfather(a),y=getfather(b);

**if** (x==y) **continue**;

update(root[i],root[i],x,y,1,n);

} **else** **if** (k==2) {

scanf("%d",&t);

t^=ans;

root[i]=root[t];

} **else** {

scanf("%d%d",&a,&b);

int x=getfather(a),y=getfather(b);

a^=ans; b^=ans;

**if** (x==y) puts("1"),ans=1;

**else** puts("0"),ans=0;

}

}

}

树套树

**// zoj2112动态第k大(这个是类似kuangbin大佬的做法按点建树，我按权值多个log...)**

**struct** node{

int l,r,cnt;

node(){l=r=cnt=0;}

}T[2500010];

int cnt;

int SIZE;

**inline** int lowbit(int x){**return** x&(-x);}

void Update(int &x,int y,int l,int r,int pos,int val){

T[++cnt]=T[y];T[cnt].cnt+=val;x=cnt;

**if** (l==r) **return**;

int mid=(l+r)/2;

**if** (mid>=pos) Update(T[x].l,T[y].l,l,mid,pos,val);

**else** Update(T[x].r,T[y].r,mid+1,r,pos,val);

}

int n,m;

int root[maxn];

void update(int x,int pos,int val){

**while** (x<=n){

Update(root[x],root[x],1,SIZE,pos,val);

x+=lowbit(x);

}

}

int ROOT[maxn];

int useL[maxn],useR[maxn];*//现在的l/r*

int Query(int l,int r,int L,int R,int pos,int pre\_L,int pre\_R){*//颜色,pos L->R*

**if** (l==r) **return** l;

int x;

int mid=(l+r)/2,nowcnt=0;

**for**(x=L-1;x;x-=lowbit(x)) nowcnt-=T[T[useL[x]].l].cnt;

**for**(x=R;x;x-=lowbit(x)) nowcnt+=T[T[useR[x]].l].cnt;

nowcnt+=T[T[pre\_R].l].cnt-T[T[pre\_L].l].cnt;

**if** (nowcnt>=pos){

**for**(x=L-1;x;x-=lowbit(x)) useL[x]=T[useL[x]].l;

**for**(x=R;x;x-=lowbit(x)) useR[x]=T[useR[x]].l;

**return** Query(l,mid,L,R,pos,T[pre\_L].l,T[pre\_R].l);

}**else**{

**for**(x=L-1;x;x-=lowbit(x)) useL[x]=T[useL[x]].r;

**for**(x=R;x;x-=lowbit(x)) useR[x]=T[useR[x]].r;

**return** Query(mid+1,r,L,R,pos-nowcnt,T[pre\_L].r,T[pre\_R].r);

}

}

int query(int L,int R,int pos){

int x;

**for**(x=L-1;x;x-=lowbit(x)) useL[x]=root[x];

**for**(x=R;x;x-=lowbit(x)) useR[x]=root[x];

**return** Query(1,SIZE,L,R,pos,ROOT[L-1],ROOT[R]);

}

char K[maxn],Q[20];

int A[maxn][4];

int a[maxn];

vector<int> H;

**inline** int getid(int x){**return** lower\_bound(H.begin(),H.end(),x)-H.begin()+1;}

void solve(){

scanf("%d%d",&n,&m);

int i;

FOR(i,1,n) scanf("%d",&a[i]),H.push\_back(a[i]);

REP(i,m){

scanf("%s",Q);

K[i]=Q[0];

**if** (K[i]=='Q') scanf("%d%d%d",&A[i][0],&A[i][1],&A[i][2]);

**if** (K[i]=='C') scanf("%d%d",&A[i][0],&A[i][1]),H.push\_back(A[i][1]);

}

sort(H.begin(),H.end());H.erase(unique(H.begin(),H.end()),H.end());

SIZE=H.size();

cnt=0;

FOR(i,1,n) Update(ROOT[i],ROOT[i-1],1,SIZE,getid(a[i]),1);

REP(i,m){

**if** (K[i]=='Q') printf("%d**\n**",H[query(A[i][0],A[i][1],A[i][2])-1]);*//l,r,pos*

**if** (K[i]=='C'){

update(A[i][0],getid(a[A[i][0]]),-1);

a[A[i][0]]=A[i][1];

update(A[i][0],getid(A[i][1]),1);

}

}

FOR(i,1,n) root[i]=0;

FOR(i,1,cnt) T[i]=node();

vector<int>().swap(H);

}

CDQ分治(套线段树)

**// CF848C CDQ分治（区间数字出现的r-l之和）**

**//将所有操作计算成为add和del,然后solve(l,r),再去除影响**

**const** LL MAX=10000007;

**struct** node{

int l,r; LL sum;

}T[MAX];

int cnt;

void Update(int &x,int pos,int val,int l,int r){

**if** (!x) x=++cnt;

T[x].sum+=val;

**if** (l==r) **return**;

int mid=(l+r)/2;

**if** (mid>=pos) Update(T[x].l,pos,val,l,mid);

**else** Update(T[x].r,pos,val,mid+1,r);

}

LL Query(int x,int l,int r,int L,int R){

**if** (!x||(l<=L&&R<=r)) **return** T[x].sum;

int mid=(L+R)/2;

LL ret=0;

**if** (mid>=l) ret+=Query(T[x].l,l,r,L,mid);

**if** (r>mid) ret+=Query(T[x].r,l,r,mid+1,R);

**return** ret;

}

int n,m;

int root[maxn];

**inline** int lowbit(int x){

**return** x&-x;

}

void update(int x,int pos,int val){

**for** (;x<=n;x+=lowbit(x)) Update(root[x],pos,val,1,n);

}

LL query(int x,int l,int r){

LL ret=0;

**for** (;x;x-=lowbit(x))

ret+=Query(root[x],l,r,1,n);*//其实还是应该是r-(l-1)的*

**return** ret;

}

int a[maxn];

set<int> S[maxn];

void ins(int pos,int val){*//固定R (L用前缀和)*

S[val].insert(pos);

set<int>::iterator it=S[val].lower\_bound(pos),itt=it;itt++;

int pre=0,suf=0;

**if** (it!=S[val].begin()) it--,pre=\*it;

**if** (itt!=S[val].end()) suf=\*itt;

**if** (pre) update(pos,pre,pos-pre);

**if** (suf) update(suf,pos,suf-pos);

**if** (pre&&suf) update(suf,pre,pre-suf);

}

void del(int pos,int val){

set<int>::iterator it=S[val].lower\_bound(pos),itt=it;itt++;

int pre=0,suf=0;

**if** (it!=S[val].begin()) it--,pre=\*it;

**if** (itt!=S[val].end()) suf=\*itt;

**if** (pre) update(pos,pre,-(pos-pre));

**if** (suf) update(suf,pos,-(suf-pos));

**if** (pre&&suf) update(suf,pre,-(pre-suf));

S[val].erase(pos);

}

int i;

int main(){

scanf("%d%d",&n,&m);

FOR(i,1,n){

scanf("%d",&a[i]);

ins(i,a[i]);

}

REP(i,m){

int k;

scanf("%d",&k);

**if** (k==1){

int p,x;

scanf("%d%d",&p,&x);

del(p,a[p]);

a[p]=x;

ins(p,a[p]);

}**else** **if** (k==2){

int l,r;

scanf("%d%d",&l,&r);

printf("%I64d**\n**",query(r,l,r));

}

}

}

SPLAY

int A[maxn];

**struct** splay\_tree {

**struct** node {

int val,min,max,add,size,son[2];*//add=lazy*

bool rev;

void init(int \_val) { *//开始时T[i].val==a[i-1](线性的);*

val=min=max=\_val; size=1;

**if** (\_val==INF) max=-INF;

add=son[0]=son[1]=0; rev=0;

}

} T[maxn\*2]; *//内存池*

int fa[maxn\*2],root,tot;

void pushup(int x) {

T[x].min=T[x].max=T[x].val; T[x].size=1;

**if** (T[x].val==INF) T[x].max=-INF;

**if** (T[x].son[0]) {

T[x].min=min(T[x].min,T[T[x].son[0]].min);

T[x].max=max(T[x].max,T[T[x].son[0]].max);

T[x].size+=T[T[x].son[0]].size;

}

**if** (T[x].son[1]) {

T[x].min=min(T[x].min,T[T[x].son[1]].min);

T[x].max=max(T[x].max,T[T[x].son[1]].max);

T[x].size+=T[T[x].son[1]].size;

}

}

void pushdown(int x) {

**if** (x==0) **return**;

**if** (T[x].add) {

**if** (T[x].son[0]) {

T[T[x].son[0]].val+=T[x].add;

T[T[x].son[0]].min+=T[x].add;

T[T[x].son[0]].max+=T[x].add;

T[T[x].son[0]].add+=T[x].add;

}

**if** (T[x].son[1]) {

T[T[x].son[1]].val+=T[x].add;

T[T[x].son[1]].min+=T[x].add;

T[T[x].son[1]].max+=T[x].add;

T[T[x].son[1]].add+=T[x].add;

}

T[x].add=0;

}

**if** (T[x].rev) {

**if** (T[x].son[0]) T[T[x].son[0]].rev^=1;

**if** (T[x].son[1]) T[T[x].son[1]].rev^=1;

swap(T[x].son[0],T[x].son[1]);

T[x].rev=0;

}

}

void rotate(int x,int kind) { *//zig(1->) zag(0<-)都行*

int y=fa[x],z=fa[y];

T[y].son[!kind]=T[x].son[kind],fa[T[x].son[kind]]=y;

T[x].son[kind]=y,fa[y]=x;

T[z].son[T[z].son[1]==y]=x,fa[x]=z;

pushup(y);

}

void splay(int x,int goal) { *//node x->goal's son*

**if** (x==goal) **return**;

**while** (fa[x]!=goal) {

int y=fa[x],z=fa[y];

pushdown(z),pushdown(y),pushdown(x);

int rx=T[y].son[0]==x,ry=T[z].son[0]==y;

**if** (z==goal) rotate(x,rx);

**else** {

**if** (rx==ry) rotate(y,ry);

**else** rotate(x,rx);

rotate(x,ry);

}

} pushup(x);

**if** (goal==0) root=x;

}

int select(int pos) { *//getnode*

int u=root;

pushdown(u);

**while** (T[T[u].son[0]].size!=pos) { *//这里由于头节点有个-INF 所以不-1*

**if** (pos<T[T[u].son[0]].size) u=T[u].son[0];

**else** {

pos-=T[T[u].son[0]].size+1;

u=T[u].son[1];

} pushdown(u);

} **return** u;

}

*//下面是自己写的一点常用?函数*

void update(int l,int r,int val) {

int u=select(l-1),v=select(r+1);

splay(u,0); splay(v,u);

T[T[v].son[0]].min+=val;

T[T[v].son[0]].max+=val;

T[T[v].son[0]].val+=val;

T[T[v].son[0]].add+=val;*//lazy*

}

void reverse(int l,int r) {

int u=select(l-1),v=select(r+1);

splay(u,0); splay(v,u);

T[T[v].son[0]].rev^=1;

}

void revolve(int l,int r,int x) { *//l~r->循环往后x位*

int u=select(r-x),v=select(r+1);

splay(u,0); splay(v,u);

int tmp=T[v].son[0]; T[v].son[0]=0;

pushup(v); pushup(u);

u=select(l-1),v=select(l);

splay(u,0); splay(v,u);

fa[tmp]=v; T[v].son[0]=tmp;

pushup(v); pushup(u);

}

void cut(int l,int r,int x) { *//l~r->去掉的x位置后 //HDU3487*

int u=select(l-1),v=select(r+1);

splay(u,0); splay(v,u);

int tmp=T[v].son[0];

T[v].son[0]=0;

pushup(v); pushup(u);

u=select(x); v=select(x+1);

splay(u,0); splay(v,u);

fa[tmp]=v; T[v].son[0]=tmp;

pushup(v); pushup(u);

}

int query\_min(int l,int r) {

int u=select(l-1),v=select(r+1);

splay(u,0); splay(v,u);

**return** T[T[v].son[0]].min;

}

void insert(int x,int val) {

int u=select(x),v=select(x+1);

splay(u,0); splay(v,u);

++tot; **if** (tot==maxn) tot=1;

T[tot].init(val); fa[tot]=v;

T[v].son[0]=tot;

pushup(v); pushup(u);

}

void delfree(int x) {*//buffer*

**if** (x==0) **return**;

bufs++; **if** (bufs==maxn) bufs=1;

nodebuff[bufs]=x;

delfree(T[x].son[0]);

delfree(T[x].son[1]);

}

void erase(int l,int r) {

int u=select(l-1),v=select(r+1);

splay(u,0); splay(v,u);

delfree(T[v].son[0]);

T[v].son[0]=0;

pushup(v); pushup(u);

}

void exchange(int l1,int r1,int l2,int r2) { *//r1-l1+1?=r2-l2+1 OK*

**if** (l1>l2) {swap(l1,l2); swap(r1,r2);}

int u=select(l1-1),v=select(r1+1);

splay(u,0); splay(v,u);

int tmp=T[v].son[0]; T[v].son[0]=0;

pushup(v); pushup(u);

l2-=T[tmp].size; r2-=T[tmp].size;

int \_u=select(l2-1),\_v=select(r2+1);

splay(\_u,0); splay(\_v,\_u);

fa[tmp]=\_v;

swap(T[\_v].son[0],tmp);

pushup(\_v); pushup(\_u);

u=select(l1-1),v=select(l1);

splay(u,0); splay(v,u);

fa[tmp]=v;

T[v].son[0]=tmp;

pushup(v); pushup(u);

}

int nodebuff[maxn],bufs;*//bufs:position*

int build(int l,int r) { *//add\_list*

**if** (l>r) **return** 0;

++tot; **if** (tot==maxn) tot=1;

int ret=nodebuff[tot];

int mid=(l+r)/2;

T[ret].init(A[mid]);

**if** (l==r) **return** ret;

int ls=build(l,mid-1);

int rs=build(mid+1,r);

**if** (ls) fa[ls]=ret,T[ret].son[0]=ls;

**if** (rs) fa[rs]=ret,T[ret].son[1]=rs;

pushup(ret);

**return** ret;

}

void init(int n) {

int i; tot=0;

REP(i,maxn) nodebuff[i]=i;

rFOR(i,1,n) A[i+1]=A[i];

A[1]=A[n+2]=-INF;

root=build(1,n+2);

fa[root]=0; T[0].init(-INF);

fa[0]=0; T[0].son[1]=root; T[0].size=0;

}

} T;

SPLAY启发式合并

**//HDU6133，一棵树的合并**

**struct** splaytree {

**struct** node {

LL val,sum;

int son[2],size;

void init(LL \_val) {

val=sum=\_val; size=1;

son[0]=son[1]=0;

}

} T[maxn]; *//编号是对应的*

int fa[maxn];

int root;

**inline** void pushup(int x) {

T[x].sum=T[x].val;

T[x].size=1;

**if** (T[x].son[0]) {

T[x].sum+=T[T[x].son[0]].sum;

T[x].size+=T[T[x].son[0]].size;

}

**if** (T[x].son[1]) {

T[x].sum+=T[T[x].son[1]].sum;

T[x].size+=T[T[x].son[1]].size;

}

}

void rotate(int x,int kind) {

int y=fa[x],z=fa[y];

T[y].son[!kind]=T[x].son[kind],fa[T[x].son[kind]]=y;

T[x].son[kind]=y,fa[y]=x;

T[z].son[T[z].son[1]==y]=x,fa[x]=z;

pushup(y);

}

void splay(int x,int goal) {

**if** (x==goal) **return**;

**while** (fa[x]!=goal) {

int y=fa[x],z=fa[y];

int rx=T[y].son[0]==x,ry=T[z].son[0]==y;

**if** (z==goal) rotate(x,rx);

**else** {

**if** (rx==ry) rotate(y,ry);

**else** rotate(x,rx);

rotate(x,ry);

}

}

pushup(x);

**if** (goal==0) root=x;

}

LL insert(int x) { *//x为原先位置*

int u=root,f=0;

**while** (u) {

f=u;

**if** (T[x].val<T[u].val) u=T[u].son[0];

**else** u=T[u].son[1];

}

**if** (T[x].val<T[f].val) T[f].son[0]=x;

**else** T[f].son[1]=x;

fa[x]=f;

splay(x,0);

**return** T[T[x].son[0]].sum+T[x].val\*(T[T[x].son[1]].size+1);

}

LL dfs(int x) {

int l=T[x].son[0],r=T[x].son[1];

LL ret=0;

T[x].init(T[x].val);

**if** (l) ret+=dfs(l);

ret+=insert(x);

**if** (r) ret+=dfs(r);

**return** ret;

}

LL merge(int x,int y,LL tmp,LL ret) {

**if** (x==y) **return** tmp;

splay(x,0); splay(y,0);

**if** (T[x].size>T[y].size) swap(x,y),swap(tmp,ret);

root=y;

ret+=dfs(x);

**return** ret;

}

int getkth(int x,int k) { *//未验证,抄的前面那个板子*

int u=root;

**while** (T[T[u].son[0]].size!=k) {

**if** (k<T[T[u].son[0]].size) u=T[u].son[0];

**else** {

k-=T[T[u].son[0]].size+1;

u=T[u].son[1];

}

}

**return** T[x].val;

}

} T;

左偏树

**struct** node{

int l,r,val,len;

node(int \_val=0){l=r=len=0; val=\_val;}

}T[maxn]; int tot;

int merge(int x,int y){*//不能直接swap x和儿子, 否则可能不满足堆性质*

**if** (!x||!y) **return** x|y;

**if** (T[x].val>T[y].val) swap(x,y);

T[x].r=merge(T[x].r,y);

**if** (T[T[x].l].len<T[T[x].r].len) swap(T[x].l,T[x].r);

T[x].len=T[T[x].r].len+1;

**return** x;

}

int pop(int x) {

T[x].val=-1;

**return** merge(T[x].l,T[x].r);

}

可持久化非旋treap

**// 内存回收: 没写, 写的话可以直接用个东西保存指向它的pointer个数**

**// 辣鸡蓝桥杯的题目, 喵的什么垃圾评测机, 全MLE是个什么东西**

**namespace** persist\_treap {

**typedef** pair<int,int> Pair;

**const** int maxn=1e7;*//maxn>=2*

**struct** node {

int l,r,len,size;

ll val,lazy,sum;

node(ll \_val=0) {

l=r=len=0; val=\_val; lazy=0;

sum=val; size=1;

}

} T[maxn];

int root;

int pool[maxn],st,ed;*//ends*

void init() {

int i; ed=maxn-1; root=0; T[0].size=0;*//0:no use*

REP(i,maxn-1) pool[i]=i+1;*//start from 1*

}

void delnode(int pos) {

**if** (ed==maxn-1) ed=0;

*// T[pos]=node();// no use*

pool[ed++]=pos;

}

int insnode(ll x) {*//value*

*// assert(st+1!=ed);// no !!!*

int pos=pool[st++];

**if** (st==maxn-1) st=0;

T[pos]=node(x);

**return** pos;

}

int persistnode(int ini) {

*// assert(st+1!=ed);// no !!!*

int pos=pool[st++];

T[pos]=T[ini];

**if** (st==maxn-1) st=0;

**return** pos;

}

void ADD(int x,ll val) { *//update*

T[x].lazy+=val; T[x].val+=val;

T[x].sum+=T[x].size\*val;

}

bool pushdown(int x) {

**if** (!T[x].lazy) **return** 0;

**if** (T[x].l) {

T[x].l=persistnode(T[x].l);

ADD(T[x].l,T[x].lazy);

}

**if** (T[x].r) {

T[x].r=persistnode(T[x].r);

ADD(T[x].r,T[x].lazy);

}

T[x].lazy=0;

**return** 1; *// changed; -1/2空间*

}

void pushup(int x) {

T[x].sum=T[x].val;

T[x].len=0; T[x].size=1;

**if** (T[x].l) {

T[x].sum+=T[T[x].l].sum;

T[x].len=max(T[x].len,T[T[x].l].len+1);

T[x].size+=T[T[x].l].size;

}

**if** (T[x].r) {

T[x].sum+=T[T[x].r].sum;

T[x].len=max(T[x].len,T[T[x].r].len+1);

T[x].size+=T[T[x].r].size;

}

}

int merge(int x,int y,bool downx=0,bool downy=0) {

**if** (!x||!y) **return** x|y;

**if** (T[x].len>T[y].len) {

**if** (!downx) x=persistnode(x);

bool okay=pushdown(x);

T[x].r=merge(T[x].r,y,okay,downy);

pushup(x); **return** x;

} **else** {

**if** (!downy) y=persistnode(y);

bool okay=pushdown(y);

T[y].l=merge(x,T[y].l,downx,okay);

pushup(y); **return** y;

}

}

pii split(int x,int k,bool down=0) {

**if** (!x) **return** make\_pair(0,0);

**if** (!down) x=persistnode(x);

bool persisted=pushdown(x); *//persist:newnode*

*// printf("split: %lld; sz=%d; k=%d\n",T[x].val,T[x].size,k);*

pii P;

**if** (!k||T[T[x].l].size>=k) {

*// printf("to\_left %d\n",T[x].l);*

P=split(T[x].l,k,persisted);

T[x].l=P.second; pushup(x); P.second=x;

} **else** {

P=split(T[x].r,k-T[T[x].l].size-1,persisted);

T[x].r=P.first; pushup(x); P.first=x;

} **return** P;

}

void print\_dfs(int x) {

**if** (!x) **return**;

print\_dfs(T[x].l);

printf("%lld ",T[x].val);

print\_dfs(T[x].r);

}

ll query(int l,int r) {*//用个东西记录一下??*

pii A=split(root,l-1);

*// print\_dfs(A.first); puts("A.first");*

pii B=split(A.second,r-l+1);

*// print\_dfs(B.first); puts("B.first");*

*// print\_dfs(B.second); puts("B.second");*

**return** T[B.first].sum;

}

void update(int l,int r,ll val) {

pii A=split(root,l-1);

pii B=split(A.second,r-l+1);

ADD(B.first,val);

root=merge(merge(A.first,B.first),B.second);

}

void insert(int k,int val) {*//after kth*

pii A=split(root,k);

int y=insnode(val);

*// print\_dfs(A.first); puts("okay");*

*// printf("root=%d\n",root);*

root=merge(A.first,merge(y,A.second));

}

void transto(int l,int r,int x,int y) {

pii A=split(root,l-1);

pii B=split(A.second,r-l+1);

pii A\_=split(root,x-1);

pii B\_=split(A\_.second,y-x+1);

root=merge(merge(A.first,B\_.first),B.second);

}

}

LCT

**//确认没写错，加边减边，改边权，查第二大值**

**//修改边权:把边当成点,mark一下,然后左右端点连边即可**

**//hdu5002, chain\_makeSame; query secondary\_max**

**namespace** LCT {

**const** int maxn=1e5+7;

**struct** info {

int size;

pii max1,max2;

info(int \_val=-INF,int \_cnt=1,int \_size=1):

size(\_size),max1(make\_pair(\_val,\_cnt)),max2(make\_pair(-INF,0)) {}

void print() {

debug(" debug: infomation: max=(%d,%d)(%d,%d) size=%d**\n**",max1.first,max1.second,max2.first,max2.second,size);

}

};

**struct** tag {

int same,add;*//same:lazy*

tag() {same=-INF; add=0;}

bool tagadd() {**return** (add!=0);}

bool tagsame() {**return** (same!=-INF);}

};

*//info\_merge*

**inline** void merge(info &x,pii value) {

**if** (x.max1.first==value.first)

**return** (void)(x.max1.second+=value.second);

**if** (x.max1<value) swap(x.max1,value);

**if** (x.max2.first==value.first)

**return** (void)(x.max2.second+=value.second);

**if** (x.max2<value) swap(x.max2,value);

}

info merge(**const** info &x,**const** info &y) {

info ret=x;

ret.size+=y.size;

merge(ret,y.max1);

merge(ret,y.max2);

**return** ret;

}

*//info\_update and tag\_update*

**inline** void MakeSame(info &\_info,int value) {

\_info.max1=make\_pair(value,\_info.size);

\_info.max2=make\_pair(-INF,0);

}

**inline** void MakeSame(tag &\_tag,int value) {

\_tag.same=value;

}

**inline** void AddValue(info &\_info,int value) {

\_info.max1.first+=value;

**if** (\_info.max2.first!=-INF) \_info.max2.first+=value;

}

**inline** void AddValue(tag &\_tag,int value) {

\_tag.add+=value;

**if** (\_tag.tagsame()) \_tag.same+=value;

}

**struct** node {

int son[2],fa;

int val;

info chain; tag chaintag;

bool rev,isroot;*//root=1:isroot*

void init(int \_val) {

val=\_val;

chain=info(val);

chaintag=tag();

rev=0; son[0]=son[1]=0;

fa=0; isroot=1;

}

} T[maxn];

void Reverse(int x) {

T[x].rev^=1;

swap(T[x].son[0],T[x].son[1]);

}

void AddValue(node &x,int val) {*//Add\_To\_Node*

x.val+=val;

AddValue(x.chain,val);

AddValue(x.chaintag,val);

}

void MakeSame(node &x,int val) {

x.val=val;

MakeSame(x.chain,val);

MakeSame(x.chaintag,val);

}

void pushup(int x) {

T[x].chain=info(T[x].val);*//clear*

**if** (T[x].son[0])

T[x].chain=merge(T[T[x].son[0]].chain,T[x].chain);

**if** (T[x].son[1])

T[x].chain=merge(T[x].chain,T[T[x].son[1]].chain);

}

void pushdown(int x) {

**if** (T[x].rev) {

**if** (T[x].son[0]) Reverse(T[x].son[0]);

**if** (T[x].son[1]) Reverse(T[x].son[1]);

T[x].rev=0;

}

**if** (T[x].chaintag.tagadd()) {

**if** (T[x].son[0]) AddValue(T[T[x].son[0]],T[x].chaintag.add);

**if** (T[x].son[1]) AddValue(T[T[x].son[1]],T[x].chaintag.add);

T[x].chaintag.add=0;

}

**if** (T[x].chaintag.tagsame()) {

**if** (T[x].son[0]) MakeSame(T[T[x].son[0]],T[x].chaintag.same);

**if** (T[x].son[1]) MakeSame(T[T[x].son[1]],T[x].chaintag.same);

T[x].chaintag.same=-INF;

}

}

void rotate(int x,int kind) {

int y=T[x].fa,z=T[y].fa;

T[y].son[!kind]=T[x].son[kind],T[T[x].son[kind]].fa=y;

T[x].son[kind]=y,T[y].fa=x;

**if** (T[y].isroot) {T[x].isroot=true; T[y].isroot=false;}

**else** T[z].son[T[z].son[1]==y]=x;

T[x].fa=z; pushup(y);

}

void PreChange(int x) {*//change\_from\_root*

**static** int ids[maxn],i,k;

**for** (k=0; !T[x].isroot; k++){

ids[k]=x,x=T[x].fa;

} ids[k++]=x;

rREP(i,k) pushdown(ids[i]);

}

void splay(int x) { *//to root*

PreChange(x);

**while** (!T[x].isroot) {

int y=T[x].fa,z=T[y].fa;

int rx=T[y].son[0]==x,ry=T[z].son[0]==y;

**if** (T[y].isroot) rotate(x,rx);

**else** {

**if** (rx==ry) rotate(y,ry);

**else** rotate(x,rx);

rotate(x,ry);

}

} pushup(x);

}

int access(int x) {

int y=0;

**for** (; x; x=T[x].fa) {

splay(x);

T[T[x].son[1]].isroot=true;

T[x].son[1]=y;

T[y].isroot=false;

y=x; pushup(x);

} **return** y;

}

bool judge(int u,int v) {

**while** (T[u].fa) u=T[u].fa;

**while** (T[v].fa) v=T[v].fa;

**return** u==v;

}

void makeroot(int x) {

access(x); splay(x);

Reverse(x);

}

bool link(int u,int v) {

**if** (judge(u,v)) **return** 1;

makeroot(u); T[u].fa=v;

**return** 0;

}

bool cut(int u,int v) {

makeroot(u); splay(v);

T[T[v].son[0]].fa=T[v].fa;

T[v].fa=0;

T[T[v].son[0]].isroot=true;

T[v].son[0]=0;

pushup(v);

**return** 0;

}

bool add(int u,int v,int val) {

makeroot(u); access(v); splay(v);

AddValue(T[v],val);

**return** 0;

}

bool change(int u,int v,int val) {

makeroot(u); access(v); splay(v);

MakeSame(T[v],val);

**return** 0;

}

pair<int,int> ask(int u,int v) {

makeroot(u); access(v); splay(v);

**return** T[v].chain.max2;

}

};

vector<int> edge[maxn];

void dfs(int x,int fa) {

LCT::T[x].fa=fa;

LCT::T[x].isroot=1;

**for** (int v:edge[x]) **if** (v!=fa) dfs(v,x);

}

int main() {

int x=0;

int T,\_; T=1;

scanf("%d",&T);

FOR(\_,1,T) {

int n,m,i;

scanf("%d%d",&n,&m);

FOR(i,1,n) {

int val;

scanf("%d",&val);

LCT::T[i].init(val);

}

REP(i,n-1) {

int u,v;

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

}

dfs(1,0);

printf("Case #%d:**\n**",++x);

**while** (m--) {

int k;

scanf("%d",&k);

int x,y;

**if** (k==1) {

int x0,y0;

scanf("%d%d%d%d",&x,&y,&x0,&y0);

LCT::cut(x,y);

LCT::link(x0,y0);

} **else** **if** (k==2) {

int val;

scanf("%d%d%d",&x,&y,&val);

LCT::change(x,y,val);

} **else** **if** (k==3) {

int val;

scanf("%d%d%d",&x,&y,&val);

LCT::add(x,y,val);

} **else** **if** (k==4) {

scanf("%d%d",&x,&y);

pair<int,int> t=LCT::ask(x,y);

**if** (t.first==-INF) puts("ALL SAME");

**else** printf("%d %d**\n**",t.first,t.second);

}

}

FOR(i,1,n) edge[i].clear();

}

}

KD树

**//线段树套KD树**

**//KD树,对于子树需要维护区间**

**//时间复杂度:nsqrt(n)**

**//最近距离的话,注意剪枝要减得多,用矩形限制**

**//可以通过对左右估值来确定query顺序**

**//(把query的东西放到外面限制)**

**namespace** KDT {

**const** double alpha=0.75;

**const** int DIM=2;

**struct** point {

int A[DIM],max[DIM],min[DIM];

int l,r; int size;

void init() {

l=r=0; initval();

}

void initval() {

int i; size=1;

REP(i,DIM) min[i]=max[i]=A[i];

}

} T[maxn\*30]; int TOT;

int Cur;

bool cmp(int x,int y) {

**return** T[x].A[Cur]<T[y].A[Cur];

}

void update(int x) {

int i; T[x].initval();

int l=T[x].l,r=T[x].r;

**if** (l) T[x].size+=T[l].size;

**if** (r) T[x].size+=T[r].size;

REP(i,DIM) {

**if** (l) {

T[x].max[i]=max(T[x].max[i],T[l].max[i]);

T[x].min[i]=min(T[x].min[i],T[l].min[i]);

}

**if** (r) {

T[x].max[i]=max(T[x].max[i],T[r].max[i]);

T[x].min[i]=min(T[x].min[i],T[r].min[i]);

}

}

}

int id[maxn],tot;

void build(int &x,int l,int r,int cur) { *//should have id*

x=0; **if** (l>r) **return**;

int m=(l+r)/2; Cur=cur;

nth\_element(id+l,id+m,id+r+1,cmp);

x=id[m];

build(T[x].l,l,m-1,cur^1);

build(T[x].r,m+1,r,cur^1);

update(x);

}

void getid(int x) { *//没有顺序=\_=*

id[++tot]=x;

**if** (T[x].l) getid(T[x].l);

**if** (T[x].r) getid(T[x].r);

}

void rebuild(int &x,int cur) {

tot=0; getid(x);

build(x,1,tot,cur);

}

void insert(int &x,int now,int cur) {

**if** (!x) {x=now; **return**;}

Cur=cur;

**if** (cmp(now,x)) insert(T[x].l,now,cur^1);

**else** insert(T[x].r,now,cur^1);

update(x);

**if** (T[x].size\*alpha+3<max(T[T[x].l].size,T[T[x].r].size))

rebuild(x,cur);

}

void addnode(int &x,int px,int py) {

TOT++; T[TOT].A[0]=px; T[TOT].A[1]=py;

T[TOT].init(); insert(x,TOT,0);

}

int x0,y0,x1,y1;*//check两个=\_=*

int check(int x,int y) {

**return** x0<=x&&x<=x1&&y0<=y&&y<=y1;

}

int ok(point &A) {

**return** check(A.A[0],A.A[1]);

}

int allin(point &A) {

**return** x0<=A.min[0]&&A.max[0]<=x1&&

y0<=A.min[1]&&A.max[1]<=y1;

}

int allout(point &A) {

**return** A.max[0]<x0||x1<A.min[0]||

A.max[1]<y0||y1<A.min[1];

}

int query(int x) {

**if** (!x) **return** 0;

**if** (allin(T[x])) **return** T[x].size;

**if** (allout(T[x])) **return** 0;

int ret=0;

**if** (ok(T[x])) ret++;

**if** (T[x].size==1) **return** ret;

ret+=query(T[x].l);

ret+=query(T[x].r);

**return** ret;

}

}

**const** int MAX=1e9+7;

**struct** Tnode {

int l,r,KD\_root;

Tnode() {l=r=KD\_root=0;}

} T[maxn\*30]; int cnt;

void update(int &x,int px,int py,int pos,int L,int R) {

**if** (!x) x=++cnt;

KDT::addnode(T[x].KD\_root,px,py);

**if** (L==R) **return**;

int mid=(L+R)/2;

**if** (pos<=mid) update(T[x].l,px,py,pos,L,mid);

**else** update(T[x].r,px,py,pos,mid+1,R);

}

int query(int x,int k,int L,int R) {

**if** (!x) **return** 0;

**if** (L==R) **return** L;

int mid=(L+R)/2;

**if** (T[x].r) {

int rk=KDT::query(T[T[x].r].KD\_root);

**if** (rk<k) **return** query(T[x].l,k-rk,L,mid);

**return** query(T[x].r,k,mid+1,R);

} **return** query(T[x].l,k,L,mid);

}

char buffer[36000000],\*buf=buffer;

void read(int &x) {

**for** (x=0; \*buf<48; ++buf);

**while** (\*buf>=48)x=x\*10+\*buf-48,++buf;

}

int n,q;

int i,j,k;

int root,lastans;

int main() {

fread(buffer,1,36000000,stdin);

read(n); read(q); KDT::TOT=0;

FOR(i,1,q) {

int op;

read(op);

**if** (op==1) {

int x,y,v;

read(x); read(y); read(v);

x^=lastans; y^=lastans; v^=lastans;

update(root,x,y,v,0,MAX);

} **else** {

int x1,y1,x2,y2,k;

read(x1); read(y1); read(x2); read(y2); read(k);

x1^=lastans; y1^=lastans;

x2^=lastans; y2^=lastans;

k^=lastans;

KDT::x0=x1; KDT::y0=y1;

KDT::x1=x2; KDT::y1=y2;

lastans=query(root,k,0,MAX);

**if** (!lastans) puts("NAIVE!ORZzyz.");

**else** printf("%d**\n**",lastans);

}

}

}

莫队

sort时可以按照&1左往右or右往左

树上莫队(套分块)

**//http://codeforces.com/gym/100962/attachments**

**//题意是求路径上最小没出现数字**

**//主要思路是分类,每个点进出各算一次可以消除影响**

**//点的直接加个lca即可**

**const** int SIZE=500;

vector<pair<int,int> > edge[maxn];

int cl[maxn],cr[maxn],val[maxn],dfn[maxn<<1];

int tot;

int dfs(int x,int fa) {

cl[x]=++tot; dfn[tot]=x;

**for** (**auto** now:edge[x]) **if** (now.first!=fa) {

dfs(now.first,x);

val[now.first]=now.second;

} cr[x]=++tot; dfn[tot]=x;

}

int block[maxn<<1];

**struct** node {

int l,r,id;

} Q[maxn];

int cmp(node a,node b) {

**if** (block[a.l]==block[b.l]) **return** a.r<b.r;

**return** block[a.l]<block[b.l];

}

bool vis[maxn];

int cnt[maxn],cur[maxn];*//block,now*

void change(int x) {

x=dfn[x]; vis[x]^=1;

**if** (vis[x]) {

**if** (!cur[val[x]]) cnt[block[val[x]]]++;

cur[val[x]]++;

} **else** {

cur[val[x]]--;

**if** (!cur[val[x]]) cnt[block[val[x]]]--;

}

}

int ans[maxn];

int L,R;

int main() {

int n,q;

int i;

scanf("%d%d",&n,&q);

FOR(i,0,n\*2+1) block[i]=i/SIZE;

REP(i,n-1) {

int u,v,len;

scanf("%d%d%d",&u,&v,&len); len=min(len,n+1);

edge[u].push\_back(make\_pair(v,len));

edge[v].push\_back(make\_pair(u,len));

}

val[1]=n+1; dfs(1,0);

REP(i,q) {

int a,b;

scanf("%d%d",&a,&b);

**if** (cl[a]>cl[b]) swap(a,b);

**if** (cr[a]>cr[b]) Q[i].l=cl[a]+1,Q[i].r=cl[b];

**else** Q[i].l=cr[a],Q[i].r=cl[b];

Q[i].id=i;

}

sort(Q,Q+q,cmp);

L=1; R=0;

REP(i,q) {

**while** (L<Q[i].l) {change(L); L++;}

**while** (R>Q[i].r) {change(R); R--;}

**while** (L>Q[i].l) {L--; change(L);}

**while** (R<Q[i].r) {R++; change(R);}

int now=0;

**while** (cnt[now]==SIZE) now++;

now\*=SIZE;

**while** (cur[now]) now++;

ans[Q[i].id]=now;

}

REP(i,q) printf("%d**\n**",ans[i]);

}

回滚莫队套分块

**//北京区域赛,题意是l->r的所有内部边的并查集啥的**

**//回滚分块(然而我没回滚,记录了一下)**

**//queries按照左端点排序(有边的要按照我这种方式来排,否则菊花图会卡死)**

**//cmpu的时候v要倒着,因为要让块外的不受左边影响(same\_l and small\_r)**

**//然后对于左端点在block内部的所有query,按右端点往右走,走到头即可**

**//这个做法就是按照左分块,然后把右边有效的加进去,再把左边的加进去就行了**

int SIZE;

**struct** node {

int u,v,id,o;

node() {};

node(int \_u,int \_v,int \_id=0):u(\_u),v(\_v),id(\_id) {};

} to[maxn],re[maxn],queries[maxn];

int BID[maxn],L[maxn];

bool cmpu(node A,node B) {

**if** (A.u!=B.u) **return** A.u<B.u;

*//区间为了避免漏掉r小的*

**if** (A.v!=B.v) **return** A.v>B.v;

**return** A.id>B.id;

}

bool cmpv(node A,node B) {

**if** (A.v!=B.v) **return** A.v<B.v;

**if** (A.u!=B.u) **return** A.u<B.u;

**return** A.id<B.id;

}

bool cmpQ(node A,node B) {

**if** (A.o!=B.o) **return** A.o<B.o;

**if** (A.v!=B.v) **return** A.v<B.v;

**if** (A.u!=B.u) **return** A.u<B.u;

**return** A.id<B.id;

}

int fa[maxn],size[maxn];

LL Ans[maxn];

**inline** int getfa(int x) {

**if** (fa[x]==x) **return** x;

**return** fa[x]=getfa(fa[x]);

}

int FA[maxn],SZ[maxn],PID[maxn];

**inline** int getFA(int x) {

**if** (FA[x]==x) **return** x;

**return** FA[x]=getFA(FA[x]);

}

**inline** void update(int u,int pid) {

**if** (PID[u]!=pid) {

int f=getfa(u);

**if** (PID[f]!=pid) {

FA[f]=f;

PID[f]=pid;

SZ[f]=size[f];

} PID[u]=pid; FA[u]=f;

}

} int tot=0;

LL now;

int main() {

int T;

scanf("%d",&T);

**while** (T--) {

int n,m,q,i,j,k;

scanf("%d%d%d",&n,&m,&q);

**if** (q==0) SIZE=m; **else** SIZE=m/sqrt(q)\*2;

**if** (!SIZE) SIZE++;

FOR(i,0,(m+1)/SIZE) L[i]=0;

FOR(i,1,m+1) {BID[i]=i/SIZE; **if** (!L[i/SIZE]) L[i/SIZE]=i;}

FOR(i,1,m) {

int u,v;

scanf("%d%d",&u,&v);

**if** (u>v) swap(u,v);

to[i]=node(u,v);

re[i]=node(u,v);

} sort(to+1,to+m+1,cmpv);

sort(re+1,re+m+1,cmpu);

FOR(i,1,m) {

to[i].o=BID[lower\_bound(re+1,re+1+m,to[i],cmpu)-re];

re[i].o=BID[i];

}

FOR(i,1,q) {

int u,v;

scanf("%d%d",&u,&v);

**if** (u>v) swap(u,v);

queries[i]=node(u,v,i);

queries[i].o=BID[lower\_bound(re+1,re+1+m,queries[i],cmpu)-re];

} sort(queries+1,queries+q+1,cmpQ);

FOR(i,1,q) {

**if** (i==1||queries[i].o!=queries[i-1].o) { *//initialize*

FOR(j,1,n) fa[j]=j,size[j]=1;

j=1; now=0;

}

**for** (; j<=m&&to[j].v<=queries[i].v; j++) {

**if** (to[j].o>queries[i].o) {*//sorted by l*

node &e=to[j];

int x=getfa(e.u),y=getfa(e.v);

**if** (x==y) **continue**; fa[x]=y;

now+=(LL)size[x]\*size[y];

size[y]+=size[x];

}

}

LL ans=now; tot++;

**for** (k=L[queries[i].o]; k<=m&&BID[k]==queries[i].o; k++) {

**if** (queries[i].u<=re[k].u&&re[k].v<=queries[i].v) {

node &e=re[k];

update(e.u,tot); update(e.v,tot);

int x=getFA(e.u),y=getFA(e.v);

**if** (x==y) **continue**; FA[x]=y;

ans+=(LL)SZ[x]\*SZ[y];

SZ[y]+=SZ[x];

}

}

Ans[queries[i].id]=ans;

}

FOR(i,1,q) printf("%lld**\n**",Ans[i]);

}

}

带修改莫队

**//change常数大时size可以增大**

**//sort时先block,改变顺序可以降低常数**

**//n^2/3,注意常数**

**//注意change时间时排的顺序**

**const** int SIZE=2500;

**struct** queries{

int l,r,t;*//pre*

queries(){};

queries(int \_l,int \_r,int \_t):l(\_l),r(\_r),t(\_t){};

}Q[maxn],S[maxn];

int n,m,q;

int i,j,k;

int a[maxn];

int BLOCK[maxn];

bool cmp(queries &A,queries &B){

**if** (BLOCK[A.l]!=BLOCK[B.l]) **return** BLOCK[A.l]<BLOCK[B.l];

**if** (BLOCK[A.r]!=BLOCK[B.r]) **return** BLOCK[A.r]<BLOCK[B.r];

**return** (A.t<B.t)^((BLOCK[A.l]^BLOCK[A.r])&1);

}vector<int> V;

**inline** int getid(int x){**return** lower\_bound(V.begin(),V.end(),x)-V.begin()+1;}

int L,R,T;

int num[maxn],cnt[maxn];

**inline** void add(int pos){

int &T=num[a[pos]];

cnt[T]--;T++;cnt[T]++;

}**inline** void del(int pos){

int &T=num[a[pos]];

cnt[T]--;T--;cnt[T]++;

}**inline** void change(int pos,int val){

**if** (L<=pos&&pos<=R){del(pos),a[pos]=val,add(pos);}

**else** a[pos]=val;

}

int ans[maxn];

int main(){

scanf("%d%d",&n,&q);

FOR(i,1,n) scanf("%d",&a[i]),V.push\_back(a[i]);

FOR(i,1,q){

int op,l,r;

scanf("%d%d%d",&op,&l,&r);

**if** (op==1){

Q[i]=queries(l,r,i);

}**if** (op==2) {

S[i]=queries(l,r,a[l]);a[l]=r;

V.push\_back(a[l]);

}

}sort(V.begin(),V.end());

V.erase(unique(V.begin(),V.end()),V.end());

FOR(i,1,n) a[i]=getid(a[i]);

FOR(i,1,q) **if** (S[i].t) S[i].r=getid(S[i].r),S[i].t=getid(S[i].t);

FOR(i,1,max(n,q)) BLOCK[i]=i/SIZE;

sort(Q+1,Q+q+1,cmp);

L=1;R=0;T=q;cnt[0]=INF;

FOR(i,1,q) **if** (Q[i].t){

**while** (T<Q[i].t){T++;**if** (S[T].t) change(S[T].l,S[T].r);}

**while** (T>Q[i].t){**if** (S[T].t) change(S[T].l,S[T].t);T--;}

**while** (L<Q[i].l){del(L);L++;}

**while** (R>Q[i].r){del(R);R--;}

**while** (L>Q[i].l){L--;add(L);}

**while** (R<Q[i].r){R++;add(R);}

int now=0;

**while** (cnt[now]) now++;

ans[Q[i].t]=now;

}FOR(i,1,q) **if** (ans[i]) printf("%d**\n**",ans[i]);

}

二次离线莫队

**// 题意: 区间A[i]%B[i]=0的对数**

**// 做法: 第十四分块, 二次离线莫队然后再离线bigsmall算贡献, 用fractional cascading(分散层叠也行)**

**// 分散层叠: 每一层保存这层所有信息+一半下一层的信息(position)**

**// 那么只lowerbound, 然后while回去就行了, 复杂度k(层数)+log(lowerbound)**

**// 还是得老老实实算, 不能按贡献算, 因为有可能向前和向后的贡献不一致(?)**

**// [l,r] to [l,r'] = [1, r(x)] to [(x+1),r'] - [1,l-1] to [r+1,r'] = sum\_l[r']-sum\_l[r]-[1,l-1] to [r+1,r']**

**// [l,r] to [l',r] =-[l(x+1),n] to [(x),l'-1] + [r+1,n] to [l,l'-1] = sum\_r[l']-sum\_r[l]+[r+1,n] to [l,l'-1](from r+1,rev)**

**// 下面式子是不对称的(贡献一致可以减常数)**

**// [l,r] to [l',r] = [1, (x)] to [(x),l'-1] - [1, r] to [l,l'-1] = sumlx[l'-1]-sumlx[l-1]-[1,r] to [l,l'-1]**

**// 二次离线的作用是如果可以O(1)查询某个点的值, O(sqrt)更新, 就可以for一遍范围直接加起来!**

**// 二次离线之后还是要再离线算其他的贡献**

**struct** node {

int l,r,pos,type;

} Q[maxn];

vector<node> Ql[maxn],Qr[maxn];

int A[maxn];

int BLOCK[maxn];*// 莫队sqrt*

**const** int SIZE=300;

int sumx[maxn],sumy[maxn];

ll suml[maxn],sumr[maxn];

ll base[maxn],ans[maxn];*//more*

int C[maxn];*//front*

int SIZE\_B=50;*// big\_small, 这个常数也太大了*

vector<int> fac[maxn];

int main() {

int i;

int n,q;

scanf("%d%d",&n,&q);

FOR(i,1,n) scanf("%d",&A[i]);

FOR(i,1,n) BLOCK[i]=i/SIZE;

FOR(i,1,q) scanf("%d%d",&Q[i].l,&Q[i].r),Q[i].pos=i;

sort(Q+1,Q+1+q,[&](node &x,node &y) {

**if** (BLOCK[x.l]!=BLOCK[y.l]) **return** BLOCK[x.l]<BLOCK[y.l];

**else** **return** bool((x.r<y.r)^(BLOCK[x.l]&1));

});

int l=1,r=0;

FOR(i,1,q) {

int l\_l=Q[i].l,r\_r=Q[i].r,base;

base=-1; **if** (r>r\_r) swap(r,r\_r),base\*=-1;

**if** (r<r\_r) Ql[l-1].push\_back(node{r+1,r\_r,i,base}); r=Q[i].r;

base=1; **if** (l>l\_l) swap(l,l\_l),base\*=-1;

**if** (l<l\_l) Qr[r+1].push\_back(node{l,l\_l-1,i,base}); l=Q[i].l;

}

int k;

FOR(i,1,100000) {

**for** (int k=i; k<=100000; k+=i) fac[k].push\_back(i);

} FOR(i,1,100000) C[i]=0;

*//first*

FOR(i,1,n) {*//x=ky or {kx=y and x>sqrt}*

suml[i]=C[A[i]];

**for** (int k:fac[A[i]]) C[k]++;

**if** (A[i]>SIZE\_B)

**for** (int k=A[i]; k<=100000; k+=A[i]) C[k]++;

**for** (**auto** now:Ql[i])

FOR(k,now.l,now.r) base[now.pos]+=C[A[k]]\*now.type;

} FOR(i,1,100000) C[i]=0;

rFOR(i,1,n) {

sumr[i]=C[A[i]];

**for** (int k:fac[A[i]]) C[k]++;

**if** (A[i]>SIZE\_B)

**for** (int k=A[i]; k<=100000; k+=A[i]) C[k]++;

**for** (**auto** now:Qr[i])

FOR(k,now.l,now.r) base[now.pos]+=C[A[k]]\*now.type;

}

*//second, BLOCK*

FOR(k,1,SIZE\_B) {

FOR(i,1,n) {*//注意r可能要另算,这个题对称所以ok*

sumx[i]=sumx[i-1]+(A[i]==k);

sumy[i]=sumy[i-1]+(A[i]%k==0);

}

FOR(i,1,n) {

**if** (A[i]%k==0) suml[i]+=sumx[i-1],sumr[i]+=sumx[n]-sumx[i];

**for** (**auto** now:Ql[i])

base[now.pos]+=(ll)sumx[i]\*(sumy[now.r]-sumy[now.l-1])\*now.type;

**for** (**auto** now:Qr[i])

base[now.pos]+=(ll)(sumx[n]-sumx[i-1])\*(sumy[now.r]-sumy[now.l-1])\*now.type;

}

}

FOR(i,1,n) suml[i]+=suml[i-1];

rFOR(i,1,n) sumr[i]+=sumr[i+1];

l=1,r=0;

FOR(i,1,q) {

int l\_l=Q[i].l,r\_r=Q[i].r,x=Q[i].pos;

base[i]=suml[r\_r]-suml[r]+sumr[l\_l]-sumr[l]+base[i]+base[i-1];

ans[x]=base[i]+r\_r-l\_l+1;

l=l\_l; r=r\_r;

}

FOR(i,1,q) printf("%lld**\n**",ans[i]);

}

用set维护凸包

**/\* 这是抄的维护上半凸壳 \*/**

**// 最大值,query的k要求>0**

bool Q;

**struct** Line {

**mutable** LL a,b,k;

bool **operator**<(**const** Line &o)**const** {

**return** Q?k<o.k:a<o.a;

}

};

**struct** convexHull:**public** multiset<Line> {

LL div(LL a,LL b) {

**return** a/b-((a^b)<0&&a%b);

}

bool getK(iterator x,iterator y) {

**if** (y==end()) {x->k=INFF; **return** 0;}

**if** (x->a==y->a) x->k=x->b>y->b?INFF:-INFF;

**else** x->k=div(y->b-x->b,x->a-y->a);

**return** x->k>=y->k;

}

void insPos(LL a,LL b) {

**auto** z=insert({a,b,0}); **auto** y=z++,x=y;

**while** (getK(y,z)) z=erase(z);

**if** (y!=begin()&&getK(--x,y)) getK(x,erase(y));

**while** ((y=x)!=begin()&&(--x)->k>=y->k)

getK(x,erase(y));

}

LL query(LL x) {

assert(size());

Q=1; **auto** now=lower\_bound({0,0,x}); Q=0;

**return** now->a\*x+now->b;

}

};

李超树

**//李超树最主要的作用在于维护线段,而不是直线!**

**//维护l<=x<=r时下放线段,时间复杂度两个log!**

**//这里是最大值**

double cross(double k1,double b1,double k2,double b2) {

**if** (abs(k1-k2)<eps) **return** INF;

**return** (b2-b1)/(k1-k2);

}

int flag[maxn\*4];

double tagk[maxn\*4],tagb[maxn\*4];

void ins(int x,double k,double b,int l,int r,int id,int L,int R) {

**if** (l<=L&&R<=r) {

**if** (!flag[x]) tagk[x]=k,tagb[x]=b,flag[x]=id;

**else** {

int mid=(L+R)/2;

double ini\_l=tagk[x]\*L+tagb[x],now\_l=k\*L+b;

double ini\_r=tagk[x]\*R+tagb[x],now\_r=k\*R+b;

**if** (ini\_l>=now\_l&&ini\_r>=now\_r) **return**;

**if** (ini\_l<=now\_l&&ini\_r<=now\_r) tagk[x]=k,tagb[x]=b,flag[x]=id;

**else** {

double pos=cross(k,b,tagk[x],tagb[x]);*//交点x坐标*

**if** ((pos<=mid&&ini\_l>=now\_l)||(pos>mid&&ini\_r>=now\_r)) { *//坐标低的下放,平的直接留下就行*

swap(tagk[x],k);

swap(tagb[x],b);

swap(flag[x],id);

} **if** (pos<=mid) ins(x<<1,k,b,l,r,id,L,mid);

**else** ins(x<<1|1,k,b,l,r,id,mid+1,R);

}

}

} **else** {

int mid=(L+R)/2;

**if** (l<=mid) ins(x<<1,k,b,l,r,id,L,mid);

**if** (mid<r) ins(x<<1|1,k,b,l,r,id,mid+1,R);

}

}

double ans; int id;

void que(int x,int pos,int L,int R) {

**if** (flag[x]) {

double now=tagk[x]\*pos+tagb[x];

**if** (now-ans>eps||(now-ans>-eps&&id>flag[x])) {

ans=now,id=flag[x];

}

}

**if** (L==R) **return**;

int mid=(L+R)/2;

**if** (pos<=mid) que(x<<1,pos,L,mid);

**else** que(x<<1|1,pos,mid+1,R);

}

线性基(套路)

**namespace** LB {

**typedef** long long BaseType;

**const** int MaxBit=63;

**struct** L\_B {

BaseType b[MaxBit]; bool have\_0;

L\_B() {clear();}

void clear() {memset(b,0,**sizeof**(b)); have\_0=0;}

BaseType XORMIN(BaseType x) {

int i;

rREP(i,MaxBit) **if** ((b[i]^x)<x) x^=b[i];

**return** x;

}

BaseType XORMAX(BaseType x) {

int i;

rREP(i,MaxBit) **if** ((b[i]^x)>x) x^=b[i];

**return** x;

}

void insert(BaseType x) {

int i;

**if** (!have\_0&&!XORMIN(x)) have\_0=1;

rREP(i,MaxBit) **if** ((x>>i)&1) {

**if** (!b[i]) b[i]=x; x^=b[i];

}

}

void rebuild() {

int i,j;

rREP(i,MaxBit) rREP(j,i) **if** ((b[i]>>j)&1) b[i]^=b[j];

}

BaseType querykth(BaseType k) {

BaseType ret=0; int i; k-=have\_0;

REP(i,MaxBit) **if** (b[i]) {**if** (k&1) ret^=b[i]; k>>=1;}

**if** (k) **return** -1;

**return** ret;

}

} A;

*//求交 merge的思路: 只要A中merge之后的线性无关组*

L\_B merge(**const** L\_B &A,**const** L\_B &B) {

int i,j; L\_B ret; ret.clear();

**static** BaseType base[MaxBit],tmp[MaxBit];*//previous\_A*

REP(i,MaxBit) tmp[i]=A.b[i],base[i]=1ll<<i;

REP(i,MaxBit) **if** (B.b[i]) {*//正者反着应该没区别*

BaseType now=B.b[i];

bool okay=1; BaseType k=0;*//base; A*

rREP(j,MaxBit) **if** ((now>>j)&1) {

**if** (tmp[j]) {

now^=tmp[j]; k^=base[j];

} **else** {

tmp[j]=now; base[j]=k; okay=0; **break**;

}

}

**if** (okay) {

BaseType should=0;

REP(j,MaxBit) **if** ((k>>j)&1) should^=A.b[j];

ret.insert(should);

}

}

**return** ret;

}

}

手写BITSET

**struct** BITSET {

vector<ULL> V;

void set(int x,int k) {

assert((int)V.size()>x/64);

**if** (k) V[x/64]|=1ull<<(x&63);

**else** V[x/64]&=~(1ull<<(x&63));

}

void resize(int x) {

V.resize((x-1)/64+1,0);

}

int get(int x) {

**return** (V[x/64]>>(x&63))&1;

}

bool **operator** < (**const** BITSET &B) **const** {

int i;

REP(i,(int)V.size()) **if** (V[i]!=B.V[i]) **return** V[i]<B.V[i];

**return** 0;

}

BITSET **const** doit(int size,int F[65536]) **const** {*//相邻两位合并*

BITSET ret; int i;

ret.resize(size/2);

REP(i,(int)V.size()) {

**if** (i&1) {

ret.V[i/2]|=((ULL)F[V[i]&65535]<<32)

|((ULL)F[(V[i]>>16)&65535]<<40)

|((ULL)F[(V[i]>>32)&65535]<<48)

|((ULL)F[(V[i]>>48)]<<56);

} **else** {

ret.V[i/2]|=((ULL)F[V[i]&65535])

|((ULL)F[(V[i]>>16)&65535]<<8)

|((ULL)F[(V[i]>>32)&65535]<<16)

|((ULL)F[(V[i]>>48)]<<24);

}

} **return** ret;

}

void print() {

int i;

REP(i,(int)V.size()) pr2(V[i],64);

}

};

杨表

**//题意: 选5个subsquence和最大**

**//杨表: 单调\*\*子序列个数最大多少个**

**//杨表做法: 直接替换原数列中比这个大的位置,**

**// 然后直接将多出来的往下放即可**

**//正确性: 可以将这个点后面连的所有东西放下边,**

**// 相当于连个边,相当于最优选择**

ll ans;

map<int,ll> MP[5];*//pos,cnts*

void update(int x,int y,int dep) {*//x,cnt*

**if** (dep==5) **return**;

**while** (y) {

map<int,ll>::iterator it=MP[dep].upper\_bound(x);

**if** (it==MP[dep].end()) {

ans+=y;

MP[dep][x]+=y; **break**;

} pair<int,int> now=\*it;

MP[dep].erase(it);

ll down=min(now.second,y);

y-=down; now.second-=down;

**if** (now.second) MP[dep][now.first]+=now.second;

MP[dep][x]+=down;

update(now.first,down,dep+1);

}

}

int main(){

int T,\_; T=1;

scanf("%d",&T);

FOR(\_,1,T){

int i,n;

scanf("%d",&n); ans=0;

REP(i,5) MP[i].clear();

FOR(i,1,n) {

int k; scanf("%d",&k);

update(k,k,0);

printf("%lld%c",ans," **\n**"[i==n]);

}

}

}

# 图论

二分图匹配

**//最小不相交路径覆盖<=>节点数-拆点以后二分图最大匹配**

**//最小相交路径覆盖<=>所有能走到的节点连边，然后节点数-拆点以后匹配**

vector<int>edge[N];

int used[N];

注意数组的标号，必须满足二分图的条件

int matching[N];

bool dfs(int u){

int v,i;

REP(i,edge[u].size()){

v=edge[u][i];

**if** (!used[v]){

used[v]=1;

**if** (matching[v]==-1||dfs(matching[v])){

matching[v]=u;

matching[u]=v;

**return** 1;

}

}

}

**return** 0;

}

int DFS(){

int ans=0;

memset(matching,-1,**sizeof**(matching));

int u;

FOR(u,1,n){

**if** (matching[u]==-1){

memset(used,0,**sizeof**(used));

**if** (dfs(u)) ans++;

}

}

**return** ans;

}

注意数组的标号，必须满足二分图的条件

queue<int> Q;

int prev[N];*//两格*

int matching[N];*//结果*

int check[N];*//matchright*

int BFS(){

int ans=0;

memset(matching,-1,**sizeof**(matching));

memset(check,-1,**sizeof**(check));

FOR(i,1,n){

**if** (matching[i]==-1){

**while** (!Q.empty()) Q.pop();

Q.push(i);

prev[i]=-1;

bool flag=false;

**while** (!Q.empty()&&!flag){

int u=Q.front();Q.pop();

**for** (j=0;!flag&&j<edge[u].size();j++){

int v=edge[u][j];

**if** (check[v]!=i){

check[v]=i;

Q.push(matching[v]);

**if** (matching[v]!=-1) prev[matching[v]]=u;

**else**{

flag=1;

int d=u,e=v;

**while** (d!=-1){

int t=matching[d];

matching[d]=e;

matching[e]=d;

d=prev[d];

e=t;

}

}

}

}

}

**if** (matching[i]!=-1) ans++;

}

}

**return** ans;

}

Hall定理

**// 题意: N个人,M个椅子,每个人只能坐[1,Li]|[Ri,M],求最多能坐多少人**

**// hall定理: 二分图; A->B (A<B)完美匹配当且仅当A中每k个在B中连着有至少k个点**

**// 引理(不常用): 如果A中每个连着最少t条边, B中每个连着最多t条边, 那么存在完美匹配; t任意**

**// 对于这个题来说: 最终选择的座位比人少; 任意座位集合A; B: [1,Lx][Rx,M]**

**// 座位当作A, 用定理, 所有区间满足: 对人的集合B, A->B ,|A|>=$加完边$的B 求下|A|-|B|>=0**

**// 枚举A的端点, 求: B的size最大值即可!**

int MIN[maxn],lazy[maxn];

**inline** void add(int x,int val){

lazy[x]+=val;MIN[x]+=val;

}void update(int x,int l,int r,int val,int L,int R){

**if** (l<=L&&R<=r){add(x,val);**return**;}

**if** (lazy[x]){

add(x<<1,lazy[x]);

add(x<<1|1,lazy[x]);

lazy[x]=0;

}int mid=(L+R)/2;

**if** (l<=mid) update(x<<1,l,r,val,L,mid);

**if** (mid<r) update(x<<1|1,l,r,val,mid+1,R);

MIN[x]=min(MIN[x<<1],MIN[x<<1|1]);

}int n,m;

vector<int> have[maxn];

int i,j,k;

int l,r;

int ans;

int main(){

scanf("%d%d",&n,&m);

FOR(i,1,n){

scanf("%d%d",&l,&r);

have[l].push\_back(r);

}

FOR(i,1,m) update(1,i,i,m-i+1,1,m+1);

ans=min(0,m-n);*//为啥会有这个问题呢*

FOR(i,0,m){

**if** (i!=0) update(1,i+1,m+1,1,1,m+1);

**for** (int r:have[i])

update(1,i+1,r,-1,1,m+1);

ans=min(ans,MIN[1]);

}printf("%d**\n**",-ans);

}

KM 二分图最大权匹配

ll g[maxn][maxn];

ll lx[maxn],ly[maxn],slack[maxn];

int linky[maxn],par[maxn];

bool visy[maxn];

void augment(int root){

std::fill(visy+1,visy+n+1,false);

std::fill(slack+1,slack+n+1,INFF);

int py; linky[py=0]=root;

**do**{

visy[py]=true;

int x=linky[py],\_y=0,y; ll d=INFF;

FOR(y,1,n) **if** (!visy[y]){

int tmp=lx[x]+ly[y]-g[x][y];

**if** (tmp<slack[y]){

slack[y]=tmp; par[y]=py;

} **if** (slack[y]<d) {

d=slack[y]; \_y=y;

}

} FOR(y,0,n){

**if** (visy[y]){

lx[linky[y]]-=d;

ly[y]+=d;

} **else** slack[y]-=d;

} py=\_y;

} **while** (linky[py]!=-1);

**do** {

int pre=par[py];

linky[py]=linky[pre];

py=pre;

} **while** (py);

}

ll KM() {

int i,y;

FOR(i,1,n) {

lx[i]=0; ly[i]=0; linky[i]=-1;

FOR(y,1,n) max\_(lx[i],g[i][y]);

} ll ret=0;

FOR(i,1,n) augment(i);

FOR(i,1,n) ret+=g[linky[i]][i];

**return** ret;

}

int main() {

int T,\_T;

scanf("%d",&T);

FOR(\_T,1,T) {

scanf("%d",&n);

int i,j;

FOR(i,1,n) FOR(j,1,n) {

int x;

scanf("%d",&x);

g[i][j]=-x;

} ll ans=-KM();

*// printf("%d\n",ans);*

printf("Case #%d: %I64d**\n**",\_T,ans);

}

}

最短路

**Dijkstra：略**

**SPFA DFS(只用于判负环)**

**struct** node{

int n,d;

node(){}

node(int a,int b):n(a),d(b){}

bool **operator**<(**const** node&a)**const**{

**if** (d==a.d) **return** n<a.n;

**return** d>a.d;

}

};

vector<node> edge[maxn];

int dis[maxn],n,m;

bool vis[maxn];

bool spfa(int u){

int i;

vis[u]=1;

REP(i,edge[u].size()){

node v=edge[u][i];

**if** (dis[u]+v.d<dis[v.n]){

dis[v.n]=dis[u]+v.d;

**if** (vis[v.n]) **return** 1;

**else** {

dis[v.n]=dis[u]+v.d;

**if** (spfa(v.n)) **return** 1;

}

}

}

vis[u]=0;

**return** 0;*//judge negative ring*

}

int s,t;

int u,v,len;

int main(){

int i,j,k;

**while** (~scanf("%d%d",&n,&m)){

FOR(i,1,n) edge[i].clear();

REP(i,m){

scanf("%d%d%d",&u,&v,&len);

edge[u].push\_back(node(v,len));

edge[v].push\_back(node(u,len));

}

FOR(i,1,n) dis[i]=INF;dis[1]=0;

FOR(i,1,n) vis[i]=0;

spfa(1);

FOR(i,2,n) printf("%d ",dis[i]==INF?-1:dis[i]);

puts("");

}

}

差分约束系统

**//主要在于建图**

**//连边u->v,len <=> val(v)-val(u)<=len**

**//其他的都要化成这种形式 int n,m;**

**//最好spfa!(可能负环)**

01分数规划

**//2017-harbin-K**

**//选出k个区间，使得这k个区间全覆盖，而且sigmaA/sigmaB最小**

**//俩log dp TLE**

**//做法：建最短路，01分数规划玄学过题**

**struct** node {

int n;

double d;

node() {}

node(int \_n,double \_d):n(\_n),d(\_d) {};

bool **operator**<(**const** node &A)**const** {

**if** (d==A.d) **return** n<A.n;

**return** d>A.d;

}

};

**struct** node\_e {

int n,A,B;

double d;

node\_e(int \_n,int \_A,int \_B,double \_d):n(\_n),A(\_A),B(\_B),d(\_d) {}

};

vector<node\_e> edge[maxn];

int dis[maxn];

int preA[maxn],preB[maxn];

void dij(int s,int n) {

int i;

FOR(i,1,n) dis[i]=INF;

dis[s]=0;

priority\_queue<node> Q;

Q.push(node(s,dis[s]));

**while** (Q.size()) {

node x=Q.top();

Q.pop();

**for** (**auto** &y:edge[x.n]) {

**if** (dis[y.n]>x.d+y.d) {

dis[y.n]=x.d+y.d;

Q.push(node(y.n,dis[y.n]));

preA[y.n]=preA[x.n]+y.A;

preB[y.n]=preB[x.n]+y.B;

}

}

}

}

int n,t;

int S[maxn],T[maxn],A[maxn],B[maxn];

double check(double x) {

int i;

double allA=0,allB=0;

FOR(i,1,t+1)

edge[i].clear();

FOR(i,1,n) {

**if** (A[i]-B[i]\*x<=0) {

allA+=A[i];

allB+=B[i];

edge[S[i]].emplace\_back(node\_e(T[i]+1,0,0,0));

} **else** edge[S[i]].emplace\_back(node\_e(T[i]+1,A[i],B[i],A[i]-B[i]\*x));

}

FOR(i,1,t)

edge[i+1].emplace\_back(node\_e(i,0,0,0));

dij(1,t+1);

allA+=preA[t+1];

allB+=preB[t+1];

**return** allA/allB;

}

int main() {

int i,j,m,x,\_T;

scanf("%d",&\_T);

**while** (\_T--) {

scanf("%d%d",&n,&t);

FOR(i,1,n)

scanf("%d%d%d%d",&S[i],&T[i],&A[i],&B[i]);

double ans=100;

**while** (1) {

double now=check(ans);

**if** (abs(now-ans)<0.001) **break**;

ans=now;

}

printf("%.3lf**\n**",ans);

}

**return** 0;

}

切比雪夫(曼哈顿)距离最小生成树

**//最小曼哈顿距离生成树**

**//按照45度4个方向排序，最近的两个点连边即可**

**//最大曼哈顿距离生成树是维护最远的点的距离（四个方向的）**

**//Kruskal(有道分治题用的Boruvka，和这个思想也类似)**

**//注意理解并查集的内涵，每次找最短的路也可以通过其他方式来找到**

**切比雪夫距离转曼哈顿距离：**

**切比雪夫距离：max(|x1-x2|,|y1-y2|);**

**曼哈顿距离：|x1-x2|+|y1-y2|**

**转化方式：旋转45度然后/2**

**(x,y)->((x+y)/2,(x-y)/2)**

**曼哈顿距离最小生成树：**

**按照45度4个方向排序，最近的两个点连边即可**

**swap方向代码：**

int a[MAXN],b[MAXN];

tot = 0;

**for** (int dir = 0; dir < 4; dir++) {

*//4种坐标变换*

**if** (dir == 1 || dir == 3) {

**for** (int i = 0; i < n; i++) swap(p[i].x,p[i].y);

} **else** **if** (dir == 2) {

**for** (int i = 0; i < n; i++) p[i].x = -p[i].x;

}

sort(p,p+n,cmp);

**for** (int i = 0; i < n; i++)

a[i] = b[i] = p[i].y - p[i].x;

sort(b,b+n);

int m = unique(b,b+n) - b;

**for** (int i = 1; i <= m; i++) bit[i].init();

**for** (int i = n-1 ; i >= 0; i--) {

int pos = lower\_bound(b,b+m,a[i]) - b + 1;

int ans = ask(pos,m);

**if** (ans != -1)

addedge(p[i].id,p[ans].id,dist(p[i],p[ans]));

update(pos,p[i].x+p[i].y,i);

}

}

笛卡尔树

**2017hdu多校1 给定区间:**

int L[maxn],R[maxn];

pii S[maxn]; int top;

int fa[maxn],id[maxn];*//id: topo*

bool buildtree(int n){*//return 1: wrong!*

**static** int i,l[maxn],r[maxn],p[maxn],top;

FOR(i,1,n) id[i]=i; top=0;

sort(id+1,id+1+n,[](int i,int j){

**if** (L[i]!=L[j]) **return** L[i]<L[j];

**return** R[i]>R[j];

}); top=1;

l[1]=1; r[1]=n; p[1]=0;

FOR(i,1,n){

**if** (L[id[i]]!=l[top]||r[top]!=R[id[i]]) **return** 1;

fa[id[i]]=p[top]; top--;

**if** (id[i]<R[id[i]]) {

++top,p[top]=id[i];

l[top]=id[i]+1,r[top]=R[id[i]];

} **if** (L[id[i]]<id[i]) {

++top,p[top]=id[i];

l[top]=L[id[i]],r[top]=id[i]-1;

}

}

*// FOR(i,1,n) printf("%d ",id[i]);puts("");*

*// FOR(i,1,n) printf("%d ",fa[i]);puts("");*

**return** 0;*//okay*

}

LL inv[1000002];*//inverse*

LL fac[1000002];*//Factorial*

LL C(int n,int m){

**return** fac[n]\*inv[m]%M\*inv[n-m]%M;

}

int sz[maxn],s[maxn];

int main() {

int \_t=0;

int i;

fac[0]=1;

FOR(i,1,1000000) fac[i]=i\*fac[i-1]%M;

inv[0]=inv[1]=1;

FOR(i,2,1000000) inv[i]=(M-M/i)\*inv[M%i]%M;

FOR(i,1,1000000) inv[i]=inv[i]\*inv[i-1]%M;*// inv(n!)*

**while** (1){

read(n);

**if** (Istream::IOerror) **break**;

int i;

FOR(i,1,n) read(L[i]);

FOR(i,1,n) read(R[i]);

int ans=1;

**if** (buildtree(n)) ans=0;

FOR(i,1,n) sz[i]=1;

rFOR(i,1,n) sz[fa[id[i]]]+=sz[id[i]];

FOR(i,1,n) s[i]=sz[i]-1;

rFOR(i,2,n) {

mul\_(ans,C(s[fa[id[i]]],sz[id[i]])),s[fa[id[i]]]-=sz[id[i]];

}

printf("Case #%d: %d**\n**",++\_t,ans);

}

}

**2018hdu多校1 给定数字:**

*// 按照A从大到小建笛卡尔树*

int A[maxn],fa[maxn],id[maxn];*//id: topo*

void buildtree(int n){

**static** int S[maxn],top,tot,i;

tot=top=0;

FOR(i,1,n){

int now=0;

**while** (top&&A[S[top]]<A[i]){

**if** (now) fa[now]=S[top],id[++tot]=now;*//pop*

now=S[top]; top--;

} S[++top]=i;

**if** (now) fa[now]=S[top],id[++tot]=now;*//pop*

} int now=0;

**while** (top){

**if** (now) fa[now]=S[top],id[++tot]=now;

now=S[top]; top--;

} fa[now]=0; id[++tot]=now;

reverse(id+1,id+1+n);*// 变成正的*

}

int inv[maxn];

int sz[maxn];*//求树的size*

int main() {

int T,\_t;

int i;

FOR(i,1,1000000) inv[i]=powMM((ll)i,M-2);

scanf("%d",&T);

FOR(\_t,1,T){

scanf("%d",&n);

FOR(i,1,n) scanf("%d",&A[i]);

buildtree(n);

int ans=(ll)n\*inv[2]%M;

FOR(i,1,n) sz[i]=1;

rFOR(i,2,n) sz[fa[id[i]]]+=sz[id[i]];

FOR(i,1,n) mul\_(ans,inv[sz[i]]);

printf("%d**\n**",ans);

}

}

强连通分量tarjan

**struct** Edge {

int to,next;

Edge(int \_to=0,int \_next=-1):to(\_to),next(\_next) {};

} edge[maxn\*2];

int head[maxn],etot;

**inline** void addedge(int u,int v) {

edge[++etot]=Edge(v,head[u]);

head[u]=etot;

}

**//lowlink是说,遇到的min**

**//无向图:**

**//u割点:low[v]>=dfn[u];(表示能到的点都在之后)**

**//u-v割边(桥):low[v]>dfn[u];(要在u-v处得到)**

**//块:low[u]==dfn[u];(最终从stack取出x)**

**//dfs时注意fa和重边处理**

**//无向图不用vis这个东西=\_=,vis是为了避免横叉边**

vector<int> nodes[maxn];

int cnt;

int dfn[maxn],low[maxn],tot;

bool vis[maxn];*//instack*

int S[maxn],top;

int id[maxn];

void tarjan(int x,int fa) {

low[x]=dfn[x]=++tot;

S[++top]=x;

vis[x]=1;

**for**(int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if**(v==fa) **continue**;

**if**(!dfn[v]) {

tarjan(v,x);

low[x]=min(low[x],low[v]);

} **else** **if**(vis[v])

low[x]=min(low[x],dfn[v]);

}

**if**(low[x]==dfn[x]) {

cnt++;

**while**(1) {

int now=S[top--];

vis[now]=0;

id[now]=cnt;

nodes[cnt].push\_back(now);

**if**(now==x) **break**;

}

}

}

支配树

**//lowlink是说,遇到的min**

**//无向图:**

**//u割点:low[v]>=dfn[u];(表示能到的点都在之后)**

**//u-v割边(桥):low[v]>dfn[u];(要在u-v处得到)**

**//块:low[u]==dfn[u];(最终从stack取出x)**

**//dfs时注意fa和重边处理**

**//有向图:**

**//DAG上的割边:u-v:cnt[u]\*cnt[v]==cnt[t](mod?)**

**//DAG上的割边是固定的,也就是说求出来以后最短路是一样长的**

**//有环割边:将边变成点,然后跑支配树即可**

**//支配树:(注意,由于可能有到达不了的节点,初始化时注意答案更新)**

**//半必经点(semi=mindep{通过非树枝边fa})定理:(semi[x]=id[temp]),**

**//temp=min(temp,dfn[pre]),dfn[x]>dfn[pre](树枝边|前向边)**

**//temp=min{temp,dfn[semi[ancestor\_pre(fa)]]}**

**//dfn[x]<dfn[pre](横叉边|后向边)**

**//必经点(idom)定理:y=id[min{dfn[z]}],z:semi\_path上的点**

**//idom[x]=semi[x],semi[x]==semi[y]**

**//idom[x]=idom[y],semi[x]!=semi[y]**

**struct** Edge {

int to,next;

Edge(int \_to=0,int \_next=-1):to(\_to),next(\_next) {};

} edge[maxn\*4];

int head[maxn],pre[maxn],dom[maxn],etot; *//edges*

**inline** void addedge(int head[],int u,int v) {

edge[++etot]=Edge(v,head[u]);

head[u]=etot;

}

int dfn[maxn],tot,par[maxn]; *//dfs-tree*

int Fa[maxn],best[maxn]; *//disjoint-set*

int semi[maxn],id[maxn],idom[maxn]; *//dom-tree*

**inline** int getfa(int x) {

**if**(Fa[x]==x) **return** x;

int F=getfa(Fa[x]);

**if**(dfn[semi[best[x]]]>dfn[semi[best[Fa[x]]]])

best[x]=best[Fa[x]];

**return** Fa[x]=F;

}

void dfs(int x) {

dfn[x]=++tot;

id[tot]=x;

**for**(int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if**(!dfn[v]) par[v]=x,dfs(v);

}

}

void tarjan(int n) {

int i;

FOR(i,1,n) dom[i]=-1;

FOR(i,1,n) best[i]=semi[i]=Fa[i]=i;

rFOR(i,2,tot) {

int x=id[i];

**for**(int j=pre[x]; ~j; j=edge[j].next) {

int v=edge[j].to;

**if**(!dfn[v]) **continue**; *//could not reach*

getfa(v); *//pre\_dfn:not changed*

**if**(dfn[semi[best[v]]]<dfn[semi[x]])

semi[x]=semi[best[v]];

}

addedge(dom,semi[x],x);

Fa[x]=par[x];

x=id[i-1];

**for**(int j=dom[x]; ~j; j=edge[j].next) { *//path*

int v=edge[j].to;

getfa(v); *//id[min{dfn[z]}];*

**if**(semi[best[v]]==x) idom[v]=x;

**else** idom[v]=best[v];

}

}

FOR(i,2,tot) {

int x=id[i];

**if**(idom[x]!=semi[x]) idom[x]=idom[idom[x]];

}

}

LL n,m;

LL CNT[maxn];

LL solve() {

LL ret=(LL)tot\*(tot-1)/2;

int i;

rFOR(i,2,tot) {

int x=id[i];

CNT[x]++;

**if**(idom[x]==1) ret-=CNT[x]\*(CNT[x]-1)/2;

**else** CNT[idom[x]]+=CNT[x];

}

**return** ret;

}

int main() {

int i;

scanf("%d%d",&n,&m);

FOR(i,1,n) head[i]=pre[i]=-1;

FOR(i,1,n) dfn[i]=id[i]=idom[i]=0;etot=tot=0;

FOR(i,1,m) {

int u,v;

scanf("%d%d",&u,&v);

addedge(head,u,v);

addedge(pre,v,u);

}

dfs(1);

tarjan(n);

*// FOR(i,1,n) printf("%2d ",par[i]);puts("");*

*// FOR(i,1,n) printf("%2d ",id[i]);puts("");*

*// FOR(i,1,n) printf("%2d ",idom[i]);puts("");*

printf("%lld**\n**",solve());

}

边双连通分量 仙人掌图

**// 2018hdu多校5A**

**// 题意: 每两个点之间只能有两条路径**

**// 也就是仙人掌**

**// 求\sum flow(i,j)^i^j, i<j**

**// 做法: 有环的话, 一定会切掉环上的一个边**

**// 所以把贡献加到其他里, lca\_sum(count)**

**// 2019nowcoder6I题意: 仙人掌, 从1号点开始选**

**// 如果存在u-v且u在某轮选了, v在下一轮选择概率w**

**// 求期望概率p. 做法是p=1-p1p2...-pipi+1...**

**// 然后仙人掌dp即可**

**// 注意一下方向**

**struct** edges {

int u,v,len;

} e[maxn];

*// vector<edges> E;*

**namespace** tarjan { *// 边双连通分量, 这里是在做仙人掌*

**struct** Edge {

int to,next,id;

Edge(int \_to=0,int \_next=-1,int \_id=0):to(\_to),next(\_next),id(\_id) {};

} edge[maxn\*2];

int head[maxn],etot;

**inline** void addedge(int u,int v,int id) {

edge[++etot]=Edge(v,head[u],id); head[u]=etot;

}

int dfn[maxn],low[maxn],tot;

bool vis[maxn],used[maxn];

int S[maxn],top;

int value[maxn];*//to\_lower*

void tarjan(int x,int fa) {

low[x]=dfn[x]=++tot; vis[x]=1;

value[x]=1;

**for** (int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (used[edge[i].id]) **continue**;

**if** (v==fa) **continue**;

S[++top]=edge[i].id;

used[edge[i].id]=1;

**if** (!dfn[v]) {

tarjan(v,x);

low[x]=min(low[x],low[v]);

**if** (dfn[x]<=low[v]) { *//割边和边双联通*

vector<int> eid,pid,basp;

int nowid=x;

**while** (1) {

int id=S[top--];

eid.push\_back(id);

nowid^=e[id].u^e[id].v;

pid.push\_back(nowid);*//last*

basp.push\_back(1);*//msut; 得正反分别一遍*

debug("mapping: %d: (%d,%d)**\n**",nowid,e[id].u,e[id].v);

**if** (id==edge[i].id) **break**;

} **if** (low[v]==dfn[x]) { *//双联通, 在这里dp*

deputs(" circle :");*//环, eid; pid*

ll nowp=1;

**for** (int \_=0; \_<(int)eid.size()-1; \_++) {*//正*

int noweid=eid[\_];

edges nowe=e[noweid];

nowp=nowp\*nowe.len%M;

basp[\_]=nowp;

debug("e: (%d-%d) : %d**\n**",nowe.u,nowe.v,nowe.len);

} nowp=1;

**for** (int \_=eid.size()-2; \_>=0; \_--) {*//倒着*

int noweid=eid[\_+1];

edges nowe=e[noweid];

nowp=nowp\*nowe.len%M;

ll base=M+1-(M+1-nowp)\*(M+1-basp[\_])%M;*//\*oth[i]*

add\_(value[x],value[pid[\_]]\*base%M);

debug("e(2): (%d-%d) %d : %d**\n**",nowe.u,nowe.v,pid[\_],nowe.len);

}

} **else** {

deputs(" tree :");

**for** (int \_=0; \_<(int)eid.size(); \_++) {

int noweid=eid[\_];

edges nowe=e[noweid];

add\_(value[x],value[pid[\_]]\*nowe.len%M);

debug("e: (%d-%d) %d : %d**\n**",nowe.u,nowe.v,pid[\_],nowe.len);

*// E.push\_back(e[now]);//割边*

}

}

debug("%d %d**\n**",low[v],dfn[x]);

}

} **else** **if** (vis[v])

low[x]=min(low[x],dfn[v]);

}

}

void init(int n,int m) {

int i;

FOR(i,1,m) used[i]=0;

FOR(i,1,n) head[i]=-1,dfn[i]=0; etot=tot=0;

FOR(i,1,m) addedge(e[i].u,e[i].v,i),addedge(e[i].v,e[i].u,i);

*// FOR(i,1,n) if (!dfn[i]) tarjan(i,0);*

}

}

int main() {

int T,\_; T=1;

scanf("%d",&T);

FOR(\_,1,T) {

int n,m,i;

scanf("%d%d",&n,&m);

FOR(i,1,m) {

int a,b;

scanf("%d%d%d%d",&e[i].u,&e[i].v,&a,&b);

e[i].len=a\*powMM(b,M-2)%M;

} tarjan::init(n,m);

tarjan::tarjan(1,0);

ll Ans=tarjan::value[1];

printf("Case #%d: %lld**\n**",\_,Ans);

}

}

环套外向树

**// wannafly挑战赛16E**

**// 题意: 给个基环内向树, 每个点每时刻走1**

**// 问你最后某时刻 某个pos有几个点**

**// 做法是基环内向树dp一下, 分两部分贡献算一下**

**struct** node {

int l,r,val;

} T[maxn\*20]; int ntot;

void ins(int &x,int pos,int L,int R) {

**if** (!x) x=++ntot; T[x].val++;

**if** (L==R) **return**;

int mid=(L+R)/2;

**if** (pos<=mid) ins(T[x].l,pos,L,mid);

**else** ins(T[x].r,pos,mid+1,R);

}

int que(int x,int l,int r,int L,int R) {

**if** (!x) **return** 0;

**if** (l<=L&&R<=r) **return** T[x].val;

int ret=0,mid=(L+R)/2;

**if** (l<=mid) ret+=que(T[x].l,l,r,L,mid);

**if** (mid<r) ret+=que(T[x].r,l,r,mid+1,R);

**return** ret;

}

int A[maxn];

vector<int> cir,edge[maxn];

map<int,int> cirnum[maxn];

int vis[maxn],cfa[maxn],circnt[maxn],dep[maxn];

int in[maxn],out[maxn],dtot,ctot;

void dfs(int x,int depth,int cir\_id) {

vis[x]=1; in[x]=++dtot; cfa[x]=cir\_id; dep[x]=depth;

**for** (int v:edge[x]) dfs(v,depth+1,cir\_id);

out[x]=dtot;

}

void solve(int x) {

cir.clear(); ctot++;

**while** (A[x]&&!vis[A[x]]) x=A[x],vis[x]=1;

**while** (A[x]&&vis[A[x]]==1) {

vis[A[x]]=2; cir.push\_back(x);

cfa[x]=ctot; x=A[x];

} int i; circnt[ctot]=cir.size();

rREP(i,cir.size()-1) dep[cir[i]]=dep[A[cir[i]]]+1;

**for** (int v:cir) **for** (int y:edge[v]) **if** (vis[y]!=2) dfs(y,1,v);

}

int n,m;

int root[maxn];

vector<pair<int,int> > t\_t[maxn];

void update(int i,int k) {

**if** (vis[k]==1) {

ins(root[i+dep[k]],in[k],1,n);

i+=dep[k]; k=cfa[k];

t\_t[i].push\_back(make\_pair(cfa[k],(i+dep[k])%circnt[cfa[k]]));

} **else** cirnum[cfa[k]][(i+dep[k])%circnt[cfa[k]]]++;

}

int getans(int i,int k) {

**if** (vis[k]==1) **return** que(root[i+dep[k]],in[k],out[k],1,n);

**else** **return** cirnum[cfa[k]][(i+dep[k])%circnt[cfa[k]]];

}

int lastans;

int main() {

int i,k;

scanf("%d",&n);

FOR(i,1,n) {

scanf("%d",&A[i]);

edge[A[i]].push\_back(i);

}

FOR(i,1,n) **if** (!vis[i]) solve(i);

scanf("%d",&m);

FOR(i,1,m) {

scanf("%d",&k);

k^=lastans;

update(i,k);

**for** (**auto** now:t\_t[i]) cirnum[now.first][now.second]++;

lastans=getans(i,k);

debug(" ans = ");

printf("%d**\n**",lastans);

}

**return** 0;

}

网络流

**最大权闭合图**

**题意:给定一个有向图,每个点有权值,求最大权闭合图(与没选的没边相连),使得sigma(val)最大**

**做法:S->+node(val);-node->T(-val);原边->INF,与S相连的最小割即为所求**

**原因:简单割=>切的全是和S,T相连的边**

**假设最终与S相连的点正的x1,负的y1;T的正的x2,负的y2,(x2=S切,y1=T切)**

**最小割C=S切的正的+T切的负的=x2+y1(即反过来)**

**要求的val=x1-y1**

**C+val=x1+x2=定值,val=x1+x2-C**

**C最小,即最大流**

**最大密度子图**

**边数/点数最大**

**这个是转化成权闭合图的做法：**

**二分答案**

**将边看成点**

**S->边,1**

**边->连着的两点,1**

**每个点->T,val**

**求完即可**

**因为 边-k\*点>=0,二分出这个即可得到答案**

**做法二：**

**s->顶点，权值m**

**顶点之间连边，权值1**

**顶点->T，m+2\*ans-d[i](度数)**

**满流就OK**

**最小割的点可以放到边上，然后考虑边！**

**做法：奇偶染色，拆点然后最小割**

**最小割填INF边的意义：使得一个矩形不可行！**

**最小路径覆盖:**

**将原图拆点成两半, 然后连成二分图(边就分开来)**

**然后求个最大匹配(当然跑网络流也行)**

**要求的路径就是, 最大匹配走的路径**

**答案是边数减去匹配的边数**

**这里输出方案有个trick, 拓扑排序感觉写起来最舒服**

**//DINIC+当前弧优化**

**namespace** maxflow {

**typedef** int type;

**const** type INF=0x3f3f3f3f;

**struct** node {

int to; type cap; int next;

node(int t=0,type c=0,int n=0):to(t),cap(c),next(n) {};

} edge[maxn\*50];

int head[maxn],tot;

void addedge(int from,int to,type cap,type rcap=0) {

edge[tot]=node(to,cap,head[from]); head[from]=tot++;

edge[tot]=node(from,rcap,head[to]); head[to]=tot++;

}

int dep[maxn],cur[maxn];*//当前弧优化*

bool bfs(int s,int t,int n) {

**static** int Q[maxn],ST,ED;

memset(dep+1,0,n\***sizeof**(int));

ST=0; ED=-1;

Q[++ED]=s; dep[s]=1;

**while** (ST<=ED) {

int u=Q[ST++];

**for** (int i=head[u]; i!=-1; i=edge[i].next) {

int v=edge[i].to;

**if** (!dep[v]&&edge[i].cap) {

Q[++ED]=v; dep[v]=dep[u]+1;

}

}

} **return** (dep[t]!=0);

}

type dfs(int x,**const** int &t,type flow=INF) {

**if** (x==t||flow==0) **return** flow;

type ret=0;

**for** (int i=cur[x]; i!=-1; i=edge[i].next) {

**if** (dep[x]+1==dep[edge[i].to]&&edge[i].cap){

type f=dfs(edge[i].to,t,min(flow,edge[i].cap));

edge[i].cap-=f; edge[i^1].cap+=f;

ret+=f; flow-=f; cur[x]=i;

**if** (flow==0) **break**;

}

} **if** (!ret) dep[x]=0;

**return** ret;

}

type maxflow(int s,int t,int n) {

type ret=0;

**while** (bfs(s,t,n)) {

type f;

memcpy(cur+1,head+1,n\***sizeof**(int));

**while** ((f=dfs(s,t))>0) ret+=f;

} **return** ret;

}

void init(int n) {

memset(head+1,0xff,n\***sizeof**(int)); tot=0;

}

}

**//ISAP**

**namespace** maxflow {

**typedef** LL type;

**const** type INF=0x3f3f3f3f3f3f3f3f;

**struct** node {

int to; type cap; int next;

node(int t=0,type c=0,int n=0):to(t),cap(c),next(n) {};

} edge[maxn\*50];

int head[maxn],tot;

void addedge(int from,int to,type cap,type rcap=0) {

edge[tot]=node(to,cap,head[from]); head[from]=tot++;

edge[tot]=node(from,rcap,head[to]); head[to]=tot++;

}

int gap[maxn],dep[maxn],cur[maxn];

void bfs(int s,int t,int n) {*//t好像没啥用啊=\_=*

**static** int Q[maxn],ST,ED;

memset(dep+1,0xff,n\***sizeof**(int));

memset(gap+1,0,n\***sizeof**(int));

gap[0]=1; dep[t]=0;

ST=0; ED=-1; Q[++ED]=t;

**while** (ST<=ED) {

int u=Q[ST++];

**for** (int i=head[u]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (dep[v]!=-1) **continue**;

Q[++ED]=v; dep[v]=dep[u]+1;

gap[dep[v]]++;

}

}

}

int S[maxn];

type sap(int s,int t,int n) {

bfs(s,t,n);

memcpy(cur+1,head+1,n\***sizeof**(int));

int top=0,u=s; type ret=0;

**while** (dep[s]<n) {

**if** (u==t) {

type MIN=INF,inser=0,i;

REP(i,top) **if** (MIN>edge[S[i]].cap)

MIN=edge[S[i]].cap,inser=i;

REP(i,top) {

edge[S[i]].cap-=MIN,edge[S[i]^1].cap+=MIN;

} ret+=MIN; top=inser; u=edge[S[top]^1].to;

**continue**;

} bool flag=0; int v;

**for** (int i=cur[u]; ~i; i=edge[i].next) {

v=edge[i].to;

**if** (edge[i].cap&&dep[v]+1==dep[u]) {

flag=1; cur[u]=i; **break**;

}

} **if** (flag) {

S[top++]=cur[u]; u=v; **continue**;

} int MIN=n;

**for** (int i=head[u]; ~i; i=edge[i].next) {

v=edge[i].to;

**if** (edge[i].cap&&dep[v]<MIN)

MIN=min(MIN,dep[v]),cur[u]=i;

} gap[dep[u]]--;

**if** (ret>INF) **return** ret;*//not okay*

**if** (!gap[dep[u]]) **return** ret;

dep[u]=MIN+1; gap[dep[u]]++;

**if** (u!=s) u=edge[S[--top]^1].to;

} **return** ret;

}

void init(int n) {

memset(head+1,0xff,n\***sizeof**(int)); tot=0;

}

}

无向图全局最小割

**无向图 分成两块最小割**

**做法:O(n^3)|O(nmlogm)**

**观察到最小割一定是两块中找个点的最小割**

**那么我们考虑每次找到S->T的最小割后缩点**

**随便找最小割的方法:O(n^2)|O(mlogm)**

**得到s,t的方法:先任意找个a开始**

**定义集合A:一些点的集合**

**定义w(A,v):v到A中所有点的sum\_value**

**每次从中找出w最大的点加入A**

**最后加入的两个点记为S,T**

**S->T的最大流的大小为最末的w**

**O(nmlogm)**

bool deleted[maxn],vis[maxn];

vector<pair<int,int> > edge[maxn];

priority\_queue<pair<int,int> > Q;

int weight[maxn];

int fa[maxn];

**inline** int getfa(int x){

**if** (fa[x]==x) **return** x;

**return** fa[x]=getfa(fa[x]);

}

int getst(int &s,int &t,int n){

int i;t=1;

**while** (Q.size()) Q.pop();

REP(i,n-1){

vis[s=t]=1;

**for** (**auto** &e:edge[s]) {

int v=getfa(e.second);

e.second=v;

**if** (!vis[v])

Q.push(make\_pair(weight[v]+=e.first,v));

}t=0;

**while** (!t&&Q.size()){

**auto** now=Q.top();Q.pop();

int v=now.second;

**if** (!vis[v]) t=v;

}**if** (!weight[t]) **return** 0;

}**return** weight[t];

}

int mincut(int n){

int ret=INF;

int s,t,i,j,k;

FOR(i,1,n) deleted[i]=0,fa[i]=i;

rFOR(i,2,n){

FOR(j,1,n) weight[j]=0,vis[j]=0;

ret=min(ret,getst(s,t,i));

**if** (!ret) **return** 0;

**for** (**auto** v:edge[t]) edge[s].push\_back(v);

int x=getfa(s),y=getfa(t);fa[y]=x;

vector<pair<int,int> >().swap(edge[t]);

}**return** ret;

}

**O(n^3)**

LL edge[507][507];

bool deleted[maxn],vis[maxn];

vector<int> id;

LL weight[maxn];

LL getst(int &s,int &t,int n){

int i;t=1;

**for** (int v:id) weight[v]=0,vis[v]=0;

REP(i,n-1){

vis[s=t]=1;

**for** (int v:id) **if** (!vis[v])

weight[v]+=edge[s][v],t=v;

**for** (int v:id) **if** (!vis[v])

**if** (weight[v]>=weight[t]) t=v;

**if** (!weight[t]) **return** 0;

}**return** weight[t];

}

LL mincut(int n){

LL ret=INFF;

int s,t,i,j,k;

FOR(i,1,n) deleted[i]=0;

rFOR(i,2,n){

j=0;id.clear();

FOR(k,1,n) **if** (!deleted[k]) id.push\_back(k);

ret=min(ret,getst(s,t,id.size()));

**if** (!ret) **return** 0;

**for** (int v:id) **if** (v!=s&&v!=t){

edge[s][v]+=edge[t][v];

edge[v][s]+=edge[v][t];

}deleted[t]=1;

}**return** ret;

}

无向图最小割树 GH-tree

**//两点的LCA\_MAX是最小割**

**namespace** gomoryhu\_tree {

**typedef** int type;

**struct** node { *//只能是双向的*

int u,v; type len;

node(int u=0,int v=0,type len=0):u(u),v(v),len(len) {};

} edge[maxn],e[maxn];

int tot,etot;

void addedge(int u,int v,int len) {

edge[++tot]=node(u,v,len);

} int n;

void solve(int l,int r,int id[]) {*//id,id+n*

**static** int tmp[maxn];

**if** (l==r) **return**;

random\_shuffle(id+l,id+r+1);

maxflow::init(n); int i,L=l,R=r;

FOR(i,1,tot) maxflow::addedge(edge[i].u,edge[i].v,edge[i].len,edge[i].len);

e[++etot]=node(id[l],id[r],maxflow::maxflow(id[l],id[r],n));

FOR(i,l,r) **if** (maxflow::dep[id[i]])

tmp[L++]=id[i]; **else** tmp[R--]=id[i];

FOR(i,l,r) id[i]=tmp[i];

solve(l,R,id); solve(L,r,id);

}

void init(int \_n) {

n=\_n; tot=etot=0;

srand(time(0));

}

}

最小费用流

**// 这个好像就是zkw费用流**

**// 拆点后可以S向入连边, 出向T连边, 然后入和出就可以保持动态平衡!**

**// 连边是为了将"获取的"和"使用的"联系起来! 大概意思就是, 使用的流量确定...**

**// 注意观察特殊性质**

**// 费用流有个"短路"的性质, 如果流到这里可能会使得其他的流量减少, 这个好像有点用**

**// cf 101492I**

**// 题意:每个点可以流无限,费用value**

**// 存在limit为l-r最多用x个**

**// 流量费用互换, 流量转化为差分约束即可**

**namespace** mincostflow {

**typedef** ll type;

**const** type INF=0x3f3f3f3f3f3f3f3fll;

**struct** node {

int to; type cap,cost; int rev;

node(int t=0,type c=0,type \_c=0,int n=0):

to(t),cap(c),cost(\_c),rev(n) {};

}; vector<node> edge[maxn];

void addedge(int from,int to,type cap,type cost,type rcap=0) {

edge[from].push\_back(node(to,cap,cost,edge[to].size()));

edge[to].push\_back(node(from,rcap,-cost,edge[from].size()-1));

}

type dis[maxn];

bool mark[maxn];

void spfa(int s,int t,int n) {

memset(dis+1,0x3f,n\***sizeof**(type));

memset(mark+1,0,n\***sizeof**(bool));

**static** int Q[maxn],ST,ED;

dis[s]=0; ST=ED=0; Q[ED++]=s;

**while** (ST!=ED) {

int v=Q[ST]; mark[v]=0;

**if** ((++ST)==maxn) ST=0;

**for** (node &e:edge[v]) {

**if** (e.cap>0&&dis[e.to]>dis[v]+e.cost) {

dis[e.to]=dis[v]+e.cost;

**if** (!mark[e.to]) {

**if** (ST==ED||dis[Q[ST]]<=dis[e.to]) {

Q[ED]=e.to,mark[e.to]=1;

**if** ((++ED)==maxn) ED=0;

} **else** {

**if** ((--ST)<0) ST+=maxn;

Q[ST]=e.to,mark[e.to]=1;

}

}

}

}

}

} int cur[maxn];

type dfs(int x,int t,type flow) {

**if** (x==t||!flow) **return** flow;

type ret=0; mark[x]=1;

int i;

rep(i,cur[x],(int)edge[x].size()) {

node &e=edge[x][i];

**if** (!mark[e.to]&&e.cap) {

**if** (dis[x]+e.cost==dis[e.to]) {

int f=dfs(e.to,t,min(flow,e.cap));

e.cap-=f; edge[e.to][e.rev].cap+=f;

ret+=f; flow-=f; cur[x]=i;

**if** (flow==0) **break**;

}

}

} mark[x]=0;

**return** ret;

}

pair<type,type> mincostflow(int s,int t,int n,type flow=INF) {

type ret=0,ans=0;

**while** (flow) {

spfa(s,t,n); **if** (dis[t]==INF) **break**;

*// 这样加当前弧优化会快, 我也不知道为啥*

memset(cur+1,0,n\***sizeof**(int));

type len=dis[t],f;

**while** ((f=dfs(s,t,flow))>0)*//while也行*

ret+=f,ans+=len\*f,flow-=f;

} **return** make\_pair(ret,ans);

}

void init(int n) {

int i; FOR(i,1,n) edge[i].clear();

}

}

int A[maxn];

int main() {

int n,m;

int i;

scanf("%d%d",&n,&m);

mincostflow::init(n+1+2);

int s=n+2,t=n+3;

FOR(i,1,n) {

scanf("%d",&A[i]);

mincostflow::addedge(s,i,A[i],0);

mincostflow::addedge(i+1,t,A[i],0);

mincostflow::addedge(i+1,i,INF,0);*//i-(i+1)<=0*

}

FOR(i,1,m) {

int l,r,c;

scanf("%d%d%d",&l,&r,&c); r++;

mincostflow::addedge(l,r,INF,c);*//r-l<=n*

}

printf("%lld**\n**",mincostflow::mincostflow(s,t,n+3,INF).second);

}

**//原始对偶dij,可以跑负边**

**namespace** mincostflow {

**typedef** int type;

**const** type INF=0x3f3f3f3f;

**struct** node {

int to; type cap,cost; int rev;

node(int t=0,type c=0,type \_c=0,int n=0):

to(t),cap(c),cost(\_c),rev(n) {};

}; vector<node> edge[maxn];

void addedge(int from,int to,type cap,type cost,type rcap=0) {

edge[from].push\_back(node(to,cap,cost,edge[to].size()));

edge[to].push\_back(node(from,rcap,-cost,edge[from].size()-1));

}

int prev[maxn],pree[maxn];*//pre\_cnt*

type dis[maxn],h[maxn];

pair<type,type> mincostflow(int s,int t,int n,type flow=INF) {

type ret=0,ans=0;

memset(h+1,0,n\***sizeof**(type));

**while** (flow) {

*// dij*

**typedef** pair<type,int> pti;

memset(dis+1,0x3f,n\***sizeof**(type));

**static** priority\_queue<pti,vector<pti>,greater<pti> > Q;

dis[s]=0; Q.push(pti(0,s));

**while** (Q.size()) {

**auto** now=Q.top(); Q.pop();

**if** (dis[now.second]<now.first) **continue**;

int i,v=now.second;

REP(i,(int)edge[v].size()) {

node &e=edge[v][i];

**if** (e.cap>0&&dis[e.to]>dis[v]+e.cost+h[v]-h[e.to]) {

dis[e.to]=dis[v]+e.cost+h[v]-h[e.to];

prev[e.to]=v; pree[e.to]=i;

Q.push(pti(dis[e.to],e.to));

}

}

} int i;

**if** (dis[t]==INF) **break**;

FOR(i,1,n) h[i]+=dis[i];

type d=flow;

**for** (int i=t; i!=s; i=prev[i])

d=min(d,edge[prev[i]][pree[i]].cap);

**if** (d==0) **break**;

flow-=d; ret+=d; ans+=d\*h[t];

**for** (int i=t; i!=s; i=prev[i]){

node &e=edge[prev[i]][pree[i]];

e.cap-=d; edge[e.to][e.rev].cap+=d;

}

} **return** make\_pair(ret,ans);

}

void init(int n) {

int i; FOR(i,1,n) edge[i].clear();

}

}

上下界网络流

**//可二分t->s边的下/上界,即可达到最大最小流**

**//最大流:t->s连边, ss->tt流, s->t正向最大流,会流掉反向建的边的流量, 流量就是ans**

**//最小流:ss->tt流, t->s连边, ss->tt流, s->t的就是最大流**

**//带权值的直接加权即可, in和out加的边val=0(只是为了限制流出可以等于流入而已)**

**namespace** pipeflow {

**typedef** int type;

int eid[maxn\*10],etot;

type in[maxn],out[maxn],flow[maxn\*10];

int s\_s,t\_t;*//S,T*

int addedge(int u,int v,int low,int high) {

eid[etot]=maxflow::addedge(u,v,high-low);

out[u]+=low; in[v]+=low; flow[etot++]=low;

**return** etot-1;

}

void init(int n) {

s\_s=n+1,t\_t=n+2; etot=0;

memset(in+1,0,n\***sizeof**(type));

memset(out+1,0,n\***sizeof**(type));

maxflow::init(n+2);

}

type solve(int n,int s,int t) {

int sum=0; int i;

FOR(i,1,n) {

sum+=max(0,in[i]-out[i]);

**if** (in[i]>out[i]) maxflow::addedge(s\_s,i,in[i]-out[i]);

**if** (in[i]<out[i]) maxflow::addedge(i,t\_t,out[i]-in[i]);

}

*// // maxflow:*

*// maxflow::addedge(t,s,INF);*

*// if (maxflow::maxflow(s\_s,t\_t,n+2)!=sum) return -1;*

*// return maxflow::maxflow(s,t,n+2);//maxflow*

*// // minflow:*

*// type first=maxflow::maxflow(s\_s,t\_t,n+2);*

*// int retpos=maxflow::addedge(t,s,INF);*

*// if (first+maxflow::maxflow(s\_s,t\_t,n+2)!=sum) return -1;*

*// return maxflow::edge[retpos^1].cap;//minflow*

*// okay flow:*

*// if (maxflow::maxflow(s\_s,t\_t,n+2)!=sum) return 0;*

REP(i,etot) flow[i]+=maxflow::edge[eid[i]^1].cap;*//edges*

*//return 1;*

}

}

树分治

**//HDU6268**

**//ccpc2017杭州**

**//树分治后 树形依赖DP**

int A[maxn];

vector<int> edge[maxn];

int sz[maxn];

bool mark[maxn];

int minweight,root;

void dfs1(int x,int fa,int n) {

int weight=0; sz[x]=1;

**for** (int v:edge[x]) {

**if** (v==fa||mark[v]) **continue**;

dfs1(v,x,n); sz[x]+=sz[v];

weight=max(weight,sz[v]);

} weight=max(weight,n-sz[x]);

**if** (weight<minweight) root=x,minweight=weight;

}

bitset<100007> now[3007],ans;*//depth*

void dfs2(int x,int fa,int dep) {

now[dep]=now[dep-1]; sz[x]=1;

**for** (int v:edge[x]) {

**if** (v==fa||mark[v]) **continue**;

dfs2(v,x,dep+1); sz[x]+=sz[v];

} now[dep-1]|=now[dep]<<A[x];

}

void calc(int x){

now[0].reset(); now[0].set(0);

dfs2(x,0,1); ans|=now[0];

}

void dfs3(int x) {

calc(x); mark[x]=1;

**for** (int v:edge[x]) {

**if** (mark[v]) **continue**;

minweight=sz[v];

dfs1(v,0,sz[v]);

dfs3(root);

}

}

int main() {

int n,m,T;

int i;

scanf("%d",&T);

**while** (T--) {

scanf("%d%d",&n,&m);

REP(i,n-1) {

int u,v;

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

} FOR(i,1,n) scanf("%d",&A[i]);

minweight=n;

dfs1(1,0,n); dfs3(root);

FOR(i,1,m) printf("%d",(int)ans[i]);

puts("");

ans.reset();

FOR(i,1,n) edge[i].clear(),mark[i]=0;

}

**return** 0;

}

动态点分治

**//题意: 动态查询到某点距离不超过x的权值和, 更改某点权值**

**//注意容斥的时候的length位置不是root~是上个root相连的位置**

**//也就是说dis得单独计算//dfs2一次比两次少一半多的常数=\_=**

**//addnode中是ids**

int BIT\_pool[maxn\*40],\*BIT[maxn],\*SUBBIT[maxn],\*st=BIT\_pool;

int size[maxn]; bool mark[maxn];

int minweight,root;

**struct** Node {

int to,next;

Node(int \_to=0,int \_next=0):to(\_to),next(\_next) {};

} edge[maxn\*2];

int head[maxn],tot;

void addedge(int u,int v) {

edge[++tot]=Node(v,head[u]); head[u]=tot;

}

void dfs1(int x,int fa,int n) {

int weight=0; size[x]=1;

**for** (int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (v==fa||mark[v]) **continue**;

dfs1(v,x,n);

size[x]+=size[v];

weight=max(weight,size[v]);

} weight=max(weight,n-size[x]);

**if** (weight<minweight) {root=x; minweight=weight;}

}

int length[maxn];

**struct** node {

int top,sub,len,next;

node() {}

node(int \_top,int \_sub,int \_len,int \_next):top(\_top),sub(\_sub),len(\_len),next(\_next) {};

} nodes[maxn\*20];

int calhead[maxn],caltot;

int maxdep;

void addnode(int x,int top,int sub,int len) {

nodes[++caltot]=node(top,sub,len,calhead[x]); calhead[x]=caltot;

}

void dfs2(int x,int fa,int top,int sub,int dep) {

addnode(x,top,sub,dep);

**for** (int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (v==fa||mark[v]) **continue**;

dfs2(v,x,top,sub,dep+1);

} maxdep=max(maxdep,dep);

}

int len[maxn],sublen[maxn];

void dfs3(int x) {

mark[x]=1; root=x;

maxdep=0; int xdep=0;

addnode(x,x,0,0);

**for** (int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (mark[v]) **continue**;

minweight=size[v]; dfs1(v,0,size[v]);

maxdep=0; dfs2(v,0,x,root,1); *//判重是x,init\_dep=1*

sublen[root]=maxdep; xdep=max(xdep,maxdep);

SUBBIT[root]=st; st+=sublen[root]+1;

dfs3(root);

} len[x]=xdep;

BIT[x]=st; st+=len[x]+1;

}

**inline** int lowbit(int x) {**return** x&-x;}

void add(int \*T,int n,int x,int val) {

x++; T--; n++;

**for** (; x<=n; x+=lowbit(x)) T[x]+=val;

} int get(int \*T,int x) {

x++; T--; int ret=0;

**for** (; x; x-=lowbit(x)) ret+=T[x];

**return** ret;

}

void update(int x,int val) {

**for** (int i=calhead[x]; ~i; i=nodes[i].next) {

int v=nodes[i].top,length=nodes[i].len;

add(BIT[v],len[v],length,val);

v=nodes[i].sub;

**if** (v) add(SUBBIT[v],sublen[v],length,val);

}

} int query(int x,int dis) {

int ret=0;

**for** (int i=calhead[x]; ~i; i=nodes[i].next) {

int v=nodes[i].top,length=nodes[i].len;

**if** (dis>=length) {

ret+=get(BIT[v],min(dis-length,len[v]));

v=nodes[i].sub;

**if** (v) ret-=get(SUBBIT[v],min(dis-length,sublen[v]));;

}

} **return** ret;

}

int n,m,T;

int i,j,k;

char op[2];

int a[maxn];

int main() {

**while** (~scanf("%d%d",&n,&m)) {

FOR(i,1,n) mark[i]=0,BIT[i]=SUBBIT[i]=**nullptr**;

memset(BIT\_pool,0,**sizeof**(int)\*(st-BIT\_pool)); st=BIT\_pool;

FOR(i,1,n) head[i]=calhead[i]=-1; tot=caltot=0;

FOR(i,1,n) scanf("%d",&a[i]);

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

addedge(u,v); addedge(v,u);

}

minweight=INF; dfs1(1,0,n);

dfs3(root);

FOR(i,1,n) update(i,a[i]);

FOR(i,1,m) {

int u,v;

scanf("%s%d%d",op,&u,&v);

**if** (op[0]=='!') update(u,v-a[u]),a[u]=v;

**else** printf("%d**\n**",query(u,v));

}

}

}

部分树上dp

**从求含某条边的最小生成树截下来的代码(当然前面sort了)合并(要记得merge咋写),先sort然后从小到大讨论,连father,之后merge**

**inline** int Union(int u,int v,int len) {

int ret=0;

**while** (u!=v&&(fa[u]!=u||fa[v]!=v)) {

**if** (fa[u]==u||fa[v]!=v&&sz[u]>sz[v])

{ret=max(ret,val[v]); v=fa[v];}

**else** {ret=max(ret,val[u]); u=fa[u];}

} **if** (u==v) **return** ret;

**if** (sz[u]>sz[v]) swap(u,v);

fa[u]=v; val[u]=len;

sz[v]+=sz[u]; ans=ans+len;

**return** len;

}

2-sat

**//重点是维护拆点后各种限制之间的关系，这个是个二分以后2-sat的**

**struct** T\_SAT {

**struct** enode {

int to,next;

enode(int \_to=0,int \_next=-1):to(\_to),next(\_next) {};

} edge[maxn\*maxn\*2];

int head[maxn\*2],etot;

void addedge(int u,int v) {

edge[++etot]=enode(v,head[u]);

head[u]=etot;

}

int dfn[maxn\*2],low[maxn\*2],belong[maxn\*2];

bool vis[maxn\*2];

int tot,cnt;

int S[maxn\*2],top;

void dfs(int x) {

dfn[x]=low[x]=++tot;

S[++top]=x;

vis[x]=1;

**for** (int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (!dfn[v]) {

dfs(v);

low[x]=min(low[x],low[v]);

} **else** **if** (vis[v])

low[x]=min(low[x],dfn[v]);

}

**if** (dfn[x]==low[x]) {

cnt++;

**while** (1) {

int now=S[top--];

vis[now]=0;

belong[now]=cnt;

**if** (now==x) **break**;

}

}

}

void init(int n) {

int i;

REP(i,2\*n) head[i]=-1;

etot=0;

}

bool solve(int n) {

int i;

tot=cnt=0;

REP(i,2\*n) dfn[i]=vis[i]=0;

REP(i,2\*n) **if** (!dfn[i]) dfs(i);

REP(i,n) **if** (belong[i]==belong[i+n]) **return** 0;

**return** 1;

}

} two\_sat;

int n,m;

int i,j;

int a1,a2,c1,c2;

int main() {

**while** (~scanf("%d%d",&n,&m)) {

two\_sat.init(n);

REP(i,m) {

scanf("%d%d%d%d",&a1,&a2,&c1,&c2);

**if** (c1==1&&c2==1) {

two\_sat.addedge(a1+n,a2);

two\_sat.addedge(a2+n,a1);

} **else** **if** (c1==0&&c2==1) {

two\_sat.addedge(a1,a2);

two\_sat.addedge(a2+n,a1+n);

} **else** **if** (c1==1&&c2==0) {

two\_sat.addedge(a1+n,a2+n);

two\_sat.addedge(a2,a1);

} **else** **if** (c1==0&&c2==0) {

two\_sat.addedge(a1,a2+n);

two\_sat.addedge(a2,a1+n);

}

}

**if** (two\_sat.solve(n)) puts("YES");

**else** puts("NO");

}

}

可持久化的2-sat 输出方案

**//对于一般点是对称的题目, 直接belong[i]<belong[i+n]输出即可**

**//否则需要拓扑排序, 破坏了本身良好的性质**

**// 题意: 给颗树, 每次给俩路径**

**// 问你m组询问, 从每个里选个路径, 是否可以不相交**

**// 做法: 可持久化建线段树然后 2-sat**

**// 输出方案需要把每个块都拓扑排序**

**namespace** T\_SAT {

**const** **static** int maxn=5e6+7;

**struct** enode {

int to,next;

enode(int \_to=0,int \_next=-1):to(\_to),next(\_next) {};

} edge[maxn\*6];

int head[maxn],etot;

void addedge(int u,int v) {

edge[++etot]=enode(v,head[u]); head[u]=etot;

}

int dfn[maxn],low[maxn],belong[maxn];

bool vis[maxn];

int tot,cnt;

int S[maxn],top;

void dfs(int x) {

dfn[x]=low[x]=++tot;

S[++top]=x; vis[x]=1;

**for** (int i=head[x]; ~i; i=edge[i].next) {

int v=edge[i].to;

**if** (!dfn[v]) {

dfs(v);

low[x]=min(low[x],low[v]);

} **else** **if** (vis[v])

low[x]=min(low[x],dfn[v]);

}

**if** (dfn[x]==low[x]) {

cnt++;

**while** (1) {

int now=S[top--];

vis[now]=0; belong[now]=cnt;

**if** (now==x) **break**;

}

}

}

void init() {

memset(head,-1,**sizeof**(head)); etot=0;

}

void solve(int n) {

int i; tot=cnt=0;

FOR(i,1,n) dfn[i]=vis[i]=0;

FOR(i,1,n) **if** (!dfn[i]) dfs(i);

}

}

int choose,remain;

int upid[maxn\*8],downid[maxn\*8],tot;

void build(int x,int L,int R) {

upid[x]=++tot; downid[x]=++tot;

**if** (downid[x>>1]) {

T\_SAT::addedge(downid[x>>1],downid[x]);

} **if** (L==R) **return**;

int mid=(L+R)/2;

build(x<<1,L,mid);

build(x<<1|1,mid+1,R);

}

bool update;

void query(int x,int l,int r,int L,int R) {

**if** (l>r) **return**;

**if** (l<=L&&R<=r) {

**if** (!update) {

T\_SAT::addedge(choose,downid[x]);

} **else** {

T\_SAT::addedge(++tot,downid[x]); downid[x]=tot;

T\_SAT::addedge(upid[x],remain);

T\_SAT::addedge(downid[x],remain);

int fa=downid[x>>1],ls=downid[x<<1],rs=downid[x<<1|1];

**if** (fa) T\_SAT::addedge(fa,downid[x]);

**if** (ls) T\_SAT::addedge(downid[x],ls);

**if** (rs) T\_SAT::addedge(downid[x],rs);

}

**return**;

} **else** **if** (!update) T\_SAT::addedge(choose,upid[x]);

int mid=(L+R)/2;

**if** (l<=mid) query(x<<1,l,r,L,mid);

**if** (mid<r) query(x<<1|1,l,r,mid+1,R);

}

**namespace** PRE\_CAL {

vector<int> edge[maxn];

int fa[maxn],son[maxn],id[maxn],tot;

int sz[maxn],top[maxn],dep[maxn];

void dfs\_1(int u,int father,int depth) {

fa[u]=father; dep[u]=depth;

int mx=-1; sz[u]=1; son[u]=0;

**for** (int v:edge[u]) {

**if** (father==v) **continue**;

dfs\_1(v,u,depth+1);

sz[u]+=sz[v];

**if** (sz[v]>mx) mx=sz[v],son[u]=v;

}

}

void dfs\_2(int u,int x) {

id[u]=++tot; top[u]=x;

**if** (son[u]) dfs\_2(son[u],x);

**for** (int v:edge[u]) {

**if** (v==fa[u]||v==son[u]) **continue**;;

dfs\_2(v,v);

}

}

void solve(int x,int y) {

**while** (top[x]!=top[y]) {

**if** (dep[top[x]]<dep[top[y]]) swap(x,y);

query(1,id[top[x]],id[x],1,n); x=fa[top[x]];

} **if** (dep[x]>dep[y]) swap(x,y);

**if** (son[x]) query(1,id[son[x]],id[y],1,n);

}

}

int chosen[maxn];

int A[maxn],B[maxn],C[maxn],D[maxn];

int TaskA() {

int i,j,m;

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

PRE\_CAL::edge[u].push\_back(v);

PRE\_CAL::edge[v].push\_back(u);

} scanf("%d",&m);

T\_SAT::init();

PRE\_CAL::dfs\_1(1,0,0);

PRE\_CAL::dfs\_2(1,1);

FOR(i,1,m) chosen[i]=++tot,++tot;

build(1,1,n);

FOR(i,1,m) scanf("%d%d%d%d",&A[i],&B[i],&C[i],&D[i]);

FOR(i,1,m) {

choose=chosen[i]; remain=chosen[i]+1;

update=0;

PRE\_CAL::solve(A[i],B[i]);

swap(choose,remain);

PRE\_CAL::solve(C[i],D[i]);

update=1; swap(choose,remain);

PRE\_CAL::solve(A[i],B[i]);

swap(choose,remain);

PRE\_CAL::solve(C[i],D[i]);

} build(1,1,n);

rFOR(i,1,m) {

choose=chosen[i]; remain=chosen[i]+1;

update=0;

PRE\_CAL::solve(A[i],B[i]);

swap(choose,remain);

PRE\_CAL::solve(C[i],D[i]);

update=1; swap(choose,remain);

PRE\_CAL::solve(A[i],B[i]);

swap(choose,remain);

PRE\_CAL::solve(C[i],D[i]);

}

T\_SAT::solve(tot);

FOR(i,1,m) **if** (T\_SAT::belong[chosen[i]]==T\_SAT::belong[chosen[i]+1]) **return** 0\*puts("NO");

puts("YES");

FOR(i,1,m) printf("%d**\n**",((T\_SAT::belong[chosen[i]]<T\_SAT::belong[chosen[i]+1])^1)+1);

**return** 0;

}

dfs序\_换根的讨论

**//http://codeforces.com/contest/916/problem/E**

**//改根,子树加,查,令人窒息的讨论**

**//有套路是dfs同时更新value啥的, 需要注意**

LL sum[maxn<<2],lazy[maxn<<2];

void update(int x,int l,int r,LL val,int L,int R) {

**if** (l>r) **return**;

**if** (l<=L&&R<=r) {lazy[x]+=val; sum[x]+=(R-L+1)\*val; **return**;}

int mid=(L+R)/2;

**if** (lazy[x]) {

lazy[x<<1]+=lazy[x];

lazy[x<<1|1]+=lazy[x];

sum[x<<1]+=(mid-L+1)\*lazy[x];

sum[x<<1|1]+=(R-mid)\*lazy[x];

lazy[x]=0;

} **if** (l<=mid) update(x<<1,l,r,val,L,mid);

**if** (mid<r) update(x<<1|1,l,r,val,mid+1,R);

sum[x]=sum[x<<1]+sum[x<<1|1];

}

LL query(int x,int l,int r,int L,int R) {

LL ret=0;

**if** (l>r) **return** 0;

**if** (l<=L&&R<=r) **return** sum[x];

int mid=(L+R)/2;

**if** (lazy[x]) {

lazy[x<<1]+=lazy[x];

lazy[x<<1|1]+=lazy[x];

sum[x<<1]+=(mid-L+1)\*lazy[x];

sum[x<<1|1]+=(R-mid)\*lazy[x];

lazy[x]=0;

} **if** (l<=mid) ret+=query(x<<1,l,r,L,mid);

**if** (mid<r) ret+=query(x<<1|1,l,r,mid+1,R);

sum[x]=sum[x<<1]+sum[x<<1|1];

**return** ret;

}

vector<int> edge[maxn];

int fa[maxn][27];

int in[maxn],out[maxn],tot,dep[maxn];

void dfs(int x,int f,int d) {

int i;

fa[x][0]=f; in[x]=++tot; dep[x]=d;

rep(i,1,20) fa[x][i]=fa[fa[x][i-1]][i-1];

**for** (int v:edge[x]) **if** (v!=f) dfs(v,x,d+1);

out[x]=tot;

}

int lca(int x,int y) {

int i;

**if** (dep[x]<dep[y]) swap(x,y);

rREP(i,20) **if** (dep[x]-dep[y]>=1<<i) x=fa[x][i];

**if** (x==y) **return** x;

rREP(i,20) **if** (fa[x][i]!=fa[y][i]) x=fa[x][i],y=fa[y][i];

**return** fa[x][0];

}

int getnthfa(int x,int k) {

int i;

rREP(i,20) **if** ((k>>i)&1) x=fa[x][i];

**return** x;

}

int root;

int n,m;

int a[maxn];

int main() {

int i,j;

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%d",&a[i]);

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

}

dfs(1,0,0);

FOR(i,1,n) update(1,in[i],in[i],a[i],1,n); root=1;

**while** (m--) {

int op,u,v,x;

scanf("%d",&op);

**if** (op==1) {

scanf("%d",&root);

} **else** **if** (op==2) {

scanf("%d%d%d",&u,&v,&x);

int f=lca(u,v)^lca(v,root)^lca(u,root);

**if** (f==root) update(1,1,n,x,1,n);

**else** **if** (lca(f,root)==f) {

int t=getnthfa(root,dep[root]-dep[f]-1);

update(1,1,in[t]-1,x,1,n);

update(1,out[t]+1,n,x,1,n);

} **else** update(1,in[f],out[f],x,1,n);

} **else** **if** (op==3) {

int x;

LL ans;

scanf("%d",&x);

**if** (x==root) ans=query(1,1,n,1,n);

**else** **if** (in[x]<=in[root]&&in[root]<=out[x]) {

int t=getnthfa(root,dep[root]-dep[x]-1);

ans=query(1,1,in[t]-1,1,n)+query(1,out[t]+1,n,1,n);

} **else** ans=query(1,in[x],out[x],1,n);

printf("%I64d**\n**",ans);

}

}

}

01序列转树

FOR(i,1,n) {

scanf("%d",&b[i]);

**if** (b[i]) {*//up*

pid[i]=now;

**if** (!fa[now]) {

root=++tot\_id; fa[now]=tot\_id;

edge[tot\_id].push\_back(now);

} now=fa[now];

} **else** {*//down*

fa[++tot\_id]=now;

edge[now].push\_back(tot\_id);

now=tot\_id; pid[i]=now;

}

}

有方向的树链剖分

int a[maxn],tot;

int mxr[maxn<<2],mxl[maxn<<2];

int mx[maxn<<2],mn[maxn<<2];

int lazy[maxn<<2];*//profit=mx-mn*

void change(int x,int val) {

lazy[x]+=val;

mx[x]+=val; mn[x]+=val;

}

void pushup(int x) {

mxr[x]=max(max(mxr[x<<1],mxr[x<<1|1]),mx[x<<1|1]-mn[x<<1]);*//->*

mxl[x]=max(max(mxl[x<<1],mxl[x<<1|1]),mx[x<<1]-mn[x<<1|1]);*//<-*

mx[x]=max(mx[x<<1],mx[x<<1|1]);

mn[x]=min(mn[x<<1],mn[x<<1|1]);

}

void pushdown(int x) {

**if** (lazy[x]) {

change(x<<1,lazy[x]);

change(x<<1|1,lazy[x]);

lazy[x]=0;

}

}

void build(int x,int l,int r) {

mxr[x]=mxl[x]=mx[x]=mn[x]=lazy[x]=0;

**if** (l==r) {

mx[x]=mn[x]=a[l];

**return**;

} int mid=(l+r)/2;

build(x<<1,l,mid);

build(x<<1|1,mid+1,r);

pushup(x);

}

int query(int x,int l,int r,bool flag,int &vmin,int &vmax,int L,int R,int val) { *//flag:-> (top->bottom yes)*

**if** (l<=L&&R<=r) {

change(x,val);

vmin=mn[x]; vmax=mx[x];

**return** flag?mxr[x]:mxl[x];

} pushdown(x);

int mid=(L+R)/2,ret=0,mx1=-INF,mx2=-INF,mn1=INF,mn2=INF;

**if** (mid>=l) ret=max(ret,query(x<<1,l,r,flag,mn1,mx1,L,mid,val));

**if** (r>mid) ret=max(ret,query(x<<1|1,l,r,flag,mn2,mx2,mid+1,R,val));

**if** (flag) ret=max(ret,mx2-mn1);

**else** ret=max(ret,mx1-mn2);

vmax=max(mx1,mx2);

vmin=min(mn1,mn2);

pushup(x);

**return** ret;

}

int n,q; int i,j,k;

int u,v,val;

int b[maxn];

vector<int> edge[maxn];

int sz[maxn],fa[maxn],dep[maxn],son[maxn],top[maxn],id[maxn];

void dfs1(int u,int from,int depth) {

int v,i,mx=-1;

sz[u]=1; fa[u]=from; dep[u]=depth; son[u]=0;

REP(i,edge[u].size()) {

v=edge[u][i];

**if** (v==from) **continue**;

dfs1(v,u,depth+1);

sz[u]+=sz[v];

**if** (sz[v]>mx) mx=sz[v],son[u]=v;

}

}

void dfs2(int u,int x) {

int v,i;

top[u]=x; id[u]=++tot;

**if** (son[u]) dfs2(son[u],x);

REP(i,edge[u].size()) {

v=edge[u][i];

**if** (v==fa[u]||v==son[u]) **continue**;

dfs2(v,v);

}

}

int Query(int x,int y,int val) {

int ret=0,mxx=-INF,mnx=INF,mxy=-INF,mny=INF,vmax,vmin;

**while** (top[x]!=top[y]) {

**if** (dep[top[x]]>dep[top[y]]) {

ret=max(ret,query(1,id[top[x]],id[x],0,vmin,vmax,1,tot,val));

ret=max(ret,vmax-mnx);

mxx=max(mxx,vmax); mnx=min(mnx,vmin);

x=fa[top[x]];

} **else** {

ret=max(ret,query(1,id[top[y]],id[y],1,vmin,vmax,1,tot,val));

ret=max(ret,mxy-vmin);

mxy=max(mxy,vmax); mny=min(mny,vmin);

y=fa[top[y]];

}

}

**if** (dep[x]>dep[y]) {

ret=max(ret,query(1,id[y],id[x],0,vmin,vmax,1,tot,val));

ret=max(ret,vmax-mnx);

mxx=max(mxx,vmax); mnx=min(mnx,vmin);

} **else** {

ret=max(ret,query(1,id[x],id[y],1,vmin,vmax,1,tot,val));

ret=max(ret,mxy-vmin);

mxy=max(mxy,vmax); mny=min(mny,vmin);

} ret=max(ret,mxy-mnx);

**return** ret;

}

int T;

int main() {

scanf("%d",&T);

**while** (T--) {

scanf("%d",&n);

FOR(i,1,n) scanf("%d",&b[i]);

FOR(i,1,n) edge[i].clear();

FOR(i,1,n-1) {

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

}

tot=0;

dfs1(1,0,1);

dfs2(1,1);

FOR(i,1,tot) a[id[i]]=b[i];

build(1,1,tot);

scanf("%d",&q);

REP(i,q) {

scanf("%d%d%d",&u,&v,&val);

printf("%d**\n**",Query(u,v,val));

}

}

}

轻重儿子分开维护

**// 题意: 更改链上的边col**

**// 更改某个链相邻的边col**

**// 查询黑点数**

**// 做法: 轻重边分开维护**

**struct** segment\_tree {

int val[maxn<<2],len[maxn<<2],lazy[maxn<<2];

void build(int x,int L,int R) {

len[x]=R-L+1; val[x]=0; lazy[x]=0;

**if** (L==R) **return**;

int mid=(L+R)/2;

build(x<<1,L,mid);

build(x<<1|1,mid+1,R);

}

void Inverse(int x) {

lazy[x]^=1; val[x]=len[x]-val[x];

}

void pushdown(int x) {

**if** (lazy[x]) {

Inverse(x<<1);

Inverse(x<<1|1);

lazy[x]=0;

}

}

void pushup(int x) {

val[x]=val[x<<1]+val[x<<1|1];

}

void update(int x,int l,int r,int L,int R) {

debug("update: %d %d %d**\n**",x,l,r);

**if** (l<=L&&R<=r) {Inverse(x); **return**;}

int mid=(L+R)/2;

pushdown(x);

**if** (l<=mid) update(x<<1,l,r,L,mid);

**if** (mid<r) update(x<<1|1,l,r,mid+1,R);

pushup(x);

}

int query(int x,int l,int r,int L,int R) {

**if** (l<=L&&R<=r) **return** val[x];

int mid=(L+R)/2,ret=0;

pushdown(x);

**if** (l<=mid) ret+=query(x<<1,l,r,L,mid);

**if** (mid<r) ret+=query(x<<1|1,l,r,mid+1,R);

pushup(x);

**return** ret;

}

} heavy,light;

vector<int> edge[maxn];

int fa[maxn],dep[maxn],sz[maxn],tot;

int top[maxn],id[maxn],son[maxn];

void dfs1(int u,int father,int depth) {

int mx=-1; sz[u]=1;

fa[u]=father; son[u]=0; dep[u]=depth;

**for** (int v:edge[u]) {

**if** (v==father) **continue**;

dfs1(v,u,depth+1); sz[u]+=sz[v];

**if** (sz[v]>mx) mx=sz[v],son[u]=v;

}

}

void dfs2(int u,int x) {

top[u]=x; id[u]=++tot;

**if** (son[u]) dfs2(son[u],x);

**for** (int v:edge[u]) {

**if** (v==fa[u]||v==son[u]) **continue**;

dfs2(v,v);

}

}

**inline** void InverseEdge(int x,int y) {

**while** (top[x]!=top[y]) {

**if** (dep[top[x]]<dep[top[y]]) swap(x,y);

heavy.update(1,id[top[x]],id[x],1,n);

x=fa[top[x]];

}

**if** (dep[x]>dep[y]) swap(x,y);

**if** (son[x]) heavy.update(1,id[son[x]],id[y],1,tot);

}

**inline** void InverseNode(int x,int y) {

**while** (top[x]!=top[y]) {

**if** (dep[top[x]]<dep[top[y]]) swap(x,y);

light.update(1,id[top[x]],id[x],1,n);

heavy.update(1,id[top[x]],id[top[x]],1,n);

**if** (son[x]) heavy.update(1,id[son[x]],id[son[x]],1,n);

x=fa[top[x]];

}

**if** (dep[x]>dep[y]) swap(x,y);

light.update(1,id[x],id[y],1,tot);

heavy.update(1,id[x],id[x],1,n);

**if** (son[y]) heavy.update(1,id[son[y]],id[son[y]],1,n);

}

**inline** int Query(int x,int y) {

int ret=0;

**while** (top[x]!=top[y]) {

**if** (dep[top[x]]<dep[top[y]]) swap(x,y);

**if** (top[x]!=x) ret+=heavy.query(1,id[son[top[x]]],id[x],1,n);

ret+=heavy.query(1,id[top[x]],id[top[x]],1,n)^light.query(1,id[fa[top[x]]],id[fa[top[x]]],1,n);

x=fa[top[x]];

}

**if** (dep[x]>dep[y]) swap(x,y);

**if** (son[x]) ret+=heavy.query(1,id[son[x]],id[y],1,n);

**return** ret;

}

int TaskA() {

int i;

scanf("%d",&n); tot=0;

FOR(i,1,n) edge[i].clear();

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

} dfs1(1,0,0); dfs2(1,1);

heavy.build(1,1,n);

light.build(1,1,n);

scanf("%d",&q);

REP(i,q) {

int op,u,v;

scanf("%d%d%d",&op,&u,&v);

**if** (op==1) InverseEdge(u,v);

**if** (op==2) InverseNode(u,v);

**if** (op==3) printf("%d**\n**",Query(u,v));

}

**return** 0;

}

链分治, 动态维护树上dp

**// f[x]:this\_ans=max(g[x]+f[heavy],0)**

**// g[x]:light\_ans=A[x]+sigma{f[light]}**

**// w[x]:dp[heavy\_son]**

**// 把轻链和重链分开维护, 在重链上一个序列上DP**

**// 题意是更改某点值, 查询联通块的最大权重和**

**struct** heap {

multiset<ll> S;

**inline** void ins(ll x) {

S.insert(x);

}

**inline** void del(ll x) {

multiset<ll>::iterator it=S.lower\_bound(x);

**if** (it!=S.end()) S.erase(it);

}

**inline** ll top() {

**if** (!S.size()) **return** 0;

**return** \*S.rbegin();

}

} SON[maxn]; *// light*

vector<int> edge[maxn];

int fa[maxn],dep[maxn],sz[maxn],tot;

int top[maxn],id[maxn],rid[maxn],son[maxn],leaf[maxn];

void dfs1(int u,int father,int depth) {

int mx=-1,i; sz[u]=1;

fa[u]=father; son[u]=0; dep[u]=depth;

REP(i,(int)edge[u].size()) {

int v=edge[u][i];

**if** (v==father) **continue**;

dfs1(v,u,depth+1); sz[u]+=sz[v];

**if** (sz[v]>mx) mx=sz[v],son[u]=v;

}

}

int A[maxn];

*// f[x]:this\_ans=max(g[x]+f[heavy],0)*

*// g[x]:light\_ans=A[x]+sigma{f[light]}*

*// w[x]:dp[heavy\_son]*

ll f[maxn],g[maxn],w[maxn];

void dfs2(int u,int x) {

top[u]=x; id[u]=++tot; rid[tot]=u;

g[u]=A[u]; f[u]=0; int i;

**if** (son[u]) dfs2(son[u],x);

REP(i,(int)edge[u].size()) {

int v=edge[u][i];

**if** (v==fa[u]||v==son[u]) **continue**;

dfs2(v,v); SON[u].ins(w[v]);

g[u]+=f[v]; max\_(w[u],w[v]);

} **if** (son[u]) {

leaf[u]=leaf[son[u]];

max\_(f[u],g[u]+f[son[u]]);

max\_(w[u],w[son[u]]);

} **else** leaf[u]=u;

max\_(f[u],g[u]); max\_(w[u],f[u]);

}

**struct** node {

ll ls,rs,sum,ans;

node(ll val=0) {sum=val; ls=rs=ans=max(0ll,val);}

} T[maxn<<2];

node merge(**const** node &A,**const** node &B) {

node ret;

ret.ls=max(A.ls,A.sum+B.ls);

ret.rs=max(B.rs,B.sum+A.rs);

ret.ans=max(A.ans,B.ans);

ret.ans=max(ret.ans,A.rs+B.ls);

ret.sum=A.sum+B.sum;

**return** ret;

}

*// f[x]:this\_ans=max(g[x]+f[heavy],0)*

*// g[x]:light\_ans=A[x]+sigma{f[light]}*

void build(int x,int L,int R) {

**if** (L==R) {

T[x]=node(g[rid[L]]);

max\_(T[x].ans,SON[rid[L]].top());

**return**;

} int mid=(L+R)/2;

build(x<<1,L,mid);

build(x<<1|1,mid+1,R);

T[x]=merge(T[x<<1],T[x<<1|1]);

}

void update(int x,int pos,int L,int R) {

**if** (L==R) {

T[x]=node(g[rid[L]]);

max\_(T[x].ans,SON[rid[L]].top());

**return**;

} int mid=(L+R)/2;

**if** (pos<=mid) update(x<<1,pos,L,mid);

**if** (mid<pos) update(x<<1|1,pos,mid+1,R);

T[x]=merge(T[x<<1],T[x<<1|1]);

}

node query(int x,int l,int r,int L,int R) {

**if** (l<=L&&R<=r) **return** T[x];

int mid=(L+R)/2;

**if** (r<=mid) **return** query(x<<1,l,r,L,mid);

**if** (mid<l) **return** query(x<<1|1,l,r,mid+1,R);

**return** merge(query(x<<1,l,r,L,mid),query(x<<1|1,l,r,mid+1,R));

}

**inline** void Update(int x,ll y) {

g[x]-=A[x]; A[x]=y; g[x]+=A[x];

**while** (x) {

update(1,id[x],1,n);

node nxtval=query(1,id[top[x]],id[leaf[x]],1,n);

ll initw=w[top[x]]; w[top[x]]=nxtval.ans;

ll initg=f[top[x]]; f[top[x]]=nxtval.ls;

x=fa[top[x]];

**if** (x) {

g[x]-=initg;

g[x]+=nxtval.ls;

SON[x].del(initw);

SON[x].ins(nxtval.ans);

}

}

}

**inline** ll Query(int x) {

**return** query(1,id[x],id[leaf[x]],1,n).ans;

}

int main() {

int i;

scanf("%d%d",&n,&q); tot=0;

FOR(i,1,n) scanf("%d",&A[i]);

FOR(i,1,n) edge[i].clear();

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

} dfs1(1,0,0); dfs2(1,1);

FOR(i,1,n) debug("%d ",id[i]); deputs("");

FOR(i,1,n) debug("%d ",rid[i]); deputs("");

build(1,1,n);

REP(i,q) {

char op[2];

scanf("%s",op);

**if** (op[0]=='M') {

int x; ll y;

scanf("%d%lld",&x,&y);

Update(x,y);

} **else** {

int x;

scanf("%d",&x);

printf("%lld**\n**",Query(x));

}

}

**return** 0;

}

DSU on tree

**//大概意思就是轻儿子记录答案, 重儿子不清空, 最后把轻儿子的贡献放到重儿子上; 如果是基于深度可合并的, 长链剖分是O(n)的**

**// CF741D 辣鸡题**

**// 问你重排能回文的最长串多长**

**// 直接上就可以了... 看下dfs顺序就行了**

vector<int> edge[maxn];

int sz[maxn],son[maxn];

void dfs1(int x) {

int mx=0; sz[x]=1;

**for** (int v:edge[x]) {

dfs1(v); sz[x]+=sz[v];

**if** (sz[v]>mx) son[x]=v,mx=sz[v];

}

}

int A[maxn],dep[maxn];

int ans[maxn],MX[1<<22|7];

map<int,int> MP[maxn];

int Merge(map<int,int> &A,map<int,int> &B,int x) { *//B->A*

int ret=0,i;

**for** (**auto** now:B) {

int p=now.first,l=now.second;

**if** (MX[p]) ret=max(ret,MX[p]+l-2\*dep[x]);

REP(i,22) {

p=now.first^(1<<i);

*// printf("now=%d; p=%d; %d %d %d\n",now.first,p,MX[p],l,dep[x]);*

**if** (MX[p]) ret=max(ret,MX[p]+l-2\*dep[x]);

}

}*//merge*

**for** (**auto** now:B) {

int p=now.first,l=now.second;

MX[p]=max(MX[p],l); A[p]=MX[p];

} map<int,int>().swap(B);

**return** ret;

}

void dfs2(int x) {

**for** (int v:edge[x]) **if** (v!=son[x]) {

dfs2(v); ans[x]=max(ans[x],ans[v]);

**for** (**auto** now:MP[v]) MX[now.first]=0;

} **if** (son[x]) {

dfs2(son[x]); ans[x]=max(ans[x],ans[son[x]]);

}*//cal*

MP[x][A[x]]=dep[x];

**if** (son[x]) {

ans[x]=max(ans[x],Merge(MP[son[x]],MP[x],x));

swap(MP[x],MP[son[x]]);

} **else** MX[A[x]]=dep[x];

**for** (int v:edge[x]) **if** (v!=son[x]) {

ans[x]=max(ans[x],Merge(MP[x],MP[v],x));

}

}

int main() {

int n,i,j,k;

char c;

scanf("%d",&n);

FOR(i,2,n) {

int fa;

scanf("%d %c",&fa,&c);

A[i]=A[fa]^(1<<(c-'a'));

dep[i]=dep[fa]+1;

edge[fa].push\_back(i);

} dfs1(1); dfs2(1);

FOR(i,1,n) printf("%d ",ans[i]);

**return** 0;

}

LCA

**树链剖分: 略**

**Tarjan:**

vector<int> edge[maxn];

int fa1[maxn],fa2[maxn];

**inline** int getfa(int \*fa,int x) {

**if** (fa[x]==x) **return** x;

**return** fa[x]=getfa(fa,fa[x]);

}

int n,m,q,i,k,u,v;

int ans[maxn];

vector<pair<int,int> > Q[maxn];*//v,id*

void dfs(int x) {

int i;

**for** (int v:edge[x]) {

dfs(v); fa2[v]=x;

}

REP(i,Q[x].size()) {

**if** (fa2[Q[x][i].first]!=Q[x][i].first)

ans[Q[x][i].second]=getfa(fa2,Q[x][i].first);

}

}

void solve() {

scanf("%d%d%d",&n,&m,&q);

FOR(i,1,n) fa1[i]=fa2[i]=i;

REP(i,m) {

scanf("%d%d",&u,&v);

edge[u].push\_back(v);

fa1[v]=u;

}

REP(i,q) {

scanf("%d%d%d",&k,&u,&v);

**if** (k==1) {

**if** (getfa(fa1,u)!=getfa(fa1,v)) ans[i]=-1;

**else** {

**if** (u==v) ans[i]=u;

**else** {

Q[u].push\_back(make\_pair(v,i));

Q[v].push\_back(make\_pair(u,i));

}

}

} **else** {

edge[u].push\_back(v);

fa1[v]=u; ans[i]=0;

}

}

FOR(i,1,n) **if** (fa1[i]==i) dfs(i);

FOR(i,1,n) edge[i].clear(),Q[i].clear();

REP(i,q) **if** (ans[i]) printf("%d**\n**",ans[i]);

}

倍增

int fa[maxn][21];

int n,i,j;

int dep[maxn];

vector<int> edge[maxn];

void dfs(int x,int depth) {

dep[x]=depth;

**for** (int v:edge[x]) dfs(v,depth+1);

}

int lca(int x,int y) {

int i;

**if** (dep[x]<dep[y]) swap(x,y);

rREP(i,20) **if** (dep[x]-dep[y]>=1<<i) x=fa[x][i];

**if** (x==y) **return** x;

rREP(i,20) **if** (fa[x][i]!=fa[y][i]) x=fa[x][i],y=fa[y][i];

**return** fa[x][0];

}

int dis(int x,int y) {

**return** dep[x]+dep[y]-2\*dep[lca(x,y)];

}

int kthfa(int x,int k) {

int i;

rREP(i,20) **if** ((k>>i)&1) x=fa[x][i];

**return** x;

}

int walk(int x,int y,int d) {

int f=lca(x,y);

**if** (dep[x]-dep[f]>=d) **return** kthfa(x,d);

**return** kthfa(y,dep[x]+dep[y]-2\*dep[f]-d);

}

虚树ST表求lca；四毛子定理

**// 题意:问最少去掉几个未标记点可以把所有的标记点全分开**

**// 做法:建虚树然后树上DP**

**// 虚树板子,注意:sort过程可以提到外边去**

**// 注意, 原先有的标记有的时候会到边上, 需要特判的, 千万不要if**

**const** int K\_K=16+1;

**namespace** ST {

**const** int maxn=1<<K\_K;

int ST[maxn][K\_K],t\_t[maxn];*//only L*

**inline** void initST(int n,**const** int id[],**const** int A[]) { *//0 to n-1*

int i,j; assert(n<=maxn);

rep(i,1,n) t\_t[i]=t\_t[i>>1]+1;

REP(i,n) {

ST[i][0]=id[i];

rep(j,1,t\_t[i+1]) {

int u=ST[i-(1<<(j-1))][j-1],v=ST[i][j-1];

ST[i][j]=A[u]>=A[v]?u:v;

}

}

}

**inline** int RMQ(int x,int y,**const** int A[]) {

int t=t\_t[y-x+1]-1;

x=ST[x+(1<<t)-1][t]; y=ST[y][t];

**return** A[x]>=A[y]?x:y;

}

}

**namespace** RMQ01 { *//start from 0; return the First Maximize Value Position*

**const** int maxn=1<<K\_K;

**inline** int B(int x) {**return** x>>4;}

int BLOCK[maxn],MAXP[1<<16|1],BMAX[maxn];

**inline** int value(int l,int r) {*//0000xxxx1111 // & mid | right*

int mid=((1<<((r&15)+1))-1)^((1<<((l&15)+1))-1);

int right=((1<<16)-1)^((1<<((r&15)+1))-1);

**return** MAXP[(BLOCK[B(l)]&mid)|right]+(l&~15);

}

void dfs\_value(int x=0,int depth=15,int maxval=0,int maxpos=0) {

**if** (maxval<=0) maxpos=0,maxval=0;

**if** (depth==0) {MAXP[x<<1]=MAXP[x<<1|1]=maxpos; **return**;}

dfs\_value(x<<1,depth-1,maxval+1,maxpos+1);

dfs\_value(x<<1|1,depth-1,maxval-1,maxpos+1);

}

void build(int n, int A[]) {*//0:+1; 1:-1 //start from (1-15); g[i]=A[i]-A[i-1]*

int i; dfs\_value(); **while** (n&15) A[n]=A[n-1]-1,n++; assert(B(n-1)<maxn);

memset(BLOCK,0,**sizeof**(BLOCK));

rrep(i,1,n) BLOCK[B(i)]=BLOCK[B(i)]<<1|(1-((A[i]-A[i-1]+1)>>1));

BLOCK[0]<<=1;

REP(i,B(n-1)+1) BMAX[i]=MAXP[BLOCK[i]]+(i<<4);

ST::initST(B(n-1)+1,BMAX,A);

}

int RMQ01(int x,int y,**const** int A[]) {

**if** (x>y) swap(x,y);

int rx=B(x)<<4|15,ly=B(y)<<4;

**if** (y<=rx) **return** value(x,y);

int u=value(x,rx),v;

**if** (rx+1!=ly) v=ST::RMQ(B(x)+1,B(y)-1,A),u=A[u]>=A[v]?u:v;

v=value(ly,y); u=A[u]>=A[v]?u:v;

**return** u;

}

}

**namespace** LCA {

**struct** Edges {

int to; ll len; int next;

Edges(int \_to=0,ll \_len=0,int \_next=0):to(\_to),len(\_len),next(\_next) {}

} edge[maxn\*2]; int etot;

int head[maxn];*//for edges*

int fa[maxn],dep[maxn];

ll uplen[maxn],len[maxn]; *//uplen: single edge*

int id[maxn],dfn[maxn],tot;*//down*

int id2[maxn\*2],dfn2[maxn],tot2,value2[maxn\*2];*//down+up; dep2=-depth; just for 4-Russian Use*

**inline** void addedge(int u,int v,ll len) {

edge[++etot]=Edges(v,len,head[u]); head[u]=etot;

}

void dfs(int x=1,int f=0,int d=0,ll l=0) {

int i; dep[x]=d; len[x]=l;

++tot; id[tot]=x; dfn[x]=tot;

id2[tot2]=x; dfn2[x]=tot2; value2[tot2]=-d; tot2++;

**for** (i=head[x]; ~i; i=edge[i].next) **if** (edge[i].to!=f) {

int v=edge[i].to;

fa[v]=x; uplen[v]=edge[i].len;

dfs(v,x,d+1,l+edge[i].len);

value2[tot2]=-d; id2[tot2++]=x;

}

}

void init\_lca(int root=1) {

dfs(root);

RMQ01::build(tot2,value2);

}

**inline** int lca(int x,int y) {

**return** id2[RMQ01::RMQ01(dfn2[x],dfn2[y],value2)];

}

**inline** ll dis(int x,int y) {

**return** len[x]+len[y]-2\*len[lca(x,y)];

}

void init(int n) {

memset(head+1,0xff,n\***sizeof**(int));

etot=tot=tot2=0;

}

}

**namespace** vtree {

int S[maxn],top;

int pid[maxn],mark[maxn];

int vid[maxn],vfa[maxn];

LL vlen[maxn];

int cmp(int x,int y) {

**return** LCA::dfn[x]<LCA::dfn[y];

}

void addedge(int u,int v) {

vfa[v]=u; vlen[v]=LCA::dis(u,v);

}

int m;

void vbuild(int n) {

int i; m=0;

sort(pid+1,pid+1+n,cmp);

S[top=1]=pid[1];

mark[pid[1]]=1;

FOR(i,2,n) {

int f=LCA::lca(pid[i-1],pid[i]);

**while** (top&&LCA::dep[S[top]]>LCA::dep[f]) {

int v; vid[++m]=v=S[top--];

**if** (top&&LCA::dep[S[top]]>LCA::dep[f]) addedge(S[top],v);

**else** addedge(f,v);

} **if** (!top||S[top]!=f) S[++top]=f;

S[++top]=pid[i]; mark[pid[i]]=1;

} **while** (top-1) addedge(S[top-1],S[top]),vid[++m]=S[top--];

vid[++m]=S[1];

reverse(vid+1,vid+m+1);

}

void vclear() {

int i;

FOR(i,1,m) mark[vfa[vid[i]]]=0;

FOR(i,1,m) mark[vid[i]]=0;

}

}

int ans;

int cnt[maxn];

void solve() {

int i;

FOR(i,1,vtree::m) cnt[vtree::vid[i]]=0;

rFOR(i,1,vtree::m) {

int x=vtree::vid[i];

**if** (vtree::mark[x]) ans+=cnt[x],cnt[x]=1;

**else** **if** (cnt[x]>1) ans++,cnt[x]=0;

**if** (i>1) cnt[vtree::vfa[x]]+=cnt[x];

}

}

int vis[maxn];

int main() {

int i;

int n,q;

scanf("%d",&n);

LCA::init(n);

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

LCA::addedge(u,v,1); LCA::addedge(v,u,1);

} LCA::init\_lca();

scanf("%d",&q);

**while** (q--) {

int m,mark=0;

scanf("%d",&m);

FOR(i,1,m) scanf("%d",&vtree::pid[i]);

FOR(i,1,m) vis[vtree::pid[i]]=1;

FOR(i,1,m) **if** (vis[LCA::fa[vtree::pid[i]]]) mark=1;

FOR(i,1,m) vis[vtree::pid[i]]=0;

**if** (mark) {puts("-1"); **continue**;}

vtree::vbuild(m);

ans=0; solve();

vtree::vclear();

printf("%d**\n**",ans);

}

}

****Ladder长链剖分 k级祖先****

**namespace** ladder {

vector<int> edge[maxn];

int id[maxn]; int tot;

int fa[maxn][21],son[maxn],top[maxn],len[maxn],dep[maxn];

vector<int> ladder[maxn];

int upp[maxn];

void dfs(int x,int father=0) {

fa[x][0]=father; id[++tot]=x;

**for** (int v:edge[x]) **if** (v!=father) dfs(v,x);

}

void buildfa() {

int i,j; dep[id[1]]=0;

FOR(i,1,tot) rep(j,1,21) fa[i][j]=fa[fa[i][j-1]][j-1],dep[i]=dep[fa[i][0]]+1;

rFOR(i,1,tot) {

int o=0,x=id[i]; top[x]=x;

ladder[x].clear();

**for** (int v:edge[x]) **if** (v!=fa[x][0]){

**if** (!o||len[o]<len[v]) o=v;

} **if** (o) len[x]=len[o]+1; **else** o=0;

son[x]=o; top[x]=x;

} FOR(i,1,tot) **if** (son[id[i]]) top[son[id[i]]]=top[id[i]];

rFOR(i,1,tot) ladder[top[id[i]]].push\_back(id[i]);

FOR(i,2,tot) {

int x=id[i];

**if** (top[x]==x) {

**for** (int y=fa[x][0],c=len[x]; y&&c; y=fa[y][0],c--)

ladder[x].push\_back(y);

}

} upp[0]=-1;

FOR(i,1,tot) upp[i]=upp[i-1]+(i==(i&-i));

}

int prev(int x,int k) {

**if** (!k) **return** x;

**if** (dep[x]<=k) **return** 0;

x=fa[x][upp[k]]; k-=1<<upp[k];

k-=dep[x]-dep[top[x]]; x=top[x];

**return** ladder[x][len[x]+k];

}

}

**using** **namespace** ladder;

****最大团****

int n,ans;

int edge[maxn][maxn],cnt[maxn],vis[maxn];*//vis:元素*

bool dfs(int u,int pos) {

int i,j;

FOR(i,u+1,n) {

**if** (cnt[i]+pos<=ans) **return** 0;

**if** (edge[u][i]) {

REP(j,pos) **if** (!edge[i][vis[j]]) **break**;

**if** (j==pos) {

vis[pos]=i;

**if** (dfs(i,pos+1)) **return** 1;

}

}

}

**if** (pos>ans) {ans=pos;**return** 1;}

**return** 0;

}

int maxclique() {

int i; ans=-1;

rFOR(i,1,n) {

vis[0]=i;

dfs(i,1);

cnt[i]=ans;

} **return** ans;

}

****最小树形图(mlogn)****

**//不定根:新加一个节点，向所有点加一条INF的边，最后减一下即可**

**//主要思路:缩点**

**//输出路径思路:缩完点记录边,然后新建边记录等价关系**

**nm的做法:**

**namespace** O {

**const** int maxn=1e5+7;

**const** int maxm=1e5+7;

**struct** Edge {*//id,pre!=0;uid:替换*

int u,v,len;*//id->usedID(new),用于新建边*

Edge(int \_u=0,int \_v=0,int \_len=0):

u(\_u),v(\_v),len(\_len) {}

} edge[maxm]; int etot;

void init() {etot=0;}

void addedge(int u,int v,int len) {

edge[++etot]=Edge(u,v,len);

}

**struct** \_info {*//pre为preuid*

int pre,len,eid;*//position(id) and length*

void init(int \_pre,int \_len,int \_eid){

pre=\_pre; len=\_len; eid=\_eid;

}

} Info[maxn\*2];

*//id:circle\_id(father);top:tree\_anc*

**inline** int getfa(int fa[],int x) {

**if** (fa[x]==x) **return** x;

**return** fa[x]=getfa(fa,fa[x]);

} int id[maxn\*2],top[maxn\*2];*//并查集;*

int idfa[maxn\*2],useid[maxn\*2];*//changes\_fa*

int used[maxm];*//output; 记录edge*

int solve(int root,int n) {

int ret=0,i,lastnid=n;

FOR(i,1,n) id[i]=top[i]=i,Info[i].len=INF;

FOR(i,1,etot) { *//initialize*

Edge &e=edge[i];

**if** (e.u!=e.v&&e.len<Info[e.v].len)

Info[e.v].init(e.u,e.len,i);

}

FOR(i,1,lastnid) {

**if** (i==root) **continue**;

**if** (Info[i].len==INF) **return** -1;

\_info &info=Info[i];

int f=getfa(top,info.pre);

*//choose; 之后再更新*

useid[i]++; ret+=info.len; idfa[i]=i;

**if** (f==i) {*//circle*

int k; ++lastnid; Info[lastnid].len=INF;

top[lastnid]=id[lastnid]=lastnid;

**for** (int x=getfa(id,info.pre); ;x=getfa(id,Info[x].pre)) {

FOR(k,1,etot) {

Edge &e=edge[k];

**if** (k==info.eid) e.len=INF;

**if** (e.len==INF) **continue**;*//removed*

**if** (getfa(id,e.v)==x) *//must\_ok*

e.len-=Info[x].len;

}

*//use and delete*

id[x]=top[x]=idfa[x]=lastnid;*//缩环*

**if** (x==i) **break**;

} *//update edges*

FOR(k,1,etot) {

Edge &e=edge[k];

**if** (e.len==INF) **continue**;*//removed*

**if** (getfa(id,e.v)==lastnid) {*//must\_ok*

**if** (getfa(id,e.u)!=lastnid&&e.len<Info[lastnid].len)

Info[lastnid].init(e.u,e.len,k);*//直接这样应该ok?*

}

}

} **else** top[getfa(id,i)]=info.pre;*//getfa=getid*

}

rFOR(i,1,lastnid) **if** (useid[i]) {*//remove\_to\_top*

**static** int ids[maxn];

int x=0,k;

**for** (int k=edge[Info[i].eid].v;k!=i;k=idfa[k]) ids[x++]=k;

REP(k,x) idfa[ids[k]]=i,useid[ids[k]]=0;

used[Info[i].eid]=1;

}

**return** ret;

}

}

**mlogn的做法：**

**namespace** heap {

**const** int maxn=1e5+7;

**struct** node {

int l,r,len;

int u,v,val,lz;

node(int \_u=0,int \_v=0,int \_val=0):

u(\_u),v(\_v),val(\_val) {l=r=len=0; lz=0;}

} T[maxn]; int tot;

*//不能直接swap x和儿子, 否则可能不满足堆性质*

void update(int x,int val) {

T[x].lz+=val; T[x].val-=val;

}

void pushdown(int x) {

**if** (T[x].lz) {

**if** (T[x].l) update(T[x].l,T[x].lz);

**if** (T[x].r) update(T[x].r,T[x].lz);

T[x].lz=0;

}

}

int merge(int x,int y) {

**if** (!x||!y) **return** x|y;

pushdown(x); pushdown(y);

**if** (T[x].val>T[y].val) swap(x,y);

T[x].r=merge(T[x].r,y);

**if** (T[T[x].l].len<T[T[x].r].len) swap(T[x].l,T[x].r);

T[x].len=T[T[x].r].len+1;

**return** x;

}

int pop(int x) {

pushdown(x);

**return** merge(T[x].l,T[x].r);

}

}

**namespace** O {

**const** int maxn=1e5+7;

**const** int maxm=1e5+7;

int Root[maxm],etot;

void init() {etot=0;}

void addedge(int u,int v,int len) {

heap::T[++etot]=heap::node(u,v,len);

}

**struct** \_info {*//pre为preuid;maxtot=maxm*

int pre,len,eid;*//position(id) and length*

void init(int \_pre,int \_len,int \_eid) {

pre=\_pre; len=\_len; eid=\_eid;

}

} Info[maxm];

*//id:circle\_id(father);top:tree\_anc*

**inline** int getfa(int fa[],int x) {

**if** (fa[x]==x) **return** x;

**return** fa[x]=getfa(fa,fa[x]);

} int id[maxm],top[maxm];*//并查集;*

bool getTopValue(int i) {

**while** (Root[i]) {

heap::node &e=heap::T[Root[i]];

**if** (getfa(id,e.u)==getfa(id,e.v)) {

Root[i]=heap::pop(Root[i]);

} **else** {

Info[i].init(e.u,e.val,Root[i]);

Root[i]=heap::pop(Root[i]);

*// printf("%d: %d; pos=%d\n",e.v,e.val,Root[i]);*

**return** 1;

}

} **return** 0;

}

int idfa[maxm],useid[maxm];*//changes\_fa*

int used[maxm];*//output; 记录edge*

int solve(int root,int n) {

int ret=0,i,lastnid=n;

FOR(i,1,n) id[i]=top[i]=i,Info[i].len=INF;

FOR(i,1,etot) { *//initialize*

heap::node &e=heap::T[i];

Root[e.v]=heap::merge(Root[e.v],i);

}

FOR(i,1,n) **if** (i!=root&&!getTopValue(i)) **return** -1;

*// puts("ok1");*

FOR(i,1,lastnid) {

**if** (i==root) **continue**;

\_info &info=Info[i];

int f=getfa(top,info.pre);

*//choose; 之后再更新*

useid[i]++; ret+=info.len; idfa[i]=i;

**if** (f==i) {*//circle*

++lastnid; Info[lastnid].len=INF;

top[lastnid]=id[lastnid]=lastnid;

**for** (int x=getfa(id,info.pre); ; x=getfa(id,Info[x].pre)) {

heap::update(Root[x],Info[x].len);

Root[lastnid]=heap::merge(Root[lastnid],Root[x]);

id[x]=top[x]=idfa[x]=lastnid;*//缩环*

**if** (x==i) **break**;

} *//update edges*

**if** (!getTopValue(lastnid)) **return** -1;

} **else** top[getfa(id,i)]=info.pre;*//getfa=getid*

}

rFOR(i,1,lastnid) **if** (useid[i]) {*//remove\_to\_top*

**static** int ids[maxn];

int x=0,k;

**for** (int k=heap::T[Info[i].eid].v; k!=i; k=idfa[k]) ids[x++]=k;

REP(k,x) idfa[ids[k]]=i,useid[ids[k]]=0;

used[Info[i].eid]=1;

}

**return** ret;

}

*//debug*

vector<pair<int,int> > check\_edge[maxn];

bool vis[maxn];

void addcheck\_edge\_check(int u,int v,int i) {

check\_edge[u].push\_back(make\_pair(v,i));

}

void bfs(int x,int n) {*//x=root*

queue<int> Q;

Q.push(x); vis[x]=1;

**while** (Q.size()) {

int x=Q.front(); Q.pop();

**for** (**auto** e:check\_edge[x]) {

int i=e.second;

**if** (vis[e.first]) **continue**;

**if** (!used[i]) **continue**;

Q.push(e.first); vis[e.first]=1;

}

} int i;

FOR(i,1,n) {

**if** (!vis[i]) debug("no! %d**\n**",i);

}

}

}

int ini[maxn];

int n,m,i;

int u,v,w;

int main() {

freopen("input.txt","r",stdin);

freopen("output.txt","w",stdout);

scanf("%d%d",&n,&m);

O::init();

FOR(i,1,m) {

**if** (scanf("%d%d%d",&u,&v,&w)!=3) {

puts("input not right");

**return** 0;

};

ini[i]=w;

O::addedge(u,v,w);

O::addcheck\_edge\_check(u,v,i);

}

int ans=O::solve(1,n);

printf("%d**\n**",ans);

**if** (ans!=-1) {

O::bfs(1,n);

FOR(i,1,m) **if** (ini[i]&&O::used[i])

printf("%d ",i),ans--;

**if** (ans) printf("**\n**notok: %d**\n**",ans);

}

}

****一般图最大匹配 带花树****

*//缩奇环*

int n,m;

vector<int> edge[maxn];

bool inQueue[maxn];

int belong[maxn];

int getbelong(int x) {

**if** (belong[x]==x) **return** x;

**return** belong[x]=getbelong(belong[x]);

}

int match[maxn],nxt[maxn],mark[maxn],vis[maxn];

int cnt;

queue<int> Q;

int used[maxn];

int lca(int u,int v) {

cnt++;

**while** (1) {

u=getbelong(u);

**if** (vis[u]==cnt) **return** u;

vis[u]=cnt;

u=nxt[match[u]];

**if** (v) swap(u,v);

}

}

void merge(int u,int p) {

**while** (u!=p) {

int mu=match[u],v=nxt[mu];

**if** (getbelong(v)!=p) nxt[v]=mu;

**if** (mark[mu]==2) mark[mu]=1,Q.push(mu);

**if** (mark[v]==2) mark[v]=1,Q.push(v);

int x,y;

x=getbelong(u),y=getbelong(mu);

**if** (x!=y) belong[x]=y;

x=getbelong(mu),y=getbelong(v);

**if** (x!=y) belong[x]=y;

u=v;

}

}

void solve(int s) { *//增广*

int i;

FOR(i,1,n) belong[i]=i,mark[i]=nxt[i]=0;

**while** (Q.size()) Q.pop();

Q.push(s);

**while** (Q.size()) {

**if** (match[s]) **return**;

int u=Q.front();

Q.pop();

**for** (int v:edge[u]) {

**if** (match[u]==v) **continue**;

**if** (getbelong(u)==getbelong(v)) **continue**;

**if** (mark[v]==2) **continue**; *//T型点*

**if** (mark[v]==1) { *//S型点,缩点*

int p=lca(u,v);

**if** (getbelong(u)!=p) nxt[u]=v;

**if** (getbelong(v)!=p) nxt[v]=u;

merge(u,p);

merge(v,p);

} **else** **if** (!match[v]) { *//增广*

nxt[v]=u;

**for** (int x=v; x;) {

int y=nxt[x],xx=match[y];

match[x]=y;

match[y]=x;

x=xx;

}

**break**;

} **else** {

nxt[v]=u;

mark[match[v]]=1;

Q.push(match[v]);

mark[v]=2;

}

}

}

}

bool E[maxn][maxn];

int ans;

int main() {

scanf("%d%d",&n,&m);

int i;

**while** (m--) {

int u,v;

scanf("%d%d",&u,&v);

**if** (u!=v&&!E[u][v]) {

edge[u].push\_back(v);

edge[v].push\_back(u);

E[u][v]=E[v][u]=1;

}

}

memset(match,0,**sizeof**(match));

FOR(i,1,n) **if** (!match[i]) solve(i);

FOR(i,1,n) **if** (match[i]) ans++;

ans/=2;

printf("%d**\n**",ans);

FOR(i,1,n) printf("%d ",match[i]);

}

****树分块 高度分块****

**//题意: 给两颗树, 树上有边**

**//问你: T1的1->i和T2的1->i的所有边加入计算后有多少联通块**

**//做法: 先把两棵树分成sqrt个块, 将query放到两个树的块上**

**//然后直接从上面转移+回滚, 最多sqrt的length**

**struct** Changes {

int x,y,ini;

Changes(int \_x=0,int \_y=0,int \_ini=0):x(\_x),y(\_y),ini(\_ini) {};

} changes[maxn\*50]; int top;

vector<int> E[maxn],V1,V2;

int fa1[maxn],fa2[maxn],SIZE;

int id1[maxn],id2[maxn];

int last[507][507];

vector<Changes> queries[507][507];

void dfs1(int x,int f,int dep) {

fa1[x]=f; id1[x]=-1;

**if** (!(dep%SIZE)) id1[x]=V1.size(),V1.push\_back(x);

**for** (int v:E[x]) **if** (v!=f) dfs1(v,x,dep+1);

} void dfs2(int x,int f,int dep) {

fa2[x]=f; id2[x]=-1;

**if** (!(dep%SIZE)) id2[x]=V2.size(),V2.push\_back(x);

**for** (int v:E[x]) **if** (v!=f) dfs2(v,x,dep+1);

}

int ux[maxn],uy[maxn],vx[maxn],vy[maxn];

int fa[maxn],sz[maxn],nowans;

void merge\_(int x,int y) {

**while** (x!=fa[x]) x=fa[x];

**while** (y!=fa[y]) y=fa[y];

**if** (x==y) **return**;

**if** (sz[x]>sz[y]) swap(x,y);

changes[++top]=Changes(x,y);

fa[x]=y; nowans--; sz[y]+=sz[x];

debug("merge\_OK (%d %d)-%d ans=%d**\n**",x,y,top,nowans);

}

void revert(int x,int y) {

debug("revert to %d %d**\n**",x,y);

**while** (top>last[id1[x]][id2[y]]) {

**auto** now=changes[top--]; nowans++;

fa[now.x]=now.x; fa[now.y]=now.y;

sz[now.y]-=sz[now.x];

debug("revert\_OK (%d %d)-%d+1 ans=%d**\n**",now.x,now.y,top,nowans);

}

}

void commit(int x,int y) {

int tx=x,ty=y;

**while** (id1[tx]==-1) tx=fa1[tx];

**while** (id2[ty]==-1) ty=fa2[ty];

*// if (tx!=lastx||ty!=lasty)*

revert(tx,ty);

**while** (x!=tx) merge\_(ux[x],vx[x]),x=fa1[x];

**while** (y!=ty) merge\_(uy[y],vy[y]),y=fa2[y];

}

int ans[maxn];

int main() {

int T;

int i,j,k;

scanf("%d",&T);

**while** (T--) {

int n,m;

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%d%d",&ux[i],&vx[i]);

SIZE=sqrt(n)\*1.1; **if** (SIZE==0) SIZE=1;

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

E[u].push\_back(v);

E[v].push\_back(u);

} V1.clear(); dfs1(1,0,0);

FOR(i,1,n) E[i].clear();

FOR(i,1,n) scanf("%d%d",&uy[i],&vy[i]);

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v);

E[u].push\_back(v);

E[v].push\_back(u);

} V2.clear(); dfs2(1,0,0);

FOR(i,1,n) E[i].clear();

*// FOR(i,1,n) debug("%d %d\n",fa1[i],id1[i]);*

*// FOR(i,1,n) debug("%d %d\n",fa2[i],id2[i]);*

FOR(i,1,n) {

int u,v,x,y;

u=v=i; x=u; y=v;

**while** (id1[x]==-1) x=fa1[x];

**while** (id2[y]==-1) y=fa2[y];

queries[id1[x]][id2[y]].push\_back(Changes(u,v,i));

} deputs("okay");

FOR(i,1,m) fa[i]=i,sz[i]=1;

top=0; nowans=m;

**for** (int x:V1) {

**for** (int y:V2) {

**if** (x==1&&y==1) {

merge\_(ux[x],vx[x]);

merge\_(uy[y],vy[y]);

} **else** **if** (y==1) {

commit(fa1[x],y);

merge\_(ux[x],vx[x]);

} **else** {

commit(x,fa2[y]);

merge\_(uy[y],vy[y]);

} last[id1[x]][id2[y]]=top;

**for** (**auto** now:queries[id1[x]][id2[y]]) {

debug("query %d %d**\n**",now.x,now.y);

commit(now.x,now.y);

ans[now.ini]=nowans;

} queries[id1[x]][id2[y]].clear();

}

}

FOR(i,1,n) printf("%d**\n**",ans[i]);

}

**return** 0;

}

****树哈希****

**// 题意: A树有多少个和B树同构的子树(B总共12个节点)**

**// 做法: 对B树进行哈希, 然后构造转移方案并转移**

**namespace** tree\_hash {

**const** int maxk=12;

vector<int> baseedge[1<<maxk|7];

void init(int n) {

int i,j;

REP(i,(1<<n)) REP(j,n) **if** ((i>>j)&1)

baseedge[i].push\_back(j);

}

**typedef** unsigned int type;

**struct** tree {

int e[maxk];*//baseedge*

void init(int n) {memset(e,0,n\***sizeof**(int));}

type dfs(int x,int fa) {*//encode*

*// printf("dfs: %d %d\n",x,fa);*

**static** int sz[maxk];

type ret=0; sz[x]=1;

vector<pair<type,int> > tmp;*//count*

**for** (int v:baseedge[e[x]]) **if** (v!=fa) {

type val=dfs(v,x)<<1|1; sz[x]+=sz[v];

tmp.push\_back(make\_pair(val,sz[v]\*2));

} sort(tmp.begin(), tmp.end());

reverse(tmp.begin(), tmp.end());

**for** (pair<type,int>v:tmp)

ret=(ret<<v.second)|v.first;

**return** ret;

}

type unrooted(int n) {

vector<type> tmp; int i;

REP(i,n) tmp.push\_back(dfs(i,i));

sort(tmp.begin(), tmp.end());

**return** tmp[0];*//字典序minimize*

}

};

set<type> S;*//hash\_answer*

**typedef** pair<tree,type> ptt;

vector<pair<tree,type> > Trees[13];*//size*

tree merge(**const** tree &x,**const** tree &y,int sizex,int sizey) {

tree ret; int i,j;

ret.init(sizex+sizey);

REP(i,sizex) REP(j,sizex) **if** ((x.e[i]>>j)&1) {

ret.e[i]|=1ull<<j;

ret.e[j]|=1ull<<i;

}

REP(i,sizey) REP(j,sizey) **if** ((y.e[i]>>j)&1) {

ret.e[i+sizex]|=1ull<<(j+sizex);

ret.e[j+sizex]|=1ull<<(i+sizex);

}

ret.e[0]|=1ull<<sizex;

ret.e[sizex]|=1ull<<0;

**return** ret;

}

int tot;

short id[1<<(maxk\*2)];*//last=0所以可以去掉*

**typedef** pair<type,int> pti;

vector<pti> edges[8007];*//这里是个暴力, 在外边再搞一次比较好*

void getall(int n) {

tree ini; ini.init(1);*//ini.dfs=0*

Trees[1].push\_back(make\_pair(ini,0));

int \_,o; tot=0;

S.insert(0); id[0]=++tot;

FOR(\_,1,n-1) {*//size*

**for** (ptt tmp:Trees[\_]) {

int i; tree &ori=tmp.first;

REP(i,\_) {

tree nxt=ori;

nxt.e[\_]=1ull<<i;

nxt.e[i]|=1ull<<\_;

type v=nxt.dfs(0,0);

**if** (S.count(v)) **continue**;

Trees[\_+1].push\_back(make\_pair(nxt,v));

S.insert(v); id[v]=++tot;

*// assert(v<(1<<(maxk\*2-1)));*

}

}

}

*// int all=0;*

*// FOR(\_,1,n) {*

*// printf("%d: %d\n",\_,(int)Trees[\_].size());*

*// all+=Trees[\_].size();*

*// // for (tree ori:Trees[\_]) pr2(ori.dfs(0,0),\_\*2-2); puts("");*

*// } printf("size\_all=%d\n",all);*

*// printf("tot=%d\n",tot);*

FOR(\_,1,n) FOR(o,1,n-\_){

**for** (ptt tmpx:Trees[\_]) {

tree &x=tmpx.first; type ox=tmpx.second;

**for** (ptt tmpy:Trees[o]) {

tree &y=tmpy.first; type oy=tmpy.second;

type o\_nxt=merge(x,y,\_,o).dfs(0,0);

edges[id[ox]].push\_back(make\_pair(id[oy],id[o\_nxt]));

}

}

}

}

}

vector<int> edge[2007];

int value[2007][7900],tmp[7900],all[7900];

void dfs(int x,int fa) {

int i; value[x][1]=1;*//base*

**for** (int v:edge[x]) **if** (v!=fa){

dfs(v,x);

FOR(i,1,tree\_hash::tot) tmp[i]=0;

FOR(i,1,tree\_hash::tot) **if** (value[x][i])

**for** (tree\_hash::pti k:tree\_hash::edges[i])

add\_(tmp[k.second],(ll)value[x][i]\*value[v][k.first]%M);

FOR(i,1,tree\_hash::tot) add\_(value[x][i],tmp[i]);

}

*// printf("dfs: %d %d\n",x,fa);*

*// FOR(i,1,tree\_hash::tot) if (value[x][i]) printf("%d: %d\n",i,value[x][i]);*

FOR(i,1,tree\_hash::tot) add\_(all[i],value[x][i]);

}

int main() {

tree\_hash::init(12);

tree\_hash::getall(12);

int n,i,q;

scanf("%d",&n);

FOR(i,1,n-1) {

int u,v; scanf("%d%d",&u,&v);

edge[u].push\_back(v);

edge[v].push\_back(u);

} dfs(1,0);

scanf("%d",&q);

**while** (q--) {

int m;

tree\_hash::tree tmp; tmp.init(12);

scanf("%d",&m);

vector<tree\_hash::type> V;

FOR(i,1,m-1) {

int u,v;

scanf("%d%d",&u,&v);

u--; v--;

tmp.e[u]|=1ull<<v;

tmp.e[v]|=1ull<<u;

} REP(i,m) V.push\_back(tree\_hash::id[tmp.dfs(i,i)]);

sort(V.begin(), V.end());

V.erase(unique(V.begin(), V.end()),V.end());

int ans=0;

**for** (int v:V) add\_(ans,all[v]);

printf("%d**\n**",ans);

}

}

****kruskal树****

**namespace** io {

**const** int SIZE = (1 << 21) + 1;

char ibuf[SIZE], \*iS, \*iT, obuf[SIZE], \*oS = obuf, \*oT = oS + SIZE - 1, c, qu[55]; int qr;

#define gc() (iS == iT ? (iT = (iS = ibuf) + fread (ibuf, 1, SIZE, stdin), (iS == iT ? EOF : \*iS ++)) : \*iS ++)

**inline** void flush () {

fwrite (obuf, 1, oS - obuf, stdout);

oS = obuf;

}

**inline** void putc (char x) {

\*oS ++ = x;

**if** (oS == oT) flush ();

}

**template** <**class** **I**>

**inline** void gi (I &x) {

**for** (c = gc(); c < '0' || c > '9'; c = gc());

**for** (x = 0; c <= '9' && c >= '0'; c = gc()) x = x \* 10 + (c & 15);

}

**template** <**class** **I**>

**inline** void print (I x) {

**if** (!x) putc ('0');

**while** (x) qu[++ qr] = x % 10 + '0', x /= 10;

**while** (qr) putc (qu[qr --]);

}

**struct** Flusher\_ {~Flusher\_(){flush();}}io\_flusher\_;

}

*// kruskal tree; 两点LCA为对应边上最小值*

**struct** kruskal\_tree {

**struct** edge{int u,v,w;};

vector<edge> e;

int tot;*//m:edges*

void addedge(int x,int y,int w) {

e.push\_back(edge{x,y,w});

}

int fa[maxn\*2][21],son[maxn\*2][2],\_Fa[maxn\*2],dtot;

void init(int n) {

int i; tot=n;

FOR(i,1,n) \_Fa[i]=i;

}

**inline** int getfa(int x) {

**if** (\_Fa[x]==x) **return** x;

**return** \_Fa[x]=getfa(\_Fa[x]);

}

int in[maxn\*2],out[maxn\*2],back[maxn\*2];

void dfs(int x) {

int i; in[x]=++dtot; back[dtot]=x;

rep(i,1,20) fa[x][i]=fa[fa[x][i-1]][i-1];

**if** (son[x][0]) fa[son[x][0]][0]=x,dfs(son[x][0]);

**if** (son[x][1]) fa[son[x][1]][0]=x,dfs(son[x][1]);

out[x]=dtot;

}

int val[maxn\*2];

void build(int op) {

int i;

memset(son,0,**sizeof**(son));

memset(fa,0,**sizeof**(fa));

*// sort(e+1,e+1+m,[&](edge &x,edge &y){return x.w<y.w;});*

**if** (op) reverse(e.begin(), e.end());

**for** (edge &\_e:e) {

int x=getfa(\_e.u),y=getfa(\_e.v);

**if** (x!=y) {

tot++; \_Fa[x]=\_Fa[y]=\_Fa[tot]=tot;

val[tot]=\_e.w;

son[tot][0]=x; son[tot][1]=y;

}

} dtot=0;

REP(i,20) fa[0][i]=0;

FOR(i,1,tot) **if** (getfa(i)==i) dfs(i);

}

int query\_range\_id(int x,int v) {

int i;

rREP(i,20) **if** (fa[x][i]&&val[fa[x][i]]<=v) x=fa[x][i];

**return** x;

}

}T[2];

int v[MAXM];

**inline** int lowbit(int x) {**return** x&-x;}

**inline** void update(int x,int val,int MAX) {

**for** (;x<=MAX;x+=lowbit(x)) v[x]+=val;

}

**inline** int query(int x) {

int ret=0;

**for** (;x;x-=lowbit(x)) ret+=v[x];

**return** ret;

}

int Ans[MAXM];

vector<pair<pii,int> > asks[MAXM];*//position; mask; id*

int main() {

int n,m,q; int i;

io::gi(n); io::gi(m); io::gi(q);

T[0].init(n); T[1].init(n);

FOR(i,1,m) {

int x,y,c;

io::gi(x),io::gi(y),io::gi(c); x++; y++;

**if** (!c) T[c].addedge(x,y,i);*//0：i-m+2*

**else** T[c].addedge(x,y,-i);*//1: 1-i*

} T[0].build(0); T[1].build(1);

FOR(i,1,q) {

int x,t;

io::gi(x); io::gi(t); x++;

int id0=T[0].query\_range\_id(x,t);

int id1=T[1].query\_range\_id(x,-t);

int l0=T[0].in[id0],r0=T[0].out[id0];

int l1=T[1].in[id1],r1=T[1].out[id1];

l0--; l1--;

asks[l0].push\_back(make\_pair(make\_pair(l1,1),i));

asks[l0].push\_back(make\_pair(make\_pair(r1,-1),i));

asks[r0].push\_back(make\_pair(make\_pair(l1,-1),i));

asks[r0].push\_back(make\_pair(make\_pair(r1,1),i));

}

FOR(i,1,q) Ans[i]=0;

FOR(i,1,T[0].dtot) {

int back=T[0].back[i];

**if** (back<=n) update(T[1].in[back],1,T[1].dtot);

**for** (**auto** now:asks[i])

Ans[now.second]+=query(now.first.first)\*now.first.second;

}

FOR(i,1,q) {io::print(Ans[i]); io::putc('\n');}

}

# 数学相关

牛顿迭代 开根

C = int(raw\_input())

for i in range(0, C):

n = int(raw\_input())

if n < 2 :

print n

continue

m = 2

tmpn, len = n, 0

while tmpn > 0:

tmpn /= 10

len += 1

base, digit, cur = 300, len / m, len % m

while (cur + m <= base) and (digit > 0):

cur += m

digit -= 1

div = 10 \*\* (digit \* m)

tmpn = n / div

x = int(float(tmpn) \*\* (1.0 / m))

x \*= (10 \*\* digit)

while True:

x0 = x

x = x + x \* (n - x \*\* m) / (n \* m)

if x == x0: break

while (x + 1) \*\* m <= n:

x = x + 1

print x % 2

逆元, kummer等基础

**// 计数题的时候, "选择"和"方案"分开算是一个不错的选择~**

**// 注意n>M时要用lucas!**

LL inv[1000002];*//inverse*

LL fac[1000002];*//Factorial*

**// 求出的是ax+by=1的解(a,b正负不限,而且挺小的);**

**// d(gcd)==1时存在逆元;(d!=1)&&(num|d)时,num\*a/d可认为逆元**

**// (x+p)%p为逆元**

**// DP:C[i][j]=(C[i-1][j-1]+C[i][j-1])%M**

void exgcd(LL a,LL b,LL &d,LL &x,LL &y) {

**if** (!b) {d=a; x=1; y=0;}

**else** {exgcd(b,a%b,d,y,x); y-=a/b\*x;}

}

**// 前面那个线性求逆元的log版2333**

int getinv(int n) {

**if** (n==1) **return** 1;

**return** (M-M/n)\*(getinv(M%n))%M;

}

LL C(int n,int m) {

**return** fac[n]\*inv[m]%M\*inv[n-m]%M;

}

**//Lucas扩展：Kummer定理：**

**//C(n,k)中的p的幂次的为p进制下n-k借位次数**

**//e.g.求C(n,0)...C(n,n)的lcm%(1e9+7)**

**//做法:考虑每个素因子,n转化为p进制后,除了最后的为p-1的都可以借位**

**//ans=pow(p,k)的乘积**

LL lucas(LL n,LL m) { *//注意MOD不能太大=\_=! Mlogn*

**return** m==0?1:1ll\*C(n%M,m%M)\*lucas(n/M,m/M)%M;

}

BSGS

***BSGS：a^x = b (mod p)***

***做法：假设m=sqrt(p)+1; x=i\*m-j(0<i<j)***

***枚举i和j，我们得到了一个sqrt(p)的做法***

Pell方程

打表求出第一项,然后下面的项可以线性递推

博弈：NIM,SG

**威佐夫博奕:**

**奇异局势:a=(b-a)\*(sqrt(5)+1)/2;(lose)**

**NIM博弈:k堆石子，两人轮流每次从某一堆拿走一些石子，问谁赢**

**做法:抽象(SG)->直接异或,sg=0(lose)**

**SG:选择的最多次数(连续,如1,2,3...)**

**选择的最多次数,main中为异或!=0**

int sg[maxm+2];//打表~~~

**这个是状态和剩余个数有关的**

map<int,int> Hash;

int SG(int mask){

**if** (Hash.count(mask)) **return** Hash[mask];

set<int> mex;

**for** (int i=0;i<maxm;++i){

**if** (!((mask>>i)&1)) **continue**;*//continue*

int tp=mask;

**for** (int j=i;j<maxm;j+=i+1)*//change*

**if** ((mask>>j)&1) tp^=1<<j;

mex.insert(SG(tp));*//dfs*

}

int ret=0;

**for** (;mex.count(ret);++ret);

**return** Hash[mask]=ret;

}

**这个是状态和剩余个数无关的**

map<LL,int> Hash[62];

int SG(int x,LL mask){

**if** (Hash[x].count(mask)) **return** Hash[x][mask];

set<int> mex;

**for** (int i=1;i<=x;++i){

**if** ((mask>>(i-1))&1) **continue**;*//continue*

int tp=mask;

tp^=1<<(i-1);*//change*

mex.insert(SG(x-i,tp));*//dfs*

}

int ret=0;

**for** (;mex.count(ret);++ret);

**return** Hash[x][mask]=ret;

}

Exgcd

//ax+by%x=y

int n,m;

int i,j,k;

void exgcd(LL a,LL b,LL &d,LL &x,LL &y){*//d==0时存在逆元 //(x+p)%p为逆元*

**if** (!b) {d=a;x=1;y=0;}

**else** {exgcd(b,a%b,d,y,x);y-=a/b\*x;}

}

bool check(LL a,LL b,LL x){

LL A,B,d;exgcd(a,b,d,A,B);

A\*=x;B\*=x;

LL T=A/b+B/a;

A%=b;B%=a;

**if** (A<0) A+=b,T--;

**if** (B<0) B+=a,T--;

**return** T>=0;

}

int solve(){

int a,b,x,y;

scanf("%lld%lld%lld%lld",&a,&b,&x,&y);

int g=gcd(a,b);

**if** (x%g||y%g) **return** 0\*puts("NO");

x/=g;y/=g;a/=g;b/=g;

**if** (!(x%a)&&!(y%b)) **return** 0\*puts("YES");

**if** (!(y%a)&&!(x%b)) **return** 0\*puts("YES");

**if** (!(x%(a\*b))&&check(a,b,y)) **return** 0\*puts("YES");

**if** (!(y%(a\*b))&&check(a,b,x)) **return** 0\*puts("YES");

**return** 0\*puts("NO");

}

K次方和, 伯努利数

*//sum{pow(i,k)}(1->n)*

ll B[maxn],pw[maxn];

ll A[maxn];

ll INV[10007];

LL inv[10002];*//inverse*

LL fac[10002];*//Factorial*

LL C(int n,int m) {

**return** fac[n]\*inv[m]%M\*inv[n-m]%M;

} ll SUM\_N\_K(int n,int k) {

ll pw=1,now=0; int i;

FOR(i,1,k+1) {

pw=pw\*(n+1)%M;

now+=INV[k+1]\*C(k+1,i)%M\*B[k+1-i]%M\*pw%M;

} mod\_M(now);

**return** now;

}

void initialize() {

int i,j;

fac[0]=1;

FOR(i,1,10000) fac[i]=i\*fac[i-1]%M;

inv[0]=inv[1]=1; INV[0]=INV[1]=1;

FOR(i,2,10000) INV[i]=inv[i]=(M-M/i)\*inv[M%i]%M;

FOR(i,1,10000) inv[i]=inv[i]\*inv[i-1]%M;*// inv(n!)*

B[0]=1;

FOR(i,1,2000) {

FOR(j,0,i-1) B[i]-=INV[i+1]\*C(i+1,j)%M\*B[j]%M; mod\_M(B[i]);

}

*// FOR(i,0,2000) printf("%lld ",B[i]);*

}

求原根 二次三次剩余(无板子)

**原根:存在:m=2,4,p^a,2\*p^a,p为奇质数,个数phi(phi(p-1))**

**查找:假设是g,从小枚举g**

**phi(m)=p1^a1\*p2^a2\*...\*pk^ak;**

**pow(g,phi(m)/pi)≡1恒成立(m质数则phi=m-1)**

**性质:pow(g,i)%p得到的答案两两不同**

**推论1 若d|(p-1),则x^d≡1(mod p)恰有d个解**

**推论2 若p为素数,d|(p-1),则阶为d (pow(x,d)≡1)**

**的最小剩余(mod p)的个数为phi(d)**

**二次剩余:x\*x≡n(mod p)**

**1.小的(a=0|p=2)直接判断**

**2.pow(n,(p-1)/2)≡1或-1(mod p)**

**pow(n,(p-1)/2)≡1则有解**

**3.由于1/2的数字有二次剩余**

**令w=a\*a-n;且pow(n,(p-1)/2)≡-1**

**struct A+B\*sqrt(w):**

**pow(a+sqrt(w),p)=pow(a,p)+pow(w,(p-1)/2)\*sqrt(w))**

**≡a-sqrt(w)**

**pow(a+sqrt(w),p+1)≡a\*a-w≡n**

**pow(a+sqrt(w),(p+1)/2)即为答案**

**三次剩余:x\*x\*x≡n(mod p)**

**1.小的(a=0|p=2,3)直接判断**

**2.p≡-1(mod 3):x≡pow(a,(2\*p-1)/3)**

**3.p≡1(mod 3):设e为三次单位根,e\*e\*e≡1(mod p)**

**pow(a,(p-1)/3)≡1(mod p)则有三次剩余**

int p[maxn],tot;

bool mark[maxn];

bool isroot(int x,int p){

**if** (!(x%p)||(x%p==1&&p!=2)) **return** 0;

**for** (ll i=2;i\*i<=p-1;i++) **if** ((p-1)%i==0)

**if** (poww(x,(p-1)/i,p)==1||poww(x,i,p)==1) **return** 0;

**return** 1;

}

int TaskA() {

int i,x;

scanf("%d%d",&n,&x);

**if** (mark[n+1]) **return** 0\*puts("-1");

rFOR(i,2,x-1){

**if** (!isroot(i,n+1)) **continue**;

**return** 0\*printf("%d**\n**",i);

} **return** 0\*puts("-1");

}

void initialize() {

int i,j;

FOR(i,2,5000001) {

**if** (!mark[i]) p[tot++]=i;

REP(j,tot) {

**if** (i\*p[j]>5000001) **break**;

mark[i\*p[j]]=1;

**if** (i%p[j]==0) **break**;

}

}

}

常系数线性递推

**M^2logn的普通版本:**

int ini[3007];

**const** int mod=998244352;

void mul(ll \*a,ll \*b,int k) {

int i,j;

**static** ll tmp[3007];

REP(i,k+k) tmp[i]=0;

REP(i,k) **if** (a[i]) REP(j,k)

((tmp[i+j]+=a[i]\*b[j])>=INFF)&&(tmp[i+j]%=mod);

rrep(i,k,k+k) **if** (tmp[i]){

tmp[i]%=mod;

REP(j,k) ((tmp[i-k+j]+=tmp[i]\*ini[k-j])>=INFF)&&(tmp[i-k+j]%=mod);

}

REP(i,k) a[i]=tmp[i]%mod;

}

ll A[3007],B[3007];

void power(int k,ll n) {

**if** (k!=1) A[1]=1; **else** A[0]=ini[1]; B[0]=1;

**for** (ll x=n; x; x>>=1) {

**if** (x&1) mul(B,A,k);

mul(A,A,k);

}

}

**Mlognlogm的fft版本:前面的fft板子后面有的**

ll A[maxn],B[maxn];

ll C[maxn],D[maxn],E[maxn];

int main() {

ll l,r;

scanf("%d%lld%lld",&n,&l,&r); int i;

FOR(i,1,n) scanf("%lld",&A[i]),A[i]=B[i]=(M-A[i]%M);

A[0]=1; **while** (A[n]==0) n--; m=n;

n+=r-l; n++;

NTT::inverse(A,A,n\*2);

m++; B[0]=1; reverse(B,B+m);*//no!*

C[1]=1; D[0]=1; int lC=2,lD=1;

**for** (ll x=l; x; x>>=1) {

**if** (x&1) {

NTT::multiply(C,D,D,lC,lD); int l1;

NTT::delivery(D,B,E,D,lC+lD-1,m,l1,lD);

} NTT::multiply(C,C,C,lC,lC); int l1;

NTT::delivery(C,B,E,C,lC+lC-1,m,l1,lC);

} reverse(D,D+lD);

NTT::multiply(D,A,E,lD,n);

FOR(i,lD-1,r-l+lD-1) printf("%lld**\n**",E[i]);

**return** 0;

}

多项式暴力求积分

**// 题意: 给n个区间, 求和与0的距离差期望**

**// 做法: 考虑每一次, 都是个区间分段积分形式**

**// 所以直接考虑每一段, 2-pointer求2^n项系数即可**

**const** int maxk=20;

int inv[maxk];

**inline** void init() {

int i; inv[0]=1;

rep(i,1,maxk) inv[i]=powMM(i,M-2);

}

**struct** poly {

int A[maxk],n;*//base*

void init(int n) {memset(A,0,**sizeof**(A)); **this**->n=n;};

poly(int n=0) {init(n);}

int getvalue(int x) {

int i; ll ret=0;

rREP(i,n+1) ret=(ret\*x+A[i])%M;

**return** ret;

}

poly integral() {*//积分; ret.A[0]需要自己算*

poly ret; int i; ret.n=n+1; ret.A[0]=0;

REP(i,n+1) ret.A[i+1]=(ll)A[i]\*inv[i+1]%M;

**return** ret;

}

poly derivative(){*//求导*

poly ret; int i; ret.n=n-1; ret.A[0]=0;

REP(i,n) ret.A[i]=(ll)A[i+1]\*(i+1)%M;

**return** ret;

}

poly move(poly base[]) {

poly ret; ret.init(n); int i,j;

REP(i,n+1) REP(j,i+1) {

add\_(ret.A[j],(ll)A[i]\*base[i].A[j]%M);

} **return** ret;

}

poly mul(**const** poly &p) {

poly ret; ret.init(n+p.n); int i,j;

REP(i,n+1) REP(j,p.n+1) {

add\_(ret.A[i+j],(ll)A[i]\*p.A[j]%M);

} **return** ret;

}

poly del(**const** poly &p) {

poly ret; ret.init(max(n,p.n)); int i;

REP(i,n+1) add\_(ret.A[i],A[i]);

REP(i,p.n+1) add\_(ret.A[i],M-p.A[i]);

**return** ret;

}

}; *//(x-r)*

int len[maxn];

**typedef** pair<int,poly> pip;*// 从first往后一段区间内的poly\_value是second*

int base=0,multi=1;

void getintegral(pip now[],pip nxt[],int n) {*//得到一个连续的积分,并从0开始算常数项*

int i;

FOR(i,1,n) {

int k=nxt[i].first=now[i].first;

nxt[i].second=now[i].second.integral();

nxt[i].second.A[0]=(nxt[i-1].second.getvalue(k)

-nxt[i].second.getvalue(k)+M\*2)%M;

}

}

void getbase(poly base[],int k,int n) {*//(x-k)^n; 用于移动整个区间*

int i; base[0].init(0);

base[0].A[0]=1;*//1*

poly mul; mul.init(1);

mul.A[0]=(M-k%M)%M; mul.A[1]=1;

FOR(i,1,n) base[i]=base[i-1].mul(mul);

}

pip now[1<<15|7],inter[1<<15|7],nxt[1<<15|7];

poly multibase[maxk];

int ans=0; int mbase=1;

int main() {

init();

int n,i;

scanf("%d",&n);

FOR(i,1,n) {

int l,r;

scanf("%d%d",&l,&r);

add\_(ans,(l+r)\*powMM(2ll,M-2)%M);*//之后加两倍负数即可*

r-=l; base-=l; len[i]=r;

**if** (r) mbase=mbase\*powMM(r,M-2)%M;

}

reverse(len+1,len+1+n);

now[0].second.init(0);

now[0].second.A[0]=0;

now[1].first=0;

now[1].second.init(0);

now[1].second.A[0]=1;*//integeal\_ed*

int m=1;*//Count*

FOR(i,1,n) {

**if** (len[i]==0) **continue**;

getintegral(now,inter,m);

*//开始积分*

getbase(multibase,len[i],n);

int l=1,r=1,nxtm=0;

**while** (l<=m||r<=m) {

int pos=INF;

**if** (l<=m) pos=min(pos,now[l].first+len[i]);

**if** (r<=m) pos=min(pos,now[r].first);

**while** (l<=m&&now[l].first+len[i]==pos) l++;*//相同pos只有一次*

**while** (r<=m&&now[r].first==pos) r++;

++nxtm; nxt[nxtm].first=pos;

nxt[nxtm].second=inter[r-1].second.del(inter[l-1].second.move(multibase));

}

swap(now,nxt); m=nxtm;

} poly x; x.init(1); x.A[1]=1;

getbase(multibase,-base,n);

FOR(i,1,m) {

now[i].first-=base;

now[i].second=now[i].second.move(multibase);

now[i].second=now[i].second.derivative();

now[i].second=now[i].second.mul(x);

} getintegral(now,inter,m);

int last=0;

FOR(i,1,m) **if** (inter[i].first<=0) last=i;

add\_(ans,(2ll\*(M-(ll)inter[last].second.getvalue(0)\*mbase%M))%M);

printf("%d**\n**",ans);

}

五边形数定理

**/\*hdu4651**

**题意：普通的整数拆分**

**限制：1 <= n <= 1e5**

**思路：五边形数定理**

**Q(x)=\mul(1-x^k) = 1-x-x^2+x^5+x^7+...**

**Q(x)=\sum\_k(-1)^k x^(k\*(3k-1)/2)**

**\*/**

#include <iostream>

#include <cstdio>

**using** **namespace** std;

#define LL \_\_int64

**const** int N=100005;

**const** int MOD=1000000007;

LL dp[N],fi[N];

LL five(LL x){ **return** (3\*x\*x-x)/2; }

*//五边形数*

void wbxs(){

dp[0]=1;

int t=1000; *//其实可以等于sqrt(N)base*

**for**(int i=-t;i<=t;++i)

fi[i+t]=five(i); *//Q*

**for**(int i=1;i<=10000;++i){

int flag=1;

**for**(int j=1;;++j){

LL a=fi[j+t],b=fi[-j+t];

**if**(a>i && b>i) **break**;

**if**(a<=i) dp[i]=(dp[i]+dp[i-a]\*flag+MOD)%MOD; *//p*

**if**(b<=i) dp[i]=(dp[i]+dp[i-b]\*flag+MOD)%MOD;

flag\*=-1;

}

}

}

FFT

**DFT式子: x\_k=\sum{x\_i\*wn[k\*i]};**

**namespace** FFT {

**const** int maxn=1<<18|7;

**struct** complex {

double a,b;

complex(double \_a=.0,double \_b=.0):a(\_a),b(\_b) {}

complex **operator**+(**const** complex x)**const** {**return** complex(a+x.a,b+x.b);}

complex **operator**-(**const** complex x)**const** {**return** complex(a-x.a,b-x.b);}

complex **operator**\*(**const** complex x)**const** {**return** complex(a\*x.a-b\*x.b,a\*x.b+b\*x.a);}

};

complex wn[maxn];

void initwn(int l) {

**static** int len=0; int i;

**if** (len==l) **return**; **else** len=l;

REP(i,len) wn[i]=complex(cos(2\*pi\*i/l),sin(2\*pi\*i/l));

}

void fft(complex \*A,int len,int inv) {

int i,j,k; initwn(len);

**for** (i=1,j=len/2; i<len-1; i++) {

**if** (i<j) swap(A[i],A[j]);

k=len/2;

**while** (j>=k) j-=k,k/=2;

**if** (j<k) j+=k;

} **for** (i=2; i<=len; i<<=1) {

**for** (j=0; j<len; j+=i) {

**for** (k=j; k<(j+i/2); k++) {

complex a,b; a=A[k];

b=A[k+i/2]\*wn[(ll)(k-j)\*len/i];

A[k]=a+b; A[k+i/2]=a-b;

}

}

} **if** (inv==-1) REP(i,len) A[i]=complex(A[i].a/len,A[i].b/len);

}

**inline** complex conj(complex &A) {**return** complex(A.a,-A.b);}

void mul(int \*A,int \*B,int \*ans,int len) { *//ans=A\*B*

**static** complex x1[maxn],x2[maxn]; int i;

REP(i,len) x1[i]=complex(A[i],B[i]);

fft(x1,len,1);

REP(i,len) {*//这个k1, b1就是前面的, 这就减掉了一半常数*

int j=(len-i)&(len-1);

complex a=(conj(x1[i])+x1[j])\*complex(0.5,0);*//dft a*

complex b=(conj(x1[i])-x1[j])\*complex(0,0.5);*//dft b*

x2[i]=a\*b;

} fft(x2,len,-1);

REP(i,len) ans[i]=x2[i].a+0.5;

}

}

多项式单log无穷背包(MTT)

**// 主要思路不是这个裸的乘法啥的啊!**

**// from picks' blog**

**// 对G(F(x))=0进行泰勒展开**

**// G'(F\_{t+1}(x))=G(F\_t(x))+G'(F\_t(x))/1\*(F\_{t+1}-F\_t(x))^1+....**

**// 后方的系数在mod x^2^t+1的意义下全是0!(因为减的那里的系数是2^t)**

**// F\_{t+1}(x)=F\_t(x)-G(F\_t(x))/G'(F\_t(x))**

**// 所以手动求个导数即可!**

**// 注意这个G(F(t))就是满足的那个式子! 注意要有常数项(否则可以全是0 =\_=!)**

**// 三角函数需要利用虚数来做, e^{iF(x)}=cos(F(x))+isin(F(x))**

**// exp(x): F\_{t+1}(x)=F\_t(x)-F\_t(x)\*((ln(F\_t(x))-P(x))\*F\_t(x))**

**// ln(x) : ln(F(x))=\int(积分) F'(x)/F(x)**

**// 注意F[0]要是0, 因为求导的时候会去掉这个贡献, 积分回来**

**// 这个是无穷背包**

**namespace** FFT {

**const** int maxn=1<<18|7;

**struct** complex {

double a,b;

complex(double \_a=.0,double \_b=.0):a(\_a),b(\_b) {}

complex **operator**+(**const** complex x)**const** {**return** complex(a+x.a,b+x.b);}

complex **operator**-(**const** complex x)**const** {**return** complex(a-x.a,b-x.b);}

complex **operator**\*(**const** complex x)**const** {**return** complex(a\*x.a-b\*x.b,a\*x.b+b\*x.a);}

};

complex wn[maxn];

void initwn(int l) {

**static** int len=0; int i;

**if** (len==l) **return**; **else** len=l;

REP(i,len) wn[i]=complex(cos(2\*pi\*i/l),sin(2\*pi\*i/l));

}

void fft(complex \*A,int len,int inv) {

int i,j,k; initwn(len);

**for** (i=1,j=len/2; i<len-1; i++) {

**if** (i<j) swap(A[i],A[j]);

k=len/2;

**while** (j>=k) j-=k,k/=2;

**if** (j<k) j+=k;

} **for** (i=2; i<=len; i<<=1) {

**for** (j=0; j<len; j+=i) {

**for** (k=j; k<(j+i/2); k++) {

complex a,b; a=A[k];

b=A[k+i/2]\*wn[(ll)(k-j)\*len/i];

A[k]=a+b; A[k+i/2]=a-b;

}

}

} **if** (inv==-1) REP(i,len) A[i]=complex(A[i].a/len,A[i].b/len);

}

**inline** complex conj(complex &A) {**return** complex(A.a,-A.b);}

void mul(int \*A,int \*B,int \*ans,int len,int mod) { *//ans=A\*B*

**static** complex x1[maxn],x2[maxn];

**static** complex x3[maxn],x4[maxn];

**static** **const** int S=1<<15 ; int i;

REP(i,len) x1[i]=complex(A[i]/S,A[i]%S);

REP(i,len) x2[i]=complex(B[i]/S,B[i]%S);

fft(x1,len,1); fft(x2,len,1);

REP(i,len) {*//这个k1, b1就是前面的, 这就减掉了一半常数*

int j=(len-i)&(len-1);

complex k1=(conj(x1[i])+x1[j])\*complex(0.5,0);*//dft k1*

complex b1=(conj(x1[i])-x1[j])\*complex(0,0.5);*//dft b1*

complex k2=(conj(x2[i])+x2[j])\*complex(0.5,0);*//dft k2*

complex b2=(conj(x2[i])-x2[j])\*complex(0,0.5);*//dft b2*

x3[i]=k1\*k2+k1\*b2\*complex(0,1);

x4[i]=b1\*k2+b1\*b2\*complex(0,1);

} fft(x3,len,-1); fft(x4,len,-1);

REP(i,len) {

ll kk=x3[i].a+0.5,kb=x3[i].b+0.5;

ll bk=x4[i].a+0.5,bb=x4[i].b+0.5;

ans[i]=((kk%mod\*S%mod+kb+bk)%mod\*S%mod+bb)%mod;

}

}

**const** ll Mod=19260817;

*// 下方的东西和ntt就根本无关, 这个模数是可以改的, 是多项式相关的东西*

*// 也就是说, 这个模数完全可以取其他的, 然后高精度的mtt来求, 不过可能会T到死*

int eInv[maxn];

void initinv(int l) {

int i; eInv[0]=eInv[1]=1;

rep(i,2,l) eInv[i]=(Mod-Mod/i)\*eInv[Mod%i]%Mod;

}

void Ftof(int \*A,int \*B,int l) {*//derivative 求导*

int i;

FOR(i,1,l) B[i-1]=(ll)A[i]\*i%Mod;

}

void ftoF(int \*A,int \*B,int l) {*//integral 积分*

int i; *// todo:get B[0], getinv*

rFOR(i,1,l) B[i]=(ll)A[i-1]\*eInv[i]%Mod;

B[0]=0;

}

void inv(int \*A,int \*B,int l) { *//B=inv(A)*

**static** int C[maxn],D[maxn];

B[0]=eInv[A[0]]; B[1]=0;

**for** (int len=2; len<=l; len<<=1) {

int i; fill(B+len,B+len+len,0);

copy(A,A+len,C); fill(C+len,C+len+len,0);

mul(C,B,D,len\*2,Mod); fill(D+len,D+len+len,0);

mul(D,B,D,len\*2,Mod);

REP(i,len) B[i]=(B[i]\*2-D[i]+Mod)%Mod;

fill(B+len,B+len+len,0);

}

}

void ln(int \*A,int \*B,int l) {

**static** int C[maxn];

inv(A,B,l); Ftof(A,C,l);

mul(B,C,B,l\*2,Mod);

ftoF(B,B,l);

}

void exp(int \*A,int \*B,int l) {

**static** int C[maxn],i;

B[0]=1; B[1]=0;

**for** (int len=2; len<=l; len<<=1) {

fill(B+len,B+len+len,0);

ln(B,C,len); fill(C+len,C+len+len,0);

REP(i,len) C[i]=(C[i]-A[i]+Mod)%Mod;

mul(B,C,C,len\*2,Mod);

REP(i,len) B[i]=(B[i]-C[i]+Mod)%Mod;

}

}

*//这里是更高一层的东西*

**static** int A[maxn],B[maxn];

void multiply(int \*a,int \*b,int \*ans,int n,int m) {*//C=A\*B(actual)*

int len=1,i;

**while** (len<n+m-1) len<<=1;

REP(i,n) A[i]=a[i]; rep(i,n,len) A[i]=0;

REP(i,m) B[i]=b[i]; rep(i,m,len) B[i]=0;

mul(A,B,ans,len,Mod);

}

void getexp(int \*a,int \*ans,int n) {

int len=1,i;

**while** (len<n) len<<=1;

REP(i,n) A[i]=a[i]; rep(i,n,len) A[i]=0;

exp(A,ans,len);

}

void solve(int \*a,int \*ans,int m) {

**static** int A[maxn];

int i,j;

FOR(i,1,m) {*//无穷背包*

int now=(ll)i\*a[i]%Mod;

**for** (j=i-1; j<=m; j+=i) A[j]=(now+A[j])%Mod;

} ftoF(A,A,m);

getexp(A,ans,m+1);

}

}

int A[maxn],ans[maxn];

int main() {

int i,k;

FFT::initinv(maxn);

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%d",&k),A[k]++;

FFT::solve(A,ans,m);

*// FOR(i,1,m) printf("%d ",ans[i]);*

int Ans=0;

FOR(i,1,m) add\_(Ans,ans[i],FFT::Mod);

printf("%d**\n**",Ans);

}

多项式开根求逆,除法取模(NTT)

**// http://codeforces.com/contest/438/problem/E**

**// 题意: 问你有多少个二叉树点权从c中取, 而且权值和是k**

**// 做法: 考虑多一个点, 所以f[x]=sigma{f[k]\*f[x-k-s],(s in c)}**

**// 所以 满足 F=F^2\*C+1, 左边是生成函数**

**// 所以 F=[1-sqrt(1-4C)]/2C=1/(1+sqrt(1-4C))**

**// 当且仅当常数项有逆元, 可以多项式求逆**

**// 求逆:C\*A≡1(mod x^n)**

**// B\*A≡1(mod x^(n/2))**

**// (B\*A-1)\*(B\*A-1)≡0(mod x^(n/2))**

**// B\*B\*A\*A-2\*A\*B+1≡0(mod x^n)**

**// B\*B\*A-2\*B+C≡0(mod x^n)**

**// C≡B\*(2-A\*B)(mod x^n)**

**// 求根:C\*C≡A(mod x^n)**

**// B\*B≡A(mod x^n/2)**

**// (B\*B-A)\*(B\*B-A)≡0(mod x^n)**

**// B\*B\*B\*B-2\*C\*C\*B\*B+C\*C\*C\*C≡0(mod x^n)**

**// (B\*B+C\*C)\*(B\*B+C\*C)≡4\*C\*C\*B\*B(mod x^n)**

**// B\*B+A≡2\*C\*B(mod x^n)**

**// C=(B\*B+A)/(2\*B)**

**namespace** NTT {

**const** int maxn=1<<20|7;

**const** ll MOD=998244353;

**const** ll g=3;

int wn[maxn],invwn[maxn];

ll mul(ll x,ll y) {

**return** x\*y%MOD;

}

ll poww(ll a,ll b) {

ll ret=1;

**for** (; b; b>>=1ll,a=mul(a,a))

**if** (b&1) ret=mul(ret,a);

**return** ret;

}

void initwn(int l) {

**static** int len=0;

**if** (len==l) **return**; len=l;

ll w=poww(g,(MOD-1)/len); int i;

ll invw=poww(w,MOD-2); wn[0]=invwn[0]=1;

rep(i,1,len) {

wn[i]=mul(wn[i-1],w);

invwn[i]=mul(invw,invwn[i-1]);

}

}

void ntt(ll \*A,int len,int inv) {

int i,j,k; initwn(len);

**for** (i=1,j=len/2; i<len-1; i++) {

**if** (i<j) swap(A[i],A[j]);

k=len/2;

**while** (j>=k) j-=k,k/=2;

**if** (j<k) j+=k;

} **for** (i=2; i<=len; i<<=1) {

**for** (j=0; j<len; j+=i) {

**for** (k=j; k<(j+i/2); k++) {

ll a,b; a=A[k];

**if** (inv==-1) b=mul(A[k+i/2],invwn[(ll)(k-j)\*len/i]);

**else** b=mul(A[k+i/2],wn[(ll)(k-j)\*len/i]);

A[k]=(a+b); (A[k]>=MOD) &&(A[k]-=MOD);

A[k+i/2]=(a-b+MOD); (A[k+i/2]>=MOD) &&(A[k+i/2]-=MOD);

}

}

} **if** (inv==-1) {

ll vn=poww(len,MOD-2);

REP(i,len) A[i]=mul(A[i],vn);

}

}

void mul(ll \*A,ll \*B,ll \*C,int len) { *//C=A\*B*

int i;

ntt(A,len,1); ntt(B,len,1);

REP(i,len) C[i]=mul(A[i],B[i]);

ntt(C,len,-1);

}

void inv(ll \*A,ll \*B,int l) { *//B=inv(A)*

**static** ll C[maxn];

B[0]=poww(A[0],MOD-2); B[1]=0;

**for** (int len=2; len<=l; len<<=1) {

int i; fill(B+len,B+len+len,0);

copy(A,A+len,C); fill(C+len,C+len+len,0);

ntt(C,len\*2,1); ntt(B,len\*2,1);

REP(i,len\*2) B[i]=mul(B[i],(MOD+2-mul(C[i],B[i])));

ntt(B,len\*2,-1); fill(B+len,B+len+len,0);

}

}

void sqrt(ll \*A,ll \*B,int l) { *//B=sqrt(A)*

**static** ll C[maxn],\_B[maxn];

B[0]=1; B[1]=0;*// 这里应该是个二次剩余*

**for** (int len=2; len<=l; len<<=1) {

int i; ll inv2=poww(2,MOD-2);

inv(B,\_B,len); fill(B+len,B+len+len,0);

copy(A,A+len,C); fill(C+len,C+len+len,0);

ntt(C,len\*2,1); ntt(\_B,len\*2,1); ntt(B,len\*2,1);

REP(i,len\*2) B[i]=mul(inv2,B[i]+mul(C[i],\_B[i]));

ntt(B,len\*2,-1); fill(B+len,B+len+len,0);

}

}

**static** ll A[maxn],B[maxn];

void multiply(ll \*a,ll \*b,ll \*ans,int n,int m) {*//C=A\*B(actual)*

int len=1,i;

**while** (len<n+m-1) len<<=1;

REP(i,n) A[i]=a[i]; rep(i,n,len) A[i]=0;

REP(i,m) B[i]=b[i]; rep(i,m,len) B[i]=0;

mul(A,B,ans,len);

}

void inverse(ll \*a,ll \*ans,int n){

int len=1,i;

**while** (len<n) len<<=1;

REP(i,n) A[i]=a[i]; rep(i,n,len) A[i]=0;

inv(A,ans,len);

}

void getsqrt(ll \*a,ll \*ans,int n){

int len=1,i;

**while** (len<n) len<<=1;

REP(i,n) A[i]=a[i]; rep(i,n,len) A[i]=0;

sqrt(A,ans,len);

}

void divide(ll \*a,ll \*b,ll \*ans,int n,int m,int &l) {

**if** (n<m) {l=1; ans[0]=0; **return**;}

int len=1,i; l=n-m+1;

**while** (len<n-m+1) len<<=1;

REP(i,n) A[i]=a[i]; reverse(A,A+n); min\_(n,l);

REP(i,m) B[i]=b[i]; reverse(B,B+m); min\_(m,l);

rep(i,m,len) B[i]=0;

inv(B,ans,len);

multiply(A,ans,ans,len,n);

reverse(ans,ans+l);

}

*//ans1:答案; ans2:余数*

void delivery(ll \*a,ll \*b,ll \*ans1,ll \*ans2,int n,int m,int &l1,int &l2) {

divide(a,b,ans1,n,m,l1); l2=m-1;

**static** ll tmp[maxn];

multiply(b,ans1,tmp,m,l1);

int i; REP(i,l2) ans2[i]=(a[i]-tmp[i]+M)%M;

}

}

ll A[maxn],ans[maxn];

int main() {

int i,k;

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%d",&k),A[k]++;

REP(i,m+1) A[i]=-4\*A[i]; A[0]++;

REP(i,m+1) mod\_(A[i]);

NTT::getsqrt(A,ans,m+1);

add\_(ans[0],1);

NTT::inverse(ans,ans,m+1);

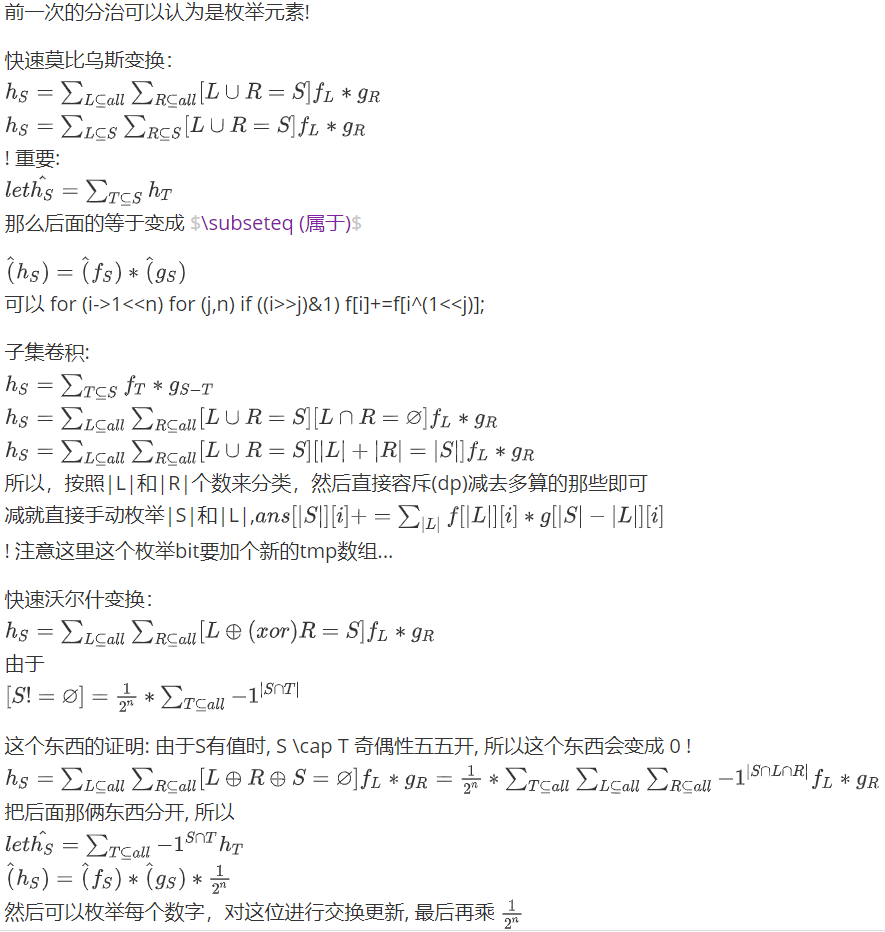
FOR(i,1,m) mul\_(ans[i],2);

FOR(i,1,m) printf("%lld**\n**",ans[i]);

}

fwt, 子集卷积

**按数学理解:**



**按位理解:**

**//or/and的理解:这里的变换是利用dp时分治来压位(写成非递归形式)实现的，时间nlogn**

**//进行组合可以将二元运算的东西都组合出来**

**//实际上or都没用**

void fwt(LL \*A,int len,int inv)*//对拍对了*

{

int i,j,k;

int div=powMM(2ll,M-2);

**for**(i=2;i<=len;i<<=1){

**for**(j=0;j<len;j+=i){

**for**(k=j;k<j+i/2;k++){

**if** (inv==1){

LL a=A[k],b=A[k+i/2];

A[k]=(a+b)%M;

A[k+i/2]=(a-b+M)%M;

*//xor:a[k]=x+y,a[k+i/2]=(x-y+mod)%mod;*

*//and:a[k]=x+y;*

*//or:a[k+i/2]=x+y;*

}**else**{

LL a=A[k],b=A[k+i/2];

A[k]=(a+b)\*div%M;

A[k+i/2]=(a-b+M)%M\*div%M;

*//xor:a[k]=(x+y)/2,a[k+i/2]=(x-y)/2;*

*//and:a[k]=x-y;*

*//or:a[k+i/2]=y-x;*

}

}

}

}

}

子集和(高维前缀和, 分治)

**大概做法是按照每一位来分类，然后往下递归获得答案**

**就是，按照这一位是1和0分成几类往下递归**

**对子集分治和高次前缀和的理解:**

**子集分治:**

**考虑维护f(x)=真正的值;**

**g(x)=子集中的贡献值(这个g可以是f进行fft得到的,也可以是与位有关的,较容易得到)**

**solve(x)的过程:**

**solve(left); 加f贡献到g;**

**solve(right); 在f=g固定right结果; 若g需要更改则更改g(如子集和就是普通的右边加左边)**

**高维前缀和:**

**高维前缀和的方式就是从小往大枚举每一位,然后一个一个维度做前缀和，从小往大for value**

**然后对于每个value,枚举前面的cnt次数,保存枚举的位数,分别前缀和**

**// f(x)=sum\_(f(y)), f(x)在某些条件下=0**

**分治做法：**

int f[1<<21|7],g[1<<21|7];

char str[1<<21|7];

void solve(int l,int r) {

**if** (l+1==r) {

**if** (!l) f[l]=1,g[l]=0;

**else** **if** (str[l-1]=='+') f[l]=0;

add\_(g[l],f[l]);

**return**;

} int i,mid=(l+r)/2;

solve(l,mid);

REP(i,mid-l) add\_(f[i+mid],g[i+l]);*//left的影响*

solve(mid,r);

REP(i,mid-l) add\_(g[i+mid],g[i+l]);

}

**高维前缀和做法：**

char str[1<<21|7];

int f[1<<21|7],g[1<<21|7][22];*//sum\_of\_previous*

int main() {

int T,\_; T=1;

scanf("%d",&T);

FOR(\_,1,T){

*/\*to solve the problem\*/*

scanf("%s",str);

int n=strlen(str),sta,i;

memset(f,0,(n+2)\***sizeof**(int));

int MAX=1;

**while** ((1<<MAX)<=n+1) MAX++;

f[0]=1; REP(i,MAX) g[0][i]=1;

FOR(sta,1,n+1) {

**if**(((n+1)&sta)==sta) {

REP(i,MAX) {

g[sta][i]=i?g[sta][i-1]:0;

**if** ((sta>>i)&1) add\_(g[sta][i],g[sta^(1<<i)][i]);

}

**if** (str[sta-1]!='+') {

f[sta]=g[sta][MAX-1];*//front*

REP(i,MAX) add\_(g[sta][i],f[sta]);

}

}

}

*// REP(i,n+3) printf("%d ",f[i]); puts("<- f");*

*// REP(i,n+3) printf("%d ",g[i][21]); puts("<- g");*

printf("%d %d**\n**",n+1,f[n+1]);

}

}

子集卷积

**//第一种做法: 按位考虑**

**//大概做法是按照每一位来分类，然后往下递归获得答案**

**就是，按照这一位是1和0分成几类往下递归**

**//http://acm.hdu.edu.cn/showproblem.php?pid=6057**

**//很容易卡T...3^18也许能过**

**//这个比2^nlog^2(n=19)的慢了快5倍**

**//这种思路这种题都能用**

**//最好像tls那样推一推然后写成非递归, 常数会减少到和2^n\*n^2差不多**

**//真\*子集卷积by TLS(卡常 )**

const int maxn = 1 << 19 | 1, mod = 998244353, seed = 1526;

int n, all, bit[maxn], a[maxn], b[maxn], ans;

inline void mod\_inc(int &x, int y) {

(x += y) >= mod && (x -= mod);

}

int main() {

while(scanf("%d", &n) == 1) {

all = (1 << n) - 1;

for(int i = 0; i <= all; ++i) scanf("%d", a + i);

for(int i = 0; i <= all; ++i) scanf("%d", b + i);

bit[0] = 1;

for(int i = 1; i <= all; ++i) {

bit[i] = bit[i >> 1] << (i & 1);

a[i] = (LL)a[i] \* bit[i] % mod;

}

ans = 0;

for(int i = all; i >= 0; --i) {

int msk = all ^ i, tim = 0;

ULL cnt = 0;

for(int j = msk; j; j = (j - 1) & msk) {

cnt += (ULL)a[j] \* b[i | j];

(++tim) == 18 && (tim = 0, cnt %= mod);

}

cnt += (ULL)a[0] \* b[i]; cnt %= mod;

ans = ((LL)seed \* ans + cnt) % mod;

}

printf("%d\n", ans);

}

return 0;

}

**//第二种做法: 数学方法**

**证明在上面小字**

**// C[k]=\sum\_{i&j=k} A[i^j] B[i|j]**

**// let i'=i^j, j'=i|j, 这样的i,j对有2^bit(i')个**

**// C[k]=\sum [j'-i'=k] [i' \subseteq j'] 2^i' \* A[i'] \* B[j']**

**// C[k]=\sum [i^j=k] [i&j=i] \* 2^i \* A[i] \* B[j] //这里的意思就是i|k=j, i+k=j**

**// C[k]=\sum [i^j=k] [bit(j)-bit(i)=bit(k)] 2^i \* A[i] \* B[j]**

**// ! 注意这里这个枚举bit要加个新的tmp数组...**

int A[20][maxn],B[20][maxn],C[20][maxn];

int bit[maxn],pw1526[maxn],ans[maxn];

int main() {

int k,i;

scanf("%d",&m); n=1<<m;

REP(i,n) bit[i]=bit[i>>1]+(i&1);

REP(i,n) scanf("%d",&A[bit[i]][i]);

REP(i,n) scanf("%d",&B[bit[i]][i]);

pw1526[0]=1;

rep(i,1,n) pw1526[i]=1526ll\*pw1526[i-1]%M;

REP(i,n) mul\_(A[bit[i]][i],powMM(2,bit[i]));

REP(i,m+1) fwt(A[i],n,1),fwt(B[i],n,1);

REP(k,m+1) REP(i,m-k+1) {

int t=0,j=i+k;

REP(t,n) add\_(C[k][t],(ll)A[i][t]\*B[j][t]%M);

} REP(i,m+1) fwt(C[i],n,-1);

REP(i,n) ans[i]=C[bit[i]][i];

int Ans=0;

REP(i,n) add\_(Ans,(ll)ans[i]\*pw1526[i]%M);

printf("%d**\n**",Ans);

}

高斯消元

**indström–Gessel–Viennot lemma:**

**使用这个定理必须是平面图;**

**这个是个求不相交路径对数的方案数的定理**

**答案是：下列矩阵行列式，其中A[i,j]表示i到j方案数**

**|A[1,1],A[1,2]|**

**|A[2,1],A[2,2]|**

**//求行列式的值**

**//%m,m为质数的积**

**//从0开始**

**template**<**typename** T>**inline** T poww(T a,T b,T M) {

T ret=1;

**for** (; b; b>>=1ll,a=1ll\*a\*a%M)

**if** (b&1) ret=1ll\*ret\*a%M;

**return** ret;

}

LL guass(LL A[107][107],int n,LL M) {

LL ret=1; int i,j,k;

REP(i,n) {

int id=i;

**if** (!A[i][i]) rep(j,i+1,n) **if** (A[j][i]) id=j;

**if** (!A[id][i]) **continue**;

**if** (id!=i) {rep(j,i,n) swap(A[i][j],A[id][j]); ret\*=-1;}

A[i][i]%=M; (A[i][i]<0) &&(A[i][i]+=M);

LL rev=poww(A[i][i],M-2,M);

rep(k,i+1,n)

rrep(j,i,n)(A[k][j]-=(LL)A[k][i]\*rev%M\*A[i][j])%=M;

} REP(i,n)(ret\*=A[i][i])%=M;

(ret<0) &&(ret+=M);

**return** ret;

}

LL A[107][107],B[107][107];

void exgcd(LL a,LL b,LL &d,LL &x,LL &y) {

**if** (!b) {d=a; x=1; y=0;}

**else** {exgcd(b,a%b,d,y,x); y-=a/b\*x;}

}

vector<LL> P;

vector<LL> Ans;

LL ans;

LL chinese\_remainder(vector<LL> &m,vector<LL> &r) {

int i; LL M=m[0],R=r[0];

rep(i,1,P.size()) {

LL x,y,d;

exgcd(M,m[i],d,x,y);

**if** ((r[i]-R)%d) **return** -1;

x=(r[i]-R)/d\*x%(m[i]/d);

R+=x\*M; M=M/d\*m[i];

R%=M; (R<0) &&(R+=M);

} **return** R;

}

int n,m;

int i,j,k;

int main() {

**while** (~scanf("%d%d",&n,&m)) {

P.clear(); Ans.clear();

REP(i,n)

REP(j,n) scanf("%lld",&A[i][j]);

**for** (i=2; i\*i<=m; i++) **if** (m%i==0) {

P.push\_back(i);

**while** (m%i==0) m/=i;

} **if** (m!=1) P.push\_back(m);

**for** (int v:P) {

REP(i,n) REP(j,n) B[i][j]=A[i][j];

Ans.push\_back((LL)guass(B,n,v));

}

ans=chinese\_remainder(P,Ans);

printf("%lld**\n**",ans);

}

}

**//辗转相除法**

REP(i,n) {

rep(j,i+1,n) {

int x=i,y=j;

**while** (a[y][i]) {

LL t=a[x][i]/a[y][i];

rep(k,i,n) a[x][k]=(a[x][k]-a[y][k]\*t)%m;

swap(x,y);

}

**if** (x!=i) {

rep(k,i,n) swap(a[i][k],a[x][k]);

ans=(-ans+m)%m;

}

}

ans=ans\*a[i][i]%m;

ans=(ans+m)%m;

}

矩阵树定理|拉格朗日插值

**// 题意:求生成树中含k条给定树边的生成树个数**

**// 做法:为给定边加不同权值,然后矩阵树定理**

**// 矩阵树定理:生成树数量=|基尔霍夫矩阵C=D-A|;**

**// D为度数矩阵,A为边矩阵**

**// 然后拉格朗日插值求出系数即可**

LL guass(LL A[107][107],int n,LL M) {

LL ret=1; int i,j,k;

REP(i,n) {

int id=i;

**if** (!A[i][i]) rep(j,i+1,n) **if** (A[j][i]) id=j;

**if** (!A[id][i]) **continue**;

**if** (id!=i) {rep(j,i,n) swap(A[i][j],A[id][j]); ret\*=-1;}

A[i][i]%=M; (A[i][i]<0) &&(A[i][i]+=M);

LL rev=poww(A[i][i],M-2,M);

rep(k,i+1,n) rrep(j,i,n)

(A[k][j]-=(LL)A[k][i]\*rev%M\*A[i][j])%=M;

} REP(i,n)(ret\*=A[i][i])%=M;

(ret<0) &&(ret+=M);

**return** ret;

}

int n,m;

int i,j,k;

int a[107][107]; LL A[107][107];

LL val[107],v\_v[107];

LL f[107],g[107],ans[107];

int main() {

scanf("%d",&n);

FOR(i,1,n-1) {

int u,v;

scanf("%d%d",&u,&v); u--; v--;

a[u][v]=a[v][u]=1;

} REP(i,n) v\_v[i]=i;

REP(k,n) {

REP(i,n) REP(j,n) A[i][j]=0;

REP(i,n) REP(j,n) **if** (i!=j) {

**if** (a[i][j]) A[i][j]=M-v\_v[k],A[i][i]+=v\_v[k];

**else** A[i][j]=M-1,A[i][i]++;

} val[k]=guass(A,n-1,M);

}

g[0]=1; REP(i,n) rFOR(j,0,i)(g[j+1]+=g[j])%=M,(g[j]\*=(M-v\_v[i]))%=M;

REP(k,n) {

LL rev=1;

rFOR(i,0,n) f[i]=(g[i+1]+f[i+1]\*v\_v[k]%M+M)%M;

REP(j,n) **if** (j!=k)(rev\*=(v\_v[k]-v\_v[j]))%=M;

(rev<0) &&(rev+=M); rev=powMM(rev,M-2);

rev=(rev\*val[k])%M;

FOR(i,0,n)(ans[i]+=(LL)f[i]\*rev%M)%=M;

} FOR(i,0,n-1) printf("%lld ",ans[i]);

}

Polya定理| Burnside引理

**//HDU3923; 颜色m, 个数n, 翻转或者置换当成一种**

**//ans=1/|G|\*sigma{pow(k(color),m(not move point 不动点数))}**

**//注意特殊形式**

**//Burnside引理:等价类个数l=sum{ci(ai)},ci是置换下的不动点数**

**//这个pow是可以变化成其他形式的**

**//注意,polya定理相当于手动算了一下Burnside引理中不动点的个数!**

int n,m;

bool mark[maxn];

int phi[maxn];

int p[maxn],tot;

int main() {

int i,j;

phi[1]=1;

FOR(i,2,1000000) {

**if** (!mark[i]) p[tot++]=i,phi[i]=i-1;

REP(j,tot) {

**if** (i\*p[j]>1000000) **break**;

*//感觉上不会爆,因为是从小往筛的*

mark[i\*p[j]]=1;

**if** (i%p[j]==0) {phi[i\*p[j]]=phi[i]\*p[j]; **break**;}

**else** phi[i\*p[j]]=phi[i]\*(p[j]-1);

}

}

int t,T;

scanf("%d",&T);

FOR(t,1,T) {

scanf("%d%d",&m,&n);

LL all=0,cnt=0;

*// FOR(i,1,n){*

*// (all+=powMM((LL)m,gcd(n,i)))%=M;*

*// (all<0)&&(all+=M);*

*// }cnt=n;*

*//置换*

FOR(i,1,n) **if** (n%i==0) {

(all+=(LL)powMM(m,i)\*phi[n/i])%=M;

(all<0) &&(all+=M);

}

cnt=n;

*//翻转*

**if** (n&1) {

(all+=(LL)n\*powMM(m,(n+1)/2))%=M;

cnt+=n;

} **else** {

(all+=(LL)n/2\*powMM(m,n/2))%=M;

(all+=(LL)n/2\*powMM(m,n/2+1))%=M;

cnt+=n;

}

*// printf("%lld %lld\n",cnt,all);*

all=all\*powMM(cnt,M-2)%M;

printf("Case #%d: %lld**\n**",t,all);

}

}

Miller\_Rabin素性测试+pollard\_rho因数分解

/\*miller\_rabin\*/

**const** int times=8;*// random\_check; 8-12 is OK*

LL mul(LL a,LL b,LL M) {

LL ret=0;

**for** (; b; b>>=1,(a+=a)>=M&&(a-=M))

**if** (b&1)(ret+=a)>=M&&(ret-=M);

**return** ret;

}

LL poww(LL a,LL b,LL M) {

LL ret=1;

**for** (; b; b>>=1,a=mul(a,a,M))

**if** (b&1) ret=mul(ret,a,M);

**return** ret;

}

bool check(LL a,LL n,LL x,LL t) {

LL ret=poww(a,x,n);

LL last=ret;

**for** (ret=mul(ret,ret,n); t--; last=ret,ret=mul(ret,ret,n))

**if** (ret==1&&last!=1&&last!=n-1) **return** true;

**if** (ret!=1) **return** true;

**return** false;

}

bool miller\_rabin(LL n) {

**if** (n<2) **return** false;

**if** (!(n&1)) **return** (n==2);

LL x=n-1,t=0;

**while** (!(x&1)) x>>=1,t++;

int i;

REP(i,times)

**if** (check(rand()%(n-1)+1,n,x,t)) **return** false;

**return** true;

}

*/\*pollard\_rho\*/*

LL pollard\_rho(LL x,LL c) {

LL x0=rand()%(x-1)+1;

LL y=x0; c%=x;

**for** (LL i=2,k=2;; i++) {

((x0=mul(x0,x0,x)+c)>=x)&&(x0-=x);

LL d=gcd(y-x0+x,x);

**if** (d!=1&&d!=x) **return** d;

**if** (y==x0) **return** x;

**if** (i==k) y=x0,k+=k;

}

}

LL factor[107]; int tot;

void findfac(LL n,int k) {

**if** (n==1) **return**;

**if** (miller\_rabin(n)) {factor[tot++]=n; **return**;}

LL p=n;

int c=k;

**while** (p>=n) p=pollard\_rho(p,c--);

findfac(p,k);

findfac(n/p,k);

}

int main() {

int T;

srand(time(0));

scanf("%d",&T);

**while** (T--) {

LL n; int i;

scanf("%I64d",&n);

**if** (miller\_rabin(n)) puts("Prime");

**else** {

tot=0;

findfac(n,107);

LL ans=factor[0];

REP(i,tot) ans=min(ans,factor[i]);

printf("%I64d**\n**",ans);

}

}

}

中国剩余定理(不一定互质)

void exgcd(LL a,LL b,LL &d,LL &x,LL &y){

**if** (!b) {d=a;x=1;y=0;}

**else** {exgcd(b,a%b,d,y,x);y-=a/b\*x;}

}

int n,m;

int i,j,k;

vector<LL> P,O;

int ans;

LL chinese\_remainder(vector<LL> &m,vector<LL> &r){

int i;LL M=m[0],R=r[0];

rep(i,1,P.size()){

LL x,y,d;

exgcd(M,m[i],d,x,y);

**if** ((r[i]-R)%d) **return** -1;

x=(r[i]-R)/d\*x%(m[i]/d);

R+=x\*M;M=M/d\*m[i];

R%=M;(R<0)&&(R+=M);

}**return** R;

}

int main(){

**while** (~scanf("%d",&n)){

P.clear();O.clear();

REP(i,n){

LL k;

scanf("%lld",&k);P.push\_back(k);

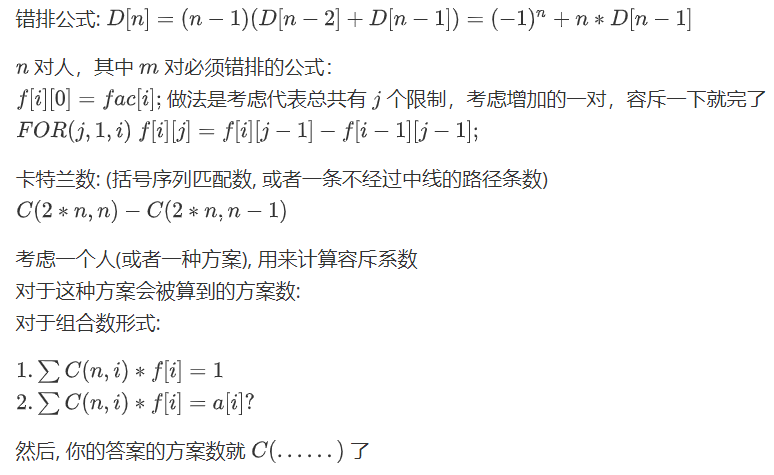
scanf("%lld",&k);O.push\_back(k);

}printf("%lld**\n**",chinese\_remainder(P,O));

}

}

广义容斥



Prime-counting function

**//这道题题意:小于n有多少个数字有4个因子**

**//(两个质数积,一个质数三次方)**

**//注意容斥减去多算的**

**//http://codeforces.com/blog/entry/44466?#comment-290036/**

**//考虑S(v,m):2...v,质因子全都>=m;那么考虑容斥:**

**//容斥掉的至少有一个p,而且没有小于p的因子**

**//很明显的,p=min(p,sqrt(v));**

**//S(v,p)=S(v,p-1)-(S(v/p,p-1)-S(p-1,p-1));(DP)**

**//那么反过来算即可;pi(n)=S(n,n);**

**//H[i]:pi(n/i);L[i]:pi(i)**

**//计算过程中,L[i]表示S(i,p),最终S(i,i)**

**//简单的这样DP,时间复杂度O(n^3/4),如果预处理n^2/3则最终n^2/3**

**//在后方,如果要容斥,FOR是很不方便的,感觉还是最好直接搞复杂度有保障**

Min\_25筛

**SPOJ DIVCNTK(sum\_ \sigma(i^k))**

**namespace** seives { *// 抄的define*

#define clr(ar) memset(ar, 0, sizeof(ar))

#define chkbit(ar, i) (((ar[(i) >> 6]) & (1 << (((i) >> 1) & 31))))

#define setbit(ar, i) (((ar[(i) >> 6]) |= (1 << (((i) >> 1) & 31))))

#define isprime(x) (( (x) && ((x)&1) && (!chkbit(ar, (x)))) || ((x) == 2))

**const** int MAXP=66666;

**const** int MAX=100010;*//euler\_seive*

**const** int maxn=100010;*//min\_25, =sqrt(n)*

int p[MAXP],tot;

ui ar[(MAX>>6)+7]= {0};

void init() {*//seives*

setbit(ar,0); setbit(ar,1);

int i,j; tot=0;

rep(i,2,MAX) {

**if** (isprime(i)) p[tot++]=i;

REP(j,tot) {

**if** (i\*p[j]>=MAX) **break**;

**if** ((i\*p[j])&1) setbit(ar,i\*p[j]);

**if** (i%p[j]==0) **break**;

}

}

}

**// 普通pcf公式: g(i,j)=g(i-1,j)-p^k\*g(i-1,j/p)**

**// 只有小于等于sqrt的p有用, 所以枚举这个, 考虑对其他答案的贡献**

**// 对于某个积性函数: (算贡献)**

**// g(i,j)=g(i-1,j)+\sum\_p^k F(p^k)\*g(i-1,j/[p^k]),还要加p^k的贡献**

**// 对于小于等于sqrt的p, 直接筛**

**// 对于大于的, 贡献只会是F(p)! 也就是...直接洲阁筛把答案的贡献加进去**

**// 这个加贡献=\_= 竟然是直接pcf求个前缀和啥的就完事了啊=\_=**

**// typedef ull ll;**

**// 注意如果想要去掉某个质数的贡献, 这个p[k]至少要筛到sqrtn...**

**// 注意F1的贡献, 是要乘的...**

**// 我这个F和G和一般的定义是反的...要先算G**

**// F和G定义是质数处的前缀和**

**// getans处记得如果质数贡献不同得改**

ll n,m;*//blocksize*

ll H[maxn],L[maxn];

void pcf() {

ll p,k;

FOR(p,1,m) L[p]=p-1,H[p]=n/p-1;

FOR(p,2,m) {

**if** (L[p]==L[p-1]) **continue**;*//not\_prime*

FOR(k,1,min(m,n/p/p)) {

**if** (p\*k<=m) H[k]-=H[p\*k]-L[p-1];

**else** H[k]-=L[n/p/k]-L[p-1];

} rFOR(k,p\*p,m) L[k]-=L[k/p]-L[p-1];

}

}

ll F[maxn],G[maxn];*//F[n/k]:H[n/k], G[i]:L[i]*

ll K;

ll getans(ll x,int i) {

**if** (x<=1||p[i]>x) **return** 0;

**if** (p[i]>m) **return** F[n/x]-G[m];

ll ans=((x<=m)?G[x]:F[n/x])-G[p[i]-1];

**for** (; (ll)p[i]\*p[i]<=x; i++) {

**for** (ll \_x=x/p[i],c=1; \_x>=p[i]; \_x/=p[i],c++)

ans+=getans(\_x,i+1)\*(c\*K+1)+((c+1)\*K+1);

} **return** ans;

}

ll solve() {

int p;

**for** (m=1; m\*m<=n; ++m); m--; pcf();

FOR(p,1,m) F[p]=H[p]\*(K+1),G[p]=L[p]\*(K+1);

**return** getans(n,0)+1;*//1:1*

}

}

**取模的质数个数筛的方式：**

ll n,m;*//blocksize*

ll H[maxn][4],L[maxn][4];

*// p%4==1那么p可以表示为两个平方因子的和*

void pcf() {

ll p,k; int i;

FOR(p,1,m) {

REP(i,4) L[p][i]=(p-i+4)/4;

REP(i,4) H[p][i]=(n/p-i+4)/4;

REP(i,2) H[p][i%4]--,L[p][i%4]--;

}

FOR(p,2,m) {

**if** (L[p][p%4]==L[p-1][p%4]) **continue**;*//not\_prime*

**static** int nxt[4];

REP(i,4) nxt[i]=i\*p%4;

FOR(k,1,min(m,n/p/p)) {

**if** (p\*k<=m) REP(i,4) H[k][nxt[i]]-=H[p\*k][i]-L[p-1][i];

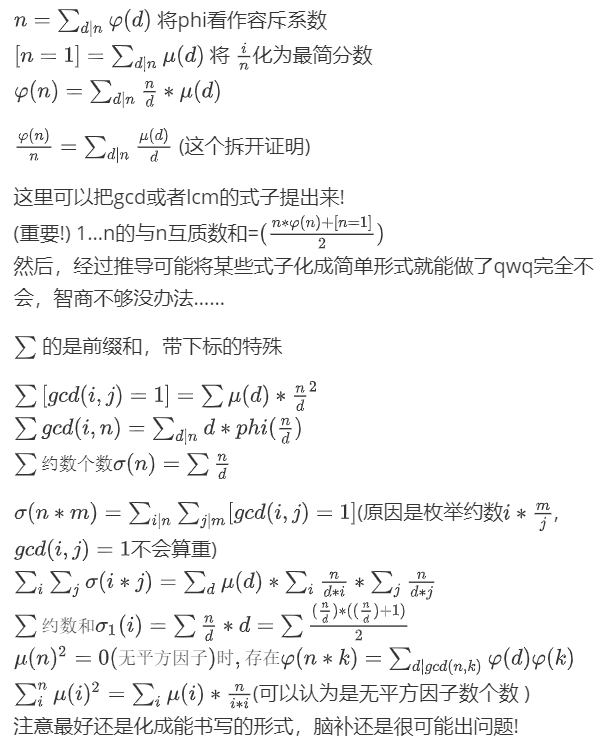
**else** REP(i,4) H[k][nxt[i]]-=L[n/p/k][i]-L[p-1][i];

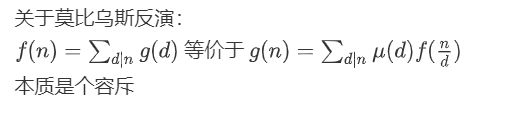
} rFOR(k,p\*p,m) REP(i,4) L[k][nxt[i]]-=L[k/p][i]-L[p-1][i];

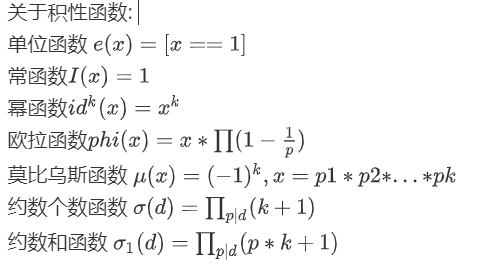
}

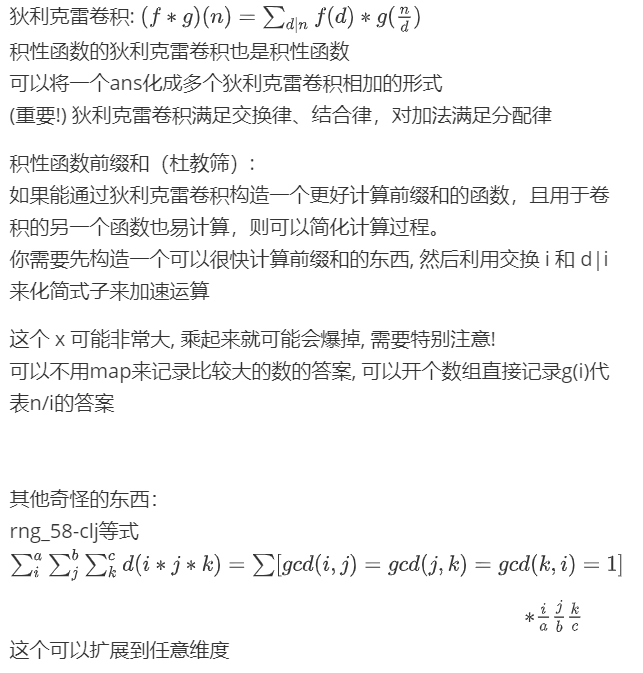
}

积性函数 前缀和 杜教筛

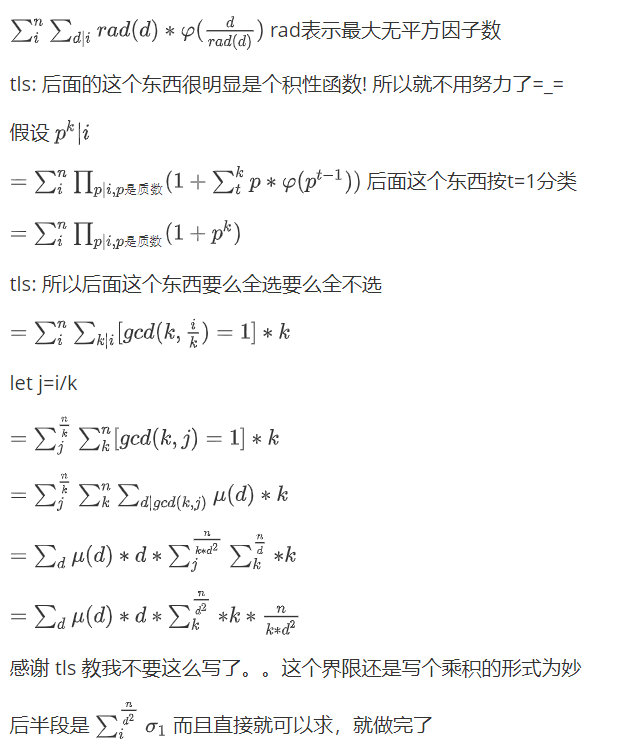








Zoj 3881



vector<int> P[maxn];

**namespace** seives { *// 抄的define*

#define clr(ar) memset(ar, 0, sizeof(ar))

#define chkbit(ar, i) (((ar[(i) >> 6]) & (1 << (((i) >> 1) & 31))))

#define setbit(ar, i) (((ar[(i) >> 6]) |= (1 << (((i) >> 1) & 31))))

#define isprime(x) (( (x) && ((x)&1) && (!chkbit(ar, (x)))) || ((x) == 2))

**const** int MAXP=666666;

**const** int MAX=2000010;

int mu[MAX],sigma1[MAX],c1[MAX],f[MAX];

int p[MAXP],tot;

ui ar[(MAX>>6)+7]= {0};

void init() {

setbit(ar,0); setbit(ar,1);

int i,j; tot=0; mu[1]=1; sigma1[1]=1;

rep(i,2,MAX) {

**if** (isprime(i)) p[tot++]=i,mu[i]=-1,sigma1[i]=i+1,c1[i]=i+1;

REP(j,tot) {

**if** (i\*p[j]>=MAX) **break**;

**if** ((i\*p[j])&1) setbit(ar,i\*p[j]);

**if** (i%p[j]==0) {

c1[i\*p[j]]=p[j]\*c1[i]+1;

sigma1[i\*p[j]]=sigma1[i]/c1[i]\*c1[i\*p[j]];

**break**;

} **else** {

c1[i\*p[j]]=p[j]+1;

sigma1[i\*p[j]]=sigma1[i]\*(p[j]+1);

mu[i\*p[j]]=-mu[i];

}

}

} rep(i,1,MAX) f[i]=sigma1[i],add\_(f[i],f[i-1]);

}

map<int,int> HASH;

int get2(ll x){

x%=M; **return** x\*(x+1)%M\*500000005%M;

}

int get\_f(ll x){*//直接sqrt也行*

**if** (x<MAX) **return** f[x];

**if** (HASH.count(x)) **return** HASH[x];

ll ret=0; ll l;

FOR(l,1,x) {

ll t=x/l,r=x/t;

add\_(ret,(get2(r)-get2(l-1)+M)%M\*(t%M)%M);

l=r;

} **return** HASH[x]=ret;

}

}

int main() {

*// startTimer();*

seives::init(); ll n;

*// printTimer();*

**while** (~scanf("%lld",&n)){

*// startTimer();*

int ans=0;

**for** (ll d=1;d\*d<=n;d++) add\_(ans,(M+seives::mu[d])\*d%M\*seives::get\_f(n/d/d)%M);

printf("%d**\n**",ans);

*// printTimer();*

}

}

类欧几里得

**一定注意前面是a,后面是b,线段树一定要注意顺序**

**f(a,b,c,n)=sigma{(ai+b)/c}; (0->n)**

**g(a,b,c,n)=sigma{(ai+b)/c\*i}; (0->n)**

**h(a,b,c,n)=sigma{((ai+b)/c)^2}; (0->n)**

**let m=(a\*n+b)/c;**

**推导f:**

**a=0:**

**return b/c\*(n+1)**

**a>=c||b>=c:有一部分是规律的;**

**return (a/c)\*n(n+1)/2+(b/c)\*(n+1)+f(a%c,b%c,c,n)**

**else:直接算,这个东西是个梯形中的点数,反过来算就可以了**

**f(a,b,c,n)=∑i=0->n ∑j=0->m-1 [(ai+b)/c>=j+1]**

**f(a,b,c,n)=∑i=0->n ∑j=0->m-1 [ai>=cj+c−b]**

**f(a,b,c,n)=∑i=0->n ∑j=0->m-1 [ai>cj+c−b−1]**

**f(a,b,c,n)=∑i=0->n ∑j=0->m-1 [i>(cj+c−b−1)/a]**

**f(a,b,c,n)=∑j=0->m (n−(cj+c−b−1)/a)**

**f(a,b,c,n)=n\*m-f(c,c-b-1,a,m-1);**

**推导g:**

**a=0:**

**return b/c\*n(n+1)/2 (sigma的是i)**

**a>=c||b>=c:有一部分是规律的;**

**g(a,b,c,n)=(a/c)\*n(n+1)(2n+1)/6+(b/c)\*n(n+1)/2+g(a%c,b%c,c,n)**

**else:**

**g(a,b,c,n)=∑i=0->n i\*∑j=0->m [(ai+b)/c>=j]**

**g(a,b,c,n)=∑i=0->n i\*∑j=0->m-1 [i>(cj+c−b−1)/a]**

**然后把这个i放进去求和**

**g(a,b,c,n)=1/2\*∑j=0->m-1 (n+1+(cj+c−b−1)/a)\*(n−(cj+c−b−1)/a)**

**g(a,b,c,n)=1/2\*∑j=0->m-1 n(n+1)−(cj+c−b−1)/a−[(cj+c−b−1)/a]^2**

**g(a,b,c,n)=1/2\*[n(n+1)\*m−f(c,c−b−1,a,m−1)−h(c,c−b−1,a,m−1)]**

**推导h:**

**a=0:**

**return (b/c)^2\*(n+1) (sigma的是i)**

**a>=c||b>=c:有一部分是规律的;**

**h(a,b,c,n)=(a/c)^2\*n(n+1)(2n+1)/6+(b/c)^2\*(n+1)+(a/c)\*(b/c)\*n(n+1)**

**+h(a%c,b%c,c,n)+2\*(a/c)\*g(a%c,b%c,c,n)+2\*(b/c)\*f(a%c,b%c,c,n)**

**else:**

**n^2=2\*n(n+1)/2−n=2(∑i=0->n i)−n**

**有了思路我们来推h**

**h(a,b,c,n)=∑i=0->n (2(∑j=1->(ai+b)/c j)−(ai+b)/c)**

**可以想到交换主体。**

**h(a,b,c,n)=∑j=0->m-1 (j+1)\*∑i=0->n [(ai+b)/c>=j+1]−f(a,b,c,n)**

**h(a,b,c,n)=∑j=0->m-1 (j+1)\*∑i=0->n [i>(cj+c−b−1)/a]−f(a,b,c,n)**

**h(a,b,c,n)=∑j=0->m-1 (j+1)\*(n−(cj+c−b−1)/a)−f(a,b,c,n)**

**h(a,b,c,n)=n\*m(m+1)−2g(c,c−b−1,a,m−1)−2f(c,c−b−1,a,m−1)−f(a,b,c,n)**

**// 题意:n%1,n%2...异或, 做法是BSGS然后类欧几里得**

**// 每块是n-n/l\*l ^ ... ^ n-n/r\*r**

**// 也就是n-(n/l)\*k,(等价于n%r+(n/l)\*k) k是0->r-l**

**// 按位计算, 就变成了个类欧几里得**

**// 玄学卡常,n<=某值直接暴力, 这里tls说是一个log的**

LL f(LL a,LL b,LL c,LL n) {

**if** (a==0) **return** b/c\*(n+1);

**if** (a>=c||b>=c) **return** (a/c)\*n\*(n+1)/2+(b/c)\*(n+1)+f(a%c,b%c,c,n);

LL m=(a\*n+b)/c;

**return** n\*m-f(c,c-b-1,a,m-1);

}

LL solve(LL l,LL c,LL n) {

LL ret=0,i;

**if** (n<=10000) REP(i,n+1) ret^=l,l+=c;

**else** REP(i,40) ret^=(f(c,l,(1ll<<i),n)&1)<<i;

**return** ret;

}

LL getans(LL n) {

LL ans=0;

**for** (LL l=1,r; l<=n;) {

r=n/(n/l);

ans^=solve(n%r,n/l,r-l);

l=r+1;

} **return** ans;

}

int main() {

int T;

int i,j,k;

scanf("%d",&T);

**while** (T--) {

LL n;

scanf("%lld",&n);

printf("%lld**\n**",getans(n));

}

**return** 0;

}

欧拉降幂公式

**//n^x(mod m)=m^(phi(m)+x%phi(m))%m (x>m)**

**//这个题让求pow(l,pow(l+1...pow(r)))**

inline int mod(LL a,int b){

if (a<b) return a;

return a%b+b;

}

inline int poww(int a,int b,int M){

int ret=1;

for (;b;b>>=1ll,a=mod(1ll\*a\*a,M))

if (b&1) ret=mod(1ll\*ret\*a,M);

return ret;

}

typedef pair<int,int> pii;

int P[maxn];

int phi(int x){

int k=x;

for (int i=2;i\*i<=k;i++) if (k%i==0){

x=x/i\*(i-1);

while (k%i==0) k/=i;

}if (k!=1) x=x/k\*(k-1);

return x;

}

int a[maxn];

int tot;

int solve(int l,int r,int pos){

if (l==r||pos==tot) return mod(a[l],P[pos]);

return poww(a[l],solve(l+1,r,pos+1),P[pos]);

}

int n,m,q,i,j,k;

int main(){

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%d",&a[i]);

P[1]=m;

for (tot=1;P[tot]!=1;tot++) P[tot+1]=phi(P[tot]);

scanf("%d",&q);

FOR(i,1,q){

int l,r;int ans=1;

scanf("%d%d",&l,&r);

printf("%d\n",solve(l,r,1)%m);

}

}

单纯形法

注意，只能做实数规划

输出方案：

**//http://www.voidcn.com/article/p-kkqovyic-qh.html**

**//m个式子,n个变量**

**//maximize A[0][i]\*value**

**//sigma<=A[i][0]**

**namespace** Simplex {

int n,m;*//n变量, m式子*

**const** int maxn=500,maxm=5000;

double A[maxm][maxn];*//*

int id[maxn+maxm];*//base*

**const** double inf=1e20;

**const** double eps=1e-7;

void pivot(int l,int e) {

int tt=id[n+l]; id[n+l]=id[e]; id[e]=tt;

int i,j; double t=A[l][e]; A[l][e]=1;

FOR(j,0,n) A[l][j]/=t;

FOR(i,0,m) **if** (i!=l && abs(A[i][e])>eps) {

t=A[i][e]; A[i][e]=0;

**for** (j=0; j<=n; j++)

A[i][j]-=A[l][j]\*t;

}

}

bool initialize() {

int i,j;

FOR(i,1,n) id[i]=i;

**while** (1) {

int e=0,l=0;

FOR(i,1,m) **if** (A[i][0]<-eps && (!l || (rand()&1))) l=i;

**if** (!l) **break**;

FOR(j,1,n) **if** (A[l][j]<-eps && (!e || (rand()&1))) e=j;

**if** (!e) **return** 0;*//Infeasible,无解*

pivot(l,e);

} **return** 1;

}

double ans[maxn],value;

bool simplex() {

int i,j;

**while** (true) {

int l=0,e=0; double minn=inf;

FOR(j,1,n) **if** (A[0][j]>eps) {e=j; **break**;}

**if** (!e) **break**;

FOR(i,1,m) **if** (A[i][e]>eps && A[i][0]/A[i][e]<minn)

minn=A[i][0]/A[i][e],l=i;

**if** (!l) **return** 0;*//Unbounded,inf*

pivot(l,e);

}

FOR(i,1,m) ans[id[n+i]]=A[i][0];

value=-A[0][0];*//maxvalue*

**return** 1;

}

}

int main() {

int n,m,i,j;

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%lf",&Simplex::A[0][i]);

FOR(i,1,m) {

int l,r,c;

scanf("%d%d%d",&l,&r,&c);

FOR(j,l,r) Simplex::A[i][j]=1;

Simplex::A[i][0]=c;

}

Simplex::n=n; Simplex::m=m;

assert(Simplex::initialize());

assert(Simplex::simplex());

*// FOR(i,1,n) printf("%.0f ",Simplex::ans[i]); puts("");*

printf("%.0f",Simplex::value);

}

不输出方案，另一种写法：

**//m个式子,n个变量**

**//maximize C[i]\*value**

**namespace** Simplex {

int n,m;*//n变量, m式子*

**const** int maxn=500,maxm=5000;

double A[maxm][maxn],B[maxm];

double C[maxn];*//base*

double v=0;

**const** double inf=1e20;

**const** double eps=1e-7;

void pivot(int l,int e) {

B[l]/=A[l][e];

**for** (int i=1; i<=n; i++)

**if** (i!=e) A[l][i]/=A[l][e];

A[l][e]=1/A[l][e];

**for** (int i=1; i<=m; i++)

**if** (i!=l&&fabs(A[i][e])>eps) {

B[i]-=B[l]\*A[i][e];

**for** (int j=1; j<=n; j++)

**if** (j!=e) A[i][j]-=A[l][j]\*A[i][e];

A[i][e]=-A[l][e]\*A[i][e];

}

v+=C[e]\*B[l];

**for** (int i=1; i<=n; i++)

**if** (i!=e) C[i]-=C[e]\*A[l][i];

C[e]=-C[e]\*A[l][e];

}

double simplex() {

int l,e;

double t;

**while** (true) {

e=n+1;

**for** (int i=1; i<=n; i++)

**if** (C[i]>eps) {e=i; **break**;}

**if** (e==n+1) **break**;

t=inf; l=0;

**for** (int i=1; i<=m; i++)

**if** (A[i][e]>eps&&t>B[i]/A[i][e]) {

t=B[i]/A[i][e]; l=i;

}

**if** (t==inf) **return** inf;

pivot(l,e);

}

**return** v;

}

}

int main() {

int n,m,i,j;

scanf("%d%d",&n,&m);

FOR(i,1,n) scanf("%lf",&Simplex::C[i]);

FOR(i,1,m) {

int l,r,c;

scanf("%d%d%d",&l,&r,&c);

FOR(j,l,r) Simplex::A[i][j]=1;

Simplex::B[i]=c;

}

Simplex::n=n; Simplex::m=m;

printf("%.0f",Simplex::simplex());

}

# 其他的东西

杜教线性递推BM板子

int \_,n;

**namespace** linear\_seq {

**const** int N=10010;

ll res[N],base[N],\_c[N],\_md[N];

vector<int> Md;

void mul(ll \*a,ll \*b,int k) {

rep(i,0,k+k) \_c[i]=0;

rep(i,0,k) **if** (a[i]) rep(j,0,k) \_c[i+j]=(\_c[i+j]+a[i]\*b[j])%mod;

**for** (int i=k+k-1;i>=k;i--) **if** (\_c[i])

rep(j,0,SZ(Md)) \_c[i-k+Md[j]]=(\_c[i-k+Md[j]]-\_c[i]\*\_md[Md[j]])%mod;

rep(i,0,k) a[i]=\_c[i];

}

int solve(ll n,VI a,VI b) { *// a 系数 b 初值 b[n+1]=a[0]\*b[n]+...*

*// printf("%d\n",SZ(b));*

ll ans=0,pnt=0;

int k=SZ(a);

assert(SZ(a)==SZ(b));

rep(i,0,k) \_md[k-1-i]=-a[i];\_md[k]=1;

Md.clear();

rep(i,0,k) **if** (\_md[i]!=0) Md.push\_back(i);

rep(i,0,k) res[i]=base[i]=0;

res[0]=1;

**while** ((1ll<<pnt)<=n) pnt++;

**for** (int p=pnt;p>=0;p--) {

mul(res,res,k);

**if** ((n>>p)&1) {

**for** (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;

rep(j,0,SZ(Md)) res[Md[j]]=(res[Md[j]]-res[k]\*\_md[Md[j]])%mod;

}

}

rep(i,0,k) ans=(ans+res[i]\*b[i])%mod;

**if** (ans<0) ans+=mod;

**return** ans;

}

VI BM(VI s) {

VI C(1,1),B(1,1);

int L=0,m=1,b=1;

rep(n,0,SZ(s)) {

ll d=0;

rep(i,0,L+1) d=(d+(ll)C[i]\*s[n-i])%mod;

**if** (d==0) ++m;

**else** **if** (2\*L<=n) {

VI T=C;

ll c=mod-d\*powmod(b,mod-2)%mod;

**while** (SZ(C)<SZ(B)+m) C.pb(0);

rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c\*B[i])%mod;

L=n+1-L; B=T; b=d; m=1;

} **else** {

ll c=mod-d\*powmod(b,mod-2)%mod;

**while** (SZ(C)<SZ(B)+m) C.pb(0);

rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c\*B[i])%mod;

++m;

}

}

**return** C;

}

int gao(VI a,ll n) {

VI c=BM(a);

c.erase(c.begin());

rep(i,0,SZ(c)) c[i]=(mod-c[i])%mod;

**for** (int v:c) printf("%d ",v);puts("");

**return** solve(n,c,VI(a.begin(),a.begin()+SZ(c)));

}

};

int main() {

int k=linear\_seq::gao(VI{7,16,25,50,84,159,277,511,906,1651,2952,5348,9601,17345,31199,56288,101341},10);

printf("%d**\n**",k);

**for** (scanf("%d",&\_);\_;\_--) {

scanf("%d",&n);

printf("%d**\n**",linear\_seq::gao(VI{0,1,1,2,3,5,8,13,21,34},n-1));

}

}

任意模数BM板子

#include <bits/stdc++.h>

**using** **namespace** std;

#ifndef ONLINE\_JUDGE

#define debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \_\_func\_\_, ##\_\_VA\_ARGS\_\_)

#else

#define debug(...)

#endif

*// given first m items init[0..m-1] and coefficents trans[0..m-1] or*

*// given first 2 \*m items init[0..2m-1], it will compute trans[0..m-1]*

*// for you. trans[0..m] should be given as that*

*// init[m] = sum\_{i=0}^{m-1} init[i] \* trans[i]*

**struct** LinearRecurrence {

**using** int64 = long long;

**using** vec = std::vector<int64>;

**static** void extand(vec &a, size\_t d, int64 value = 0) {

**if** (d <= a.size()) **return**;

a.resize(d, value);

}

**static** vec BerlekampMassey(**const** vec &s, int64 mod) {

std::function<int64(int64)> inverse = [&](int64 a) {

**return** a == 1 ? 1 : (int64)(mod - mod / a) \* inverse(mod % a) % mod;

};

vec A = {1}, B = {1};

int64 b = s[0];

**for** (size\_t i = 1, m = 1; i < s.size(); ++i, m++) {

int64 d = 0;

**for** (size\_t j = 0; j < A.size(); ++j) {

d += A[j] \* s[i - j] % mod;

}

**if** (!(d %= mod)) **continue**;

**if** (2 \* (A.size() - 1) <= i) {

**auto** temp = A;

extand(A, B.size() + m);

int64 coef = d \* inverse(b) % mod;

**for** (size\_t j = 0; j < B.size(); ++j) {

A[j + m] -= coef \* B[j] % mod;

**if** (A[j + m] < 0) A[j + m] += mod;

}

B = temp, b = d, m = 0;

} **else** {

extand(A, B.size() + m);

int64 coef = d \* inverse(b) % mod;

**for** (size\_t j = 0; j < B.size(); ++j) {

A[j + m] -= coef \* B[j] % mod;

**if** (A[j + m] < 0) A[j + m] += mod;

}

}

}

**return** A;

}

**static** void exgcd(int64 a, int64 b, int64 &g, int64 &x, int64 &y) {

**if** (!b)

x = 1, y = 0, g = a;

**else** {

exgcd(b, a % b, g, y, x);

y -= x \* (a / b);

}

}

**static** int64 crt(**const** vec &c, **const** vec &m) {

int n = c.size();

int64 M = 1, ans = 0;

**for** (int i = 0; i < n; ++i) M \*= m[i];

**for** (int i = 0; i < n; ++i) {

int64 x, y, g, tm = M / m[i];

exgcd(tm, m[i], g, x, y);

ans = (ans + tm \* x \* c[i] % M) % M;

}

**return** (ans + M) % M;

}

**static** vec ReedsSloane(**const** vec &s, int64 mod) {

**auto** inverse = [](int64 a, int64 m) {

int64 d, x, y;

exgcd(a, m, d, x, y);

**return** d == 1 ? (x % m + m) % m : -1;

};

**auto** L = [](**const** vec& a, **const** vec& b) {

int da = (a.size() > 1 || (a.size() == 1 && a[0])) ? a.size() - 1 : -1000;

int db = (b.size() > 1 || (b.size() == 1 && b[0])) ? b.size() - 1 : -1000;

**return** std::max(da, db + 1);

};

**auto** prime\_power = [&](**const** vec& s, int64 mod, int64 p, int64 e) {

*// linear feedback shift register mod p^e, p is prime*

std::vector<vec> a(e), b(e), an(e), bn(e), ao(e), bo(e);

vec t(e), u(e), r(e), to(e, 1), uo(e), pw(e + 1);

;

pw[0] = 1;

**for** (int i = pw[0] = 1; i <= e; ++i) pw[i] = pw[i - 1] \* p;

**for** (int64 i = 0; i < e; ++i) {

a[i] = {pw[i]}, an[i] = {pw[i]};

b[i] = {0}, bn[i] = {s[0] \* pw[i] % mod};

t[i] = s[0] \* pw[i] % mod;

**if** (t[i] == 0) {

t[i] = 1, u[i] = e;

} **else** {

**for** (u[i] = 0; t[i] % p == 0; t[i] /= p, ++u[i])

;

}

}

**for** (size\_t k = 1; k < s.size(); ++k) {

**for** (int g = 0; g < e; ++g) {

**if** (L(an[g], bn[g]) > L(a[g], b[g])) {

ao[g] = a[e - 1 - u[g]];

bo[g] = b[e - 1 - u[g]];

to[g] = t[e - 1 - u[g]];

uo[g] = u[e - 1 - u[g]];

r[g] = k - 1;

}

}

a = an, b = bn;

**for** (int o = 0; o < e; ++o) {

int64 d = 0;

**for** (size\_t i = 0; i < a[o].size() && i <= k; ++i) {

d = (d + a[o][i] \* s[k - i]) % mod;

}

**if** (d == 0) {

t[o] = 1, u[o] = e;

} **else** {

**for** (u[o] = 0, t[o] = d; t[o] % p == 0; t[o] /= p, ++u[o])

;

int g = e - 1 - u[o];

**if** (L(a[g], b[g]) == 0) {

extand(bn[o], k + 1);

bn[o][k] = (bn[o][k] + d) % mod;

} **else** {

int64 coef = t[o] \* inverse(to[g], mod) % mod \* pw[u[o] - uo[g]] % mod;

int m = k - r[g];

extand(an[o], ao[g].size() + m);

extand(bn[o], bo[g].size() + m);

**for** (size\_t i = 0; i < ao[g].size(); ++i) {

an[o][i + m] -= coef \* ao[g][i] % mod;

**if** (an[o][i + m] < 0) an[o][i + m] += mod;

}

**while** (an[o].size() && an[o].back() == 0) an[o].pop\_back();

**for** (size\_t i = 0; i < bo[g].size(); ++i) {

bn[o][i + m] -= coef \* bo[g][i] % mod;

**if** (bn[o][i + m] < 0) bn[o][i + m] -= mod;

}

**while** (bn[o].size() && bn[o].back() == 0) bn[o].pop\_back();

}

}

}

}

**return** std::make\_pair(an[0], bn[0]);

};

std::vector<std::tuple<int64, int64, int>> fac;

**for** (int64 i = 2; i \* i <= mod; ++i) {

**if** (mod % i == 0) {

int64 cnt = 0, pw = 1;

**while** (mod % i == 0) mod /= i, ++cnt, pw \*= i;

fac.emplace\_back(pw, i, cnt);

}

}

**if** (mod > 1) fac.emplace\_back(mod, mod, 1);

std::vector<vec> as;

size\_t n = 0;

**for** (**auto**&& x : fac) {

int64 mod, p, e;

vec a, b;

std::tie(mod, p, e) = x;

**auto** ss = s;

**for** (**auto**&& x : ss) x %= mod;

std::tie(a, b) = prime\_power(ss, mod, p, e);

as.emplace\_back(a);

n = std::max(n, a.size());

}

vec a(n), c(as.size()), m(as.size());

**for** (size\_t i = 0; i < n; ++i) {

**for** (size\_t j = 0; j < as.size(); ++j) {

m[j] = std::get<0>(fac[j]);

c[j] = i < as[j].size() ? as[j][i] : 0;

}

a[i] = crt(c, m);

}

**return** a;

}

LinearRecurrence(**const** vec &s, **const** vec &c, int64 mod) : init(s), trans(c), mod(mod), m(s.size()) {}

LinearRecurrence(**const** vec &s, int64 mod, bool is\_prime = true) : mod(mod) {

vec A;

**if** (is\_prime)

A = BerlekampMassey(s, mod);

**else**

A = ReedsSloane(s, mod);

**if** (A.empty()) A = {0};

m = A.size() - 1;

trans.resize(m);

**for** (int i = 0; i < m; ++i) {

trans[i] = (mod - A[i + 1]) % mod;

}

std::reverse(trans.begin(), trans.end());

init = {s.begin(), s.begin() + m};

}

int64 calc(int64 n) {

**if** (mod == 1) **return** 0;

**if** (n < m) **return** init[n];

vec v(m), u(m << 1);

int msk = !!n;

**for** (int64 m = n; m > 1; m >>= 1) msk <<= 1;

v[0] = 1 % mod;

**for** (int x = 0; msk; msk >>= 1, x <<= 1) {

std::fill\_n(u.begin(), m \* 2, 0);

x |= !!(n & msk);

**if** (x < m)

u[x] = 1 % mod;

**else** {

*// can be optimized by fft/ntt*

**for** (int i = 0; i < m; ++i) {

**for** (int j = 0, t = i + (x & 1); j < m; ++j, ++t) {

u[t] = (u[t] + v[i] \* v[j]) % mod;

}

}

**for** (int i = m \* 2 - 1; i >= m; --i) {

**for** (int j = 0, t = i - m; j < m; ++j, ++t) {

u[t] = (u[t] + trans[j] \* u[i]) % mod;

}

}

}

v = {u.begin(), u.begin() + m};

}

int64 ret = 0;

**for** (int i = 0; i < m; ++i) {

ret = (ret + v[i] \* init[i]) % mod;

}

**return** ret;

}

vec init, trans;

int64 mod;

int m;

};

**const** int mod = 1e9;

**typedef** long long ll;

ll Pow(ll a, ll n, ll mod) {

ll t = 1;

**for** (; n; n >>= 1, (a \*= a) %= mod)

**if** (n & 1)(t \*= a) %= mod;

**return** t;

}

int main() {

int n, m;

cin >> n >> m;

std::vector<long long> f = {0, 1};

**for** (int i = 2; i < m \* 2 + 5; i++)

f.push\_back((f[i - 1] + f[i - 2]) % mod);

**for** (**auto** &t : f) t = Pow(t, m, mod);

**for** (int i = 1; i < m \* 2 + 5; i++)

f[i] = (f[i - 1] + f[i]) % mod;

LinearRecurrence solver(f, mod, false);

printf("%lld**\n**", solver.calc(n));

}

自适应 simpson积分

double simpson(double a,double b) {

double c = a + (b-a)/2;

return (F(a) + 4\*F(c) + F(b))\*(b-a)/6;

}

double asr(double a,double b,double eps,double A) {

double c = a + (b-a)/2;

double L = simpson(a,c), R = simpson(c,b);

if (fabs(L + R - A) <= 15\*eps)

return L + R + (L + R - A)/15.0;

return asr(a,c,eps/2,L) + asr(c,b,eps/2,R);

}

double asr(double a,double b,double eps) {

return asr(a,b,eps,simpson(a,b));

}

杜教多项式插值

#include<stdio.h>

#include<string.h>

#include<algorithm>

#include<assert.h>

**using** **namespace** std;

**typedef** long long ll;

**const** int mod = 1e9+7;

**namespace** polysum {

#define rep(i,a,n) for (int i=a;i<n;i++)

#define per(i,a,n) for (int i=n-1;i>=a;i--)

**const** int D=2010;*//最高幂次*

ll a[D],f[D],g[D],p[D],p1[D],p2[D],b[D],h[D][2],C[D];

ll powmod(ll a,ll b){ll res=1;a%=mod;assert(b>=0);**for**(;b;b>>=1){**if**(b&1)res=res\*a%mod;a=a\*a%mod;}**return** res;}

ll calcn(int d,ll \*a,ll n) { *// a[0].. a[d] a[n]*

**if** (n<=d) **return** a[n];

p1[0]=p2[0]=1;

rep(i,0,d+1) {

ll t=(n-i+mod)%mod;

p1[i+1]=p1[i]\*t%mod;

}

rep(i,0,d+1) {

ll t=(n-d+i+mod)%mod;

p2[i+1]=p2[i]\*t%mod;

}

ll ans=0;

rep(i,0,d+1) {

ll t=g[i]\*g[d-i]%mod\*p1[i]%mod\*p2[d-i]%mod\*a[i]%mod;

**if** ((d-i)&1) ans=(ans-t+mod)%mod;

**else** ans=(ans+t)%mod;

}

**return** ans;

}

void init(int M) {*//最高幂次*

f[0]=f[1]=g[0]=g[1]=1;

rep(i,2,M+5) f[i]=f[i-1]\*i%mod;

g[M+4]=powmod(f[M+4],mod-2);

per(i,1,M+4) g[i]=g[i+1]\*(i+1)%mod;

}

ll polysum(ll m,ll \*a,ll n) { *// a[0].. a[m] \sum\_{i=0}^{n-1} a[i]*

ll b[D];

**for**(int i=0;i<=m;i++) b[i]=a[i];

b[m+1]=calcn(m,b,m+1);

rep(i,1,m+2) b[i]=(b[i-1]+b[i])%mod;

**return** calcn(m+1,b,n-1);

}

ll qpolysum(ll R,ll n,ll \*a,ll m) { *// a[0].. a[m] \sum\_{i=0}^{n-1} a[i]\*R^i*

**if** (R==1) **return** polysum(n,a,m);

a[m+1]=calcn(m,a,m+1);

ll r=powmod(R,mod-2),p3=0,p4=0,c,ans;

h[0][0]=0;h[0][1]=1;

rep(i,1,m+2) {

h[i][0]=(h[i-1][0]+a[i-1])\*r%mod;

h[i][1]=h[i-1][1]\*r%mod;

}

rep(i,0,m+2) {

ll t=g[i]\*g[m+1-i]%mod;

**if** (i&1) p3=((p3-h[i][0]\*t)%mod+mod)%mod,p4=((p4-h[i][1]\*t)%mod+mod)%mod;

**else** p3=(p3+h[i][0]\*t)%mod,p4=(p4+h[i][1]\*t)%mod;

}

c=powmod(p4,mod-2)\*(mod-p3)%mod;

rep(i,0,m+2) h[i][0]=(h[i][0]+h[i][1]\*c)%mod;

rep(i,0,m+2) C[i]=h[i][0];

ans=(calcn(m,C,n)\*powmod(R,n)-c)%mod;

**if** (ans<0) ans+=mod;

**return** ans;

}

} *// polysum::init();*

求x^2+y^2=n的(x,y)对数

typedef long long ll;

const ll inf = 1e9+7;

const ll maxn = 2e5+7;

int solve(int n){

int sum=0;

for(int i=1;i\*i<=n;i++){

if(n%i==0){

if(i%4==1)sum++;

else if(i%4==3)sum--;

if(i\*i!=n){

if(n/i%4==1)sum++;

else if(n/i%4==3)sum--;

}

}

}

return sum\*4;

}

int solve2(int n){

while(n%2==0)n/=2;

int res=4;

for(int i=2;i\*i<=n;i++){

if(n%i==0){

int sum=0;

while(n%i==0)n/=i,sum++;

if(i%4==1)

res=res\*(sum+1);

else if(i%4==3&&sum%2==1)

return 0;

}

}

if(n>1){

if(n%4==1)

res=res\*2;

}

return res;

}